

(L8) (April 24, 2020) | A Power sum is a formal expression $S(x) = \sum_{i=1}^r p_i(x) d_i^x$, where x is a formal variable, $r \in \mathbb{N}_0$, $p_i \in \mathbb{C}[x]$ are $\neq 0$ polynomials and $d_i \in \mathbb{C}$ are $\neq 0$ and mutually distinct numbers. For $r=0$ we define

$S(x) := 0$. For example, $1^x + (-1)^x$ is a power sum, and so is Binet's formula

$$\frac{1}{\sqrt{5}} \left(\frac{1+\sqrt{5}}{2} \right)^x - \frac{1}{\sqrt{5}} \left(\frac{1-\sqrt{5}}{2} \right)^x$$

from the last lecture.

Every polynomial $p \in \mathbb{C}[x]$, $p \neq 0$, is a power sum $p(x) \cdot 1^x$.

Substituting $x := n \in \mathbb{N}$ we make $S(x)$ a function or sequence $S(n)$, $S: \mathbb{N} \rightarrow \mathbb{C}$.

Theorem (another \iff) for

If $(a_n) = (a_1, a_2, \dots) \subset \mathbb{C}$ is a LRS [LRS]

then \exists a power sum $S(x)$ s.t. $\forall n \in \mathbb{N}: S(n) = a_n$.

If $S(x)$ is a power sum then $S: \mathbb{N} \rightarrow \mathbb{C}$ is a LRS. [Proof. \Rightarrow] We assume that

$a_{n+k} = \sum_{i=0}^{k-1} c_i a_{n+i}$ for $\forall n \in \mathbb{N}$ and some $c_0, c_1, \dots, c_{k-1} \in \mathbb{C}$, $c_0 \neq 0$. We
 strands $\{c_i\}_{i=0}^{k-1}$ and $c_0, c_1, \dots, c_{k-1} \in \mathbb{C}$, $c_0 \neq 0$. We
 first extend our sequence by one more term a_k to
~~the~~ sequence $a: \mathbb{N}_0 \rightarrow \mathbb{C}$, ~~so that it is better~~ so that
 we can work better with the GF of (a_n) . In
 fact, we extend the sequence to the sequence
~~by~~
~~a: \mathbb{Z} \rightarrow \mathbb{C}, following the backward direction~~

$$a_n = -\frac{c_1}{c_0} a_{n+1} - \dots - \frac{c_{k-1}}{c_0} a_{n+k-1} + \frac{a_{n+k}}{c_0}.$$

So we have ^a sequence (a_0, a_1, a_2, \dots) and take its
 GF $A(x) = \sum_{n=0}^{\infty} a_n x^n$. We know from the
 previous theorem

$$\begin{aligned}
 A(x) &= \frac{P(x)}{1 - x c_{k-1} - x^2 c_{k-2} - \dots - x^{k-1} c_0} \\
 &=: q(x)
 \end{aligned}$$

where $P, q \in \mathbb{C}[x]$ and, as is easy to see,
 $P \neq 0$ or $\deg(P) < k$. We factorize the denominator
 of $q(x) = 1 - c_{k-1}x - c_{k-2}x^2 - \dots - c_0x^k$ as

$$q(x) = \prod_{i=1}^j (1-d_i x)^{q_i} \text{ where } d_i \in \mathbb{C} \setminus \{0\} \quad (3)$$

mutually distinct and $q_i \in \mathbb{N}$, with $q_1 + q_2 + \dots + q_j = q$. Here the d_i are roots of the char. polynomial of the recurrence. So

$$A(x) = \sum_{n=0}^{\infty} a_n x^n = \frac{P(x)}{\prod_{i=1}^j (1-d_i x)^{q_i}} \quad \begin{array}{l} P \in \mathbb{C}[x] \\ \text{with } P \neq 0 \\ \text{or } \deg(P) < q \end{array}$$

We expand in partial fractions:

$$= \sum_{i=1}^j \sum_{l=1}^{q_i} \frac{B_{il}}{(1-d_i x)^l} \quad \text{where } B_{il} \in \mathbb{C}.$$

We have the generalized geometric series identity

$$\frac{1}{(1-dx)^l} = \sum_{n=0}^{\infty} \binom{n+l-1}{n} d^n x^n \quad [\text{it follows by } dx]$$

