

NDMI113 – Extremal Combinatorics

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Lecture 2 - LYM inequality, set shifting.

Sperner's Theorem tells us that the size of a Sperner family $\mathcal{F}$ on $[n]$ is bounded by $\binom{n}{\lfloor n/2 \rfloor}$. The following generalization of this theorem, known as the LYM inequality, was proven independently by Bollobás (1965), Lubell (1966), Meshalkin (1963), and Yamamoto (1954).

Theorem 1 (LYM Inequality) *If $\mathcal{F} \subseteq 2^{[n]}$ is Sperner, then*

$$\sum_{k=0}^n \frac{|\mathcal{F}_k|}{\binom{n}{k}} \leq 1,$$

where $\mathcal{F}_k := \mathcal{F} \cap \binom{[n]}{k}$, i.e. $\mathcal{F}_k$ is the subset of $\mathcal{F}$ consisting of precisely the sets of size k .

We shall give two proofs of the LYM inequality. The first one is due to Frankl, and is an ingenious application of the probabilistic method.

Proof Let $\mathcal{C}$ be chain of length $n + 1$ picked uniformly at random (among the $n!$ choices). Let $\mathcal{C} = \{C_0, C_1, \dots, C_n\}$ where $\emptyset = C_0 \subsetneq C_1 \subsetneq \dots \subsetneq C_n = [n]$. We have

$$1 \geq \mathbb{P}[\mathcal{C} \cap \mathcal{F} \neq \emptyset] = \sum_{k=0}^n \mathbb{P}[\mathcal{C} \cap \mathcal{F}_k \neq \emptyset] = \sum_{k=0}^n \frac{|\mathcal{F}_k|}{\binom{n}{k}}.$$

Here the first equality comes from noting that the events $\mathcal{C} \cap \mathcal{F}_i \neq \emptyset$ and $\mathcal{C} \cap \mathcal{F}_j \neq \emptyset$ for $i \neq j$ are mutually exclusive, as a chain can intersect an antichain in at most one element. The second equality comes from

$$\mathbb{P}[\mathcal{C} \cap \mathcal{F}_k \neq \emptyset] = \frac{|\mathcal{F}_k|}{\binom{n}{k}},$$

since, by symmetry of $2^{[n]}$, the chain is equally likely to hit any given element of $\binom{[n]}{k}$. □

Our second proof of LYM requires the following fact known as the “Local LYM inequality”.

Definition 1 *Let $\mathcal{F} \subseteq \binom{[n]}{k}$ be a uniform set system. The shadow of $\mathcal{F}$ is defined as*

$$\partial\mathcal{F} := \{B \in \binom{[n]}{k-1} : B \subseteq A, \text{ for some } A \in \mathcal{F}\}.$$

Theorem 2 (Local LYM) *For any uniform set system $\mathcal{F} \subseteq \binom{[n]}{k}$ we have*

$$\frac{|\partial\mathcal{F}|}{\binom{n}{k-1}} \geq \frac{|\mathcal{F}|}{\binom{n}{k}},$$

with equality if and only if $\mathcal{F} = \binom{[n]}{k}$ or $\mathcal{F} = \emptyset$.

Proof We argue by double counting the edges between $\mathcal{F}$ and $\partial\mathcal{F}$ in the inclusion graph (see Lecture 1). Each set in $\mathcal{F}$ is a superset of precisely k sets in $\partial\mathcal{F}$, and thus there are exactly $k \cdot |\mathcal{F}|$ edges between $\mathcal{F}$ and $\partial\mathcal{F}$. On the other hand, each set in $\partial\mathcal{F}$ is a subset of at most $n - (k - 1)$ sets in $\mathcal{F}$ (it is a subset of exactly $n - (k - 1)$ sets in $\binom{[n]}{k}$, some of which may not be in $\mathcal{F}$), so the edge set has cardinality at most $(n - k + 1) \cdot |\partial\mathcal{F}|$. We obtain

$$|\partial\mathcal{F}| \geq |\mathcal{F}| \cdot \frac{k}{n - k + 1} = |\mathcal{F}| \frac{\binom{n}{k-1}}{\binom{n}{k}}.$$

