## **Combinatorics**

Exercise 6 – List colouring

## Problems

- 1. Let G be a d-degenerate graph. Prove that  $ch(G) \leq d+1$ .
- 2. Let G be a graph such that every subgraph  $H \subseteq G$  satisfies  $|E(H)| \leq k|V(H)|$ . Prove that G has an orientation with maximum outdegree  $\leq k$
- 3. For every k, find a bipartite graph  $G_k$  with  $ch(G_k) \ge k$ . Conclude that one cannot bound ch(G) from above by any function of  $\chi(G)$ .
- 4. Let G be a bipartite graph with n vertices. Prove that  $ch(G) \leq \lfloor \log_2(n) \rfloor + 1$ .
- 5. Prove that every orientation of a bipartite graph has a kernel.
- 6. Let  $G_n$  be the star with n+1 vertices (i.e. the graph  $K_{1,n}$ ).
  - (a) Observe that  $ch(G_n) = 2$ .
  - (b) Let G be a graph, let k be an integer and consider the following "game" played on G by two players, *Painter* and *Paint shop clerk*. There are k rounds. In each of the k rounds, Painter points at a vertex v of her choice and Paint shop clerk offers her a new colour which has not yet been offered for v. That is, after k rounds, Painter has a list of some colours for every vertex and the sum of the sizes of these lists is k. After this, Painter wins if G has a list colouring using these lists. If k is the smallest number for which Painter has a winning strategy on G in k rounds, we write  $ich(G) = \frac{k}{|V(G)|}$ .
    - i. Observe that  $ich(G) \leq ch(G)$ .
    - ii. Prove that  $ich(G_n) \leq 1 + \frac{2\sqrt{n}}{n+1}$  (in other words, one needs just a little bit more than 1 colour per vertex on average).