Annotation

- A-sort,
- Sparse polynomials and matrices,
- Low-level Access to Memory,
- Hashing,
- Heaps,
- Arithmetic expressions, notations and conversion between them,
- Graphs and their representation,
- Graph-algorithms.

A-sort

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- Finger points at the leaf (of the B-tree) where we inserted last time.
- We are not inserting from the root but from the "finger".
- Good results when the input is pre-sorted (we don't bubble too often to the root).

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- Head can be used to find that we reached the end of the polynomial.

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- Linear list of the elements (ordered in both dimensions),
- linear list of linear lists (list of rows consisting of list of columns),
- Dividing into quarters (divide the matrix into four parts – left top, right top, left bottom, right bottom). If the submatrix is "too large"and non-zero, we divide again.

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- `function MemAvail: longint;` – returns number of available bytes on heap
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- `function MaxAvail: longint;` – returns size of the largest free block (largest allocable size)
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- `FreeMem` – deallocates memory allocated by `GetMem` – (DTTO)

Example of GetMem/FreeMem

We create an array of uncertain length

```
type parr=^tarr;
     tarr=array[1..10000] of longint;
var arr:parr;
begin
      GetMem(arr,500);{get 500 bytes}
      arr^[10]:=1000;{This is OK}
      arr^[500]:=1024;{Problem -- array too small!}
      FreeMem(arr,500);{FreeMem(arr); should suffice}
end.
```

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- This is called the *hashing*.
- It may happen that more candidates want to seat the same cell in the table. This is called a *collision*.

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- Where to place the element in collision? There are many possibilities. Either we pick next free cell or we design a function that proposes next cell.
- If we know the size of the data, we may try to implement *perfect hashing*, i.e., hashing without collisions. Hashing should be in more detail explained in the lecture of Algorithms and Data Structures (proofs).

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- $((10 \ 5 \ +) \ (15 \ 4 \ -) \ *) \ 2 \ /$ (with superfluous parentheses),
- by a tree: Each node contains an operator (and has two sons - operands) or a value (leaf).
- Evaluation will be only sketched, pseudocode has to be creatively interpreted!

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- Is it possible to evaluate expressions in all these notations?
- Can we convert one notation into another one?
- Yes, e.g., using the tree.

Evaluating the prefix notation

We use recursion: `function evaluate:integer;`

```
begin
      if (we read a number) then
            evaluate:=value_of_the_input_number
      else
      begin operator:=read_operator();
            arg1:=evaluate;
            arg2:=evaluate;
            evaluate:=perform(operator,arg1,arg2);
      end;
end;
```

Tree from the prefix notation

```
function pref_tree:tree;
begin
      if(we read a number) then
            pref_tree:=leaf(value_of_input);
      else
      begin tmp:=inner_node(operator);
            tmp.arg1:=evaluate;
            tmp.arg2:=evaluate;
            evaluate:=tmp;
      end;
end;
```

function `leaf` creates a leaf,
function `inner_node` creates a node of out_deg 2,
vertex of out_deg 2 has sons `arg1` a `arg2`.

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- in one phase we search right sone and in one phase we write the operator.
- All three notations arise by correct ordering of these phases.
- Even always we visit left son before the right one, thus the only change is the time of outputting the operator!

Generating prefix notation

```
procedure gen_pref(v:tree);
begin
      if(leaf(v)) then
            output(v);
      else
      begin output(v);
            gen_pref(v.arg1);
            gen_pref(v.arg2);
      end;
end;
```

Function `output` outputs the operator or number (resp.), function `leaf` decides whether a given node is a leaf.

Generating postfix notation

```
procedure gen_pref(v:tree);
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Generating infix notation

almost correctly!

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Generating infix notation

ugly but correctly!

```
procedure gen_pref(v:tree);
begin
      if(leaf(v)) then
            output(v);
      else
      begin write('(');
            gen_pref(v.arg1);
            output(v);
            gen_pref(v.arg2);
            write(')');
      end;
end;
```

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Evaluating postfix notation

...towards the solution

Revision of our knowledge:

Buffer is a data structure equipped by operations:

- `push` – insert onto the buffer-top,
- `pop` – remove from the buffer-top,
- i.e., last in, first out.

Evaluating postfix notation

```
function eval_post:integer;
begin
     while not eof do
     begin if (we read a number) then
                    push(number);
          if (we read an operator) then
          begin arg2:=pop;
                    arg1:=pop;
                    push(operator(arg1,arg2));
          end;
     end;
     writeln(pop);{Result is on the buffer-top}
end;
```

Tree from the prefix notation

