1. Show that VERTEX COVER is FPT with parameter treewidth. Meaning that given  $k \in \mathbb{N}$ , we can in time  $f(t)n^{\mathcal{O}(1)}$  determine whether an input graph G with t = tw(G) has a vertex cover of size at most k for some computable function f.

Can you also do it in time  $c^t n^{\mathcal{O}(1)}$  for some constant c?

- In the PLANAR VERTEX COVER problem, we are given a planar graph G and an integer k. The goal is to determine whether it has a vertex cover of size at most k.
  Can you do it in time 2<sup>O(√k)</sup>n<sup>O(1)</sup> for some constant?
- Show that HAMILTONIAN CYCLE is FPT with parameter treewidth.
  Hint. In the dynamic program, you will need more than just the information from the subtree.
- 4. Show that STEINER TREE is FPT with parameter treewidth.