Equality occurs if for all $A \in \mathcal{F}$, $a \in A$, and $x \notin A$ we have $(A \setminus \{a\}) \cup \{x\} \in \mathcal{F}$, which can only happen if $\mathcal{F} = \binom{[n]}{k}$ or $\mathcal{F} = \emptyset$. $\square$

Proof [of the LYM inequality.] We may assume that $\mathcal{F} \neq \emptyset$. We first prove the following claim by downward induction. Define $\mathcal{G}_n = \mathcal{F}_n$ and for each $\ell < n$ define $\mathcal{G}_\ell = \mathcal{F}_\ell \cup \partial\mathcal{G}_{\ell+1}$. We claim that for all $\ell \leq n$ we have

$$\frac{|\mathcal{G}_\ell|}{\binom{n}{\ell}} \geq \sum_{k=\ell}^n \frac{|\mathcal{F}_k|}{\binom{n}{k}}.$$

The base case $\ell = n$ is trivial. Now we suppose the claim has been proven for $\ell + 1$ and we shall prove it holds for ℓ . Notice first that $\mathcal{F}_\ell$ and $\partial\mathcal{G}_{\ell+1}$ are disjoint since each set in $\partial\mathcal{G}_{\ell+1}$ is contained in some element of $\mathcal{F}$, but $\mathcal{F}$ is an antichain. Therefore,

$$\frac{|\mathcal{G}_\ell|}{\binom{n}{\ell}} = \frac{|\mathcal{F}_\ell|}{\binom{n}{\ell}} + \frac{|\partial\mathcal{G}_{\ell+1}|}{\binom{n}{\ell}}.$$

Applying Local LYM yields

$$\frac{|\mathcal{G}_\ell|}{\binom{n}{\ell}} \geq \frac{|\mathcal{F}_\ell|}{\binom{n}{\ell}} + \frac{|\mathcal{G}_{\ell+1}|}{\binom{n}{\ell+1}}.$$

We now apply the induction hypothesis to the last term to get

$$\frac{|\mathcal{G}_\ell|}{\binom{n}{\ell}} \geq \frac{|\mathcal{F}_\ell|}{\binom{n}{\ell}} + \sum_{k=\ell+1}^n \frac{|\mathcal{F}_k|}{\binom{n}{k}} = \sum_{k=\ell}^n \frac{|\mathcal{F}_k|}{\binom{n}{k}}.$$

Thus, the claim is proven. Now consider the case $\ell = 0$ to get

$$1 = \frac{|\mathcal{G}_0|}{\binom{n}{0}} \geq \sum_{k=0}^n \frac{|\mathcal{F}_k|}{\binom{n}{k}},$$

as desired. $\square$

Observation 1 *The LYM inequality implies Sperner's Theorem. Indeed, we have*

$$1 \geq \sum_{k=0}^n \frac{|\mathcal{F}_k|}{\binom{n}{k}} \geq \sum_{k=0}^n \frac{|\mathcal{F}_k|}{\binom{n}{\lfloor n/2 \rfloor}} = \frac{|\mathcal{F}|}{\binom{n}{\lfloor n/2 \rfloor}}.$$

Furthermore, equality in Sperner implies equality in LYM and in the second inequality above. Hence $\mathcal{F} = \binom{[n]}{\lfloor n/2 \rfloor}$ and $\binom{[n]}{\lceil n/2 \rceil}$ are the only equality cases for Sperner's theorem.

Let us now discuss the shadows in more detail. In particular, for a k -uniform set system $\mathcal{F}$, how to minimize $|\partial\mathcal{F}|$, given $|\mathcal{F}|$. We may apply Local LYM to get

$$\frac{|\partial\mathcal{F}|}{\binom{n}{k-1}} \geq \frac{|\mathcal{F}|}{\binom{n}{k}}.$$

This bound is very good when $|\mathcal{F}|$ is close to $\binom{n}{k}$, but becomes really bad for small (e.g. constant) $|\mathcal{F}|$. So, what to do instead? For $|\mathcal{F}| = \binom{t}{k}$ for some $t \leq n$ it seems plausible that the family $\binom{[t]}{k}$ minimizes the shadow (e.g. this easy to verify for $k = 2$). If we believe this, then for general $|\mathcal{F}|$ we need to find a good way to interpolate between $\binom{[t]}{k}$ and $\binom{[t+1]}{k}$. This is achieved by the colexicographic ordering.