```
function tree_post:tree;
begin
      while not eof do
      begin if (we read a number) then
                  push(leaf(number));
            if (we read an operator) then
            begin pom:=node(operator);
                  pom.arg2:=pop;
                  pom.arg1:=pop;
                  push(pom);
            end;
      end;
      tree_post:=pop;{Result is on the buffer-top}
end;
```

Evaluating the tree

should be clear, but let's go:

```
function eval_tree(v:tree);
begin
      if(leaf(v)) then
            eval_tree:=value(v)
      else
      begin arg1:=eval_tree(v.arg1);
            arg2:=eval_tree(v.arg2);
            op:=operator(v);
            eval_tree:=op(arg1,arg2);
      end;
end;
```

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5. in cases 2^i (i.e., 1, 2, 4) employ recursion.

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- How to represent a graph while programming?

Graphs – remarks to definition

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- Define basic notions (walk, trail, path, connectivity, trees).

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- Advantages and disadvantages?
- Can we convert these representations?

Converting A_G to B_G and back

```
init_with_0s(B_G);
edge_index:=1;
for i:=1 to n do begin
      for j:=i+1 to n do begin
      if(A_G[i,j]=1) then
      begin
            B_G[i,edge_index]:=1;
            B_G[j,edge_index]:=1;
            inc(edge_index);
      end;
end;
```

B_G to A_G

Either we analyze the incidence matrix (in a similar way) or:

$A_G := B_G \times {B_G}^T$;

`for i:=1 to n do`

$\qquad A_G[i,i] := 0$;

Důkaz.

Exercise in Combinatorics and Graph Theory I $\square$

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- `vertices`,
- `edges` or `edge(u,v)` – we can find it through `vertices` and `find_neighbors`,

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Further graph representations

- List of vertices and edges incident to individual vertices.
- I.e., we are keeping a list of edges incident to each vertex in the graph.
- We employ linear- (linked-)lists. If we employ an array, what do we get?
- The adjacency matrix!

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- Advantages/disadvantages?
- In oriented case we have to modify the representation.

Walk, trail, path, circle

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- A trail is a circle if it starts and ends in the same vertex and each vertex occurs there exactly once.

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- How do we decide whether a graph is a tree?
- Similarly!

Graph connectivity

Graph is connected iff from one (fixed) vertex we can reach all the other vertices.

```
for i in vertices do
      unvisit(i); {so far we visited nothing}
i:=start_vertex;
queue:={i};{for reachable vertices}
while nonempty(queue) do begin
      visit(i);
      queue:=queue+unvisited_neighbors(i);
end;
connected:=true;
for i in vertices do begin
      if unvisited(i) then
            connected:=false;
```

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- for-cycle takes place at most *n*-times.

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- Thus a good representation yields the complexity $\Theta(m+n)$.

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- DFS may get implemented using recursion (and thus without an auxiliary data-structure).
- BFS visits the vertex using the shortest path.

Looking for a cycle

A graph has a cycle if we return to a particular vertex while searching the graph.

```
cycle:=false; {so far no cycle}
for i in vertices do unvisit(i);
for i in vertices do
    if unvisited(i) then{new component}
    begin queue:={i};
        while(nonempty(queue)) do
        begin dequeue_from_queue_and_assing_into(i);
if(visited(i)) then
                cycle:=true;
            else  for j in neighbors(i) do
                begin queue:=queue+{j};
                    erase_edge({i,j});
                end;
    end; end;
```

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- Or we test cycle-freeness and connectivity (one component).
- Or we test connectivity (or cycle-freeness) and an appropriate number of edges (Euler's formular).

Shortest path

When looking for the shortest path, it depends on the representation:

- Perform BFS (considering the list of vertices and edges),

Theorem

In ${A_G}^k$ position i, j gives number of walks with length k from (vertex) i to j.

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Dijkstra's algorithm

Looks for the shortest path from a given vertex into all other vertices

Input: Graph with `nonnegatively` evaluated edges.

- We keep the "queue" for vertices ordered by the shortest so far found path.
- At the beginning we inicialize the distances to all vertices [except start] by infinity [large-enough value], distance to start is 0.
- We add `start` into the queue for reachable vertices.
- Remove the first vertex of the "queue" and inspect its neighbors.
- Repeat this while the "queue" is non-empty.

Extending the path

When extending the path, for a vertex v in distance $d(v)$ we try for each edge $\{v, w\}$ whether

$$d(w) > d(v) + length(\{v, w\}).$$

If so, let $d(w) := d(v) + length(\{v, w\})$ and correct the position of w in the "queue".

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- The algorithm is kind of modification of BFS!
- Complexity depends on the representation of the graph and of the 'queue'!

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Thank you for your attention...