[differentiating ~~of~~ $\frac{1}{1-dx} = 1 + dx + d^2 x^2 + \dots$]. We substitute it in (3), compare coeffs of x^n on both sides and get for a_n the expression

*) Recall the precise statement and proof of the proposition on Partial fractions.

$$a_n = \sum_{i=1}^j \sum_{l=1}^{d_i} \binom{n+l-1}{l} d_i^n, \text{ which is a power, } \quad (4)$$

Sum expression $a_n = S(n)$ because $\binom{n+l-1}{l} = \binom{n+l-1}{l-1} = \frac{(n+l-1)(n+l-2)\cdots(n+1)}{(l-1)!}$ is a polynomial in n of degree $l-1$.

Let $S(x) = \sum_{i=1}^k p_i(x) d_i$ be a power sum. We want to show that $(S(1), S(2), \dots)$ is a LRS. We have to make $(S(n))$ a LRS similar to the following well known (?) Lemma from linear algebra.

Lemma (L) Let K be a field, $m, n \in \mathbb{N}$ with $m < n$, $a_{i,j} \in K$ for $i=1, 2, \dots, m$, $j=1, 2, \dots, n$ be elements of the field. Then $\exists d_1, d_2, \dots, d_n \in K$, not all equal to 0_K , such that $\forall i=1, 2, \dots, m: \sum_{j=1}^n a_{i,j} d_j = 0_K$. In other words: every homogeneous linear system of equations with more unknowns than equations has a nontrivial solution. Proof. Exercise!

Let $d := \max_{1 \leq i \leq r} \deg(p_i)$ and $\alpha := r(1+d)$. (5)

Since α is the shifted p-sum $S(t+\alpha) =$

$$= \sum_{i=1}^r d_i^{t+\alpha} p_i(t+\alpha) d_i^{-t} = \sum_{i=1}^r p_i(t+\alpha) d_i^{t+\alpha}$$

is a bin. combining func.

[expanded by the binomial formula] (with coeffs. in \mathbb{C}) of the monomials $x^{\alpha} d_i^t$ for

$0 \leq \alpha \leq d$, $1 \leq i \leq r$ — there are $r(1+d)$ of these monomials. Hence

$S(t) =$ lin comb. L_0 of these monomials

$$S(t+1) = -1/ - L_1 \quad \dots \quad 1/ \dots$$

$$\vdots \quad \vdots \quad \vdots$$

$$(S(t+\alpha)) = -1/ - L_\alpha \quad \dots \quad 1/ \dots$$

By the lemma L there exist coeffs. $b_0, b_1, \dots, b_\alpha \in \mathbb{C}$, not all 0, s.t. $\sum_{i=0}^\alpha b_i S(t+i) = 0$.

Let $b_{\alpha'}^*$ be the first

$$\boxed{\sum_{i=0}^\alpha b_i S(t+i) = 0}$$

last ≠ 0 coeff. b_i ; so $0 \leq \alpha' \leq \alpha'' \leq \alpha$. We define

$$v := \alpha'' - \alpha', c_0 := -\frac{b_{\alpha'}}{b_{\alpha''}}, c_1 := -\frac{b_{\alpha'+1}}{b_{\alpha''}}, \dots,$$

$c_{v-i} := -\frac{b_{q''}}{b_q''}$. It follows from $\boxed{\dots}$ (6)

that further: $S(u+v) = \sum_{i=0}^{v-1} c_i S(u+i)$.
 Thus $(S(u))$ is a LRS. $(c_0 = -b_{q''}/b_q'' \neq 0)$ \square

Lemma L can be used to prove the next interesting result on LRS (which we prove yet time).

Proposition

Let KCL be an exte sign of fields and let

(a_n) $\subset K$ satisfy $a_{n+q} = \sum_{i=0}^{q-1} c_i a_{n+i}$ for such and some constants $q \in \mathbb{N}_0$, $c_0, c_1, \dots, c_{q-1} \in L$.

Then there exist $q' \in \mathbb{N}_0$ and $c'_0, c'_1, \dots, c'_{q'-1} \in K$ s.t. then: $a_{n+q'} = \sum_{i=0}^{q'-1} c'_i a_{n+i}$.

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So, for

example, if (a_n) $\subset \mathbb{Q}$ is a LRS of fractions then we can assume that (a_n) is determined by a recurrence relation whose coeffs are fractions as well. As I wrote, we prove it in the next lecture.