Definition 2 The colexicographic (colex, for short) order on $\binom{[n]}{k}$ is the total order in which

$$A < B \text{ if } A \neq B \text{ and } \max(A \triangle B) \in B.$$

This can also be expressed as

$$A < B \text{ if } \sum_{i \in A} 2^i < \sum_{i \in B} 2^i.$$

Example 1 Let $k = 3$. The first 10 elements in the colex order in uniformity k are (omitting parentheses):

$$123, 124, 134, 234, 125, 135, 235, 145, 245, 345.$$

Remark 1 Note that for $m = \binom{n}{k}$ the initial segment of length m of colex on $\binom{[n]}{k}$ is precisely $\binom{[n]}{k}$. Thus, colex is also an ordering on $\binom{[n]}{k}$, which is consistent for a fixed k between different n .

Remark 2 Colex is similar to the better known lexicographic order (lex, for short), in which

$$A < B \text{ if } A \neq B \text{ and } \min(A \triangle B) \in A.$$

This is just the familiar dictionary order. Note though, that for $k \geq 2$ lex, unlike colex, does not form a well-ordering of $\binom{[n]}{k}$ of ordinal type ω . Also, lex on $\binom{[n]}{k}$ is not consistent for a fixed k between different n .

With this in hand, we can state the Kruskal-Katona Theorem on minimizing the shadow.

Theorem 3 (Kruskal-Katona) Let $\mathcal{F} \subseteq \binom{[n]}{k}$, and let $\mathcal{A}$ be the family consisting of the first $|\mathcal{F}|$ elements of $\binom{[n]}{k}$ in colex order. Then

$$|\partial\mathcal{F}| \geq |\partial\mathcal{A}|.$$

A major technique in the proof will be set shifting.

Definition 3 For a set $A \subseteq [n]$ and distinct $i, j \in [n]$ we define the shifting operator $C_{j \rightarrow i}(A)$ as

$$C_{j \rightarrow i}(A) = \begin{cases} (A \setminus \{j\}) \cup \{i\} & \text{if } j \in A \text{ and } i \notin A, \\ A & \text{otherwise.} \end{cases}$$

For a family $\mathcal{F}$, we define

$$C_{j \rightarrow i}(\mathcal{F}) = \{C_{j \rightarrow i}(A) \mid A \in \mathcal{F}\} \cup \{A \in \mathcal{F} \mid C_{j \rightarrow i}(A) \in \mathcal{F}\}.$$

For $i < j$, we call $C_{j \rightarrow i}$ a left-shift, or a left-compression.

Example 2 Let $n = 5, k = 3$ and (omitting parentheses) $\mathcal{F} = \{123, 134, 234, 235\}$. Let $i = 1$ and $j = 2$. Then

$$\begin{aligned} C_{j \rightarrow i}(123) &= 123, \\ C_{j \rightarrow i}(134) &= 134, \\ C_{j \rightarrow i}(234) &= 134 \in \mathcal{F}, \text{ and} \\ C_{j \rightarrow i}(235) &= 135 \notin \mathcal{F}. \end{aligned}$$

Therefore,

$$C_{j \rightarrow i}(\mathcal{F}) = \{123, 134, 234, 135\}.$$

Observation 2 The following hold for any $A \subseteq [n]$ and $\mathcal{F} \subseteq 2^{[n]}$:

1. $|C_{j \rightarrow i}(A)| = |A|$,
2. $|C_{j \rightarrow i}(\mathcal{F})| = |\mathcal{F}|$,
3. $C_{j \rightarrow i}(C_{j \rightarrow i}(\mathcal{F})) = C_{j \rightarrow i}(\mathcal{F})$.