## **Lecture 2.** $C_n = \frac{1}{n} \binom{2n-2}{n-1}$ again—the right GF proof. D-finite and algebraic FPS

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Last time we justified the equality

(the 
$$f \in \mathbb{R}[[x]]$$
 s. t.  $f(x)^2 = 1 - 4x$  and  $f(0) = 1 = \sum_{n=0}^{\infty} {\binom{1/2}{n}} (-4)^n x^n$ 

which was needed to get the formula  $C_n = \frac{1}{n} \binom{2n-2}{n-1}$  by solving the quadratic equation  $C^2 - C + x = 0$ . Literature on enumerative combinatorics glosses over this justification. Now we give the "right" simple derivation of the formula by generating functions. The proper way to make use of the equation  $C^2 - C + x = 0$  is not to apply the quadratic formula (!) but to differentiate it. Recall that  $C = C(x) = \sum_{n\geq 1} C_n x^n$  is the generating function of Catalan numbers.

The third proof of Theorem 1  $(C_n = \frac{1}{n} \binom{2n-2}{n-1})$  in Lecture 1. From  $C^2 - C + x = 0$  we have that  $-C^2 + C = x$ . Differentiating  $C^2 - C + x = 0$  by x we thus get

$$2CC' - C' + 1 = 0 \rightsquigarrow C' = \frac{1}{1 - 2C} = \frac{\frac{1}{2}C - \frac{1}{4}}{-C^2 + C - \frac{1}{4}} = \frac{2C - 1}{4x - 1}$$

and (4x-1)C'-2C+1=0. Since  $C = \sum_{n\geq 1} C_n x^n$  and  $C' = \sum_{n\geq 1} nC_n x^{n-1}$ , equating for  $n \in \mathbb{N}$  the coefficient of  $x^n$  on the left side to zero gives the equation

$$4nC_n - (n+1)C_{n+1} - 2C_n = 0.$$

Thus  $C_1 = 1$  and  $C_{n+1} = \frac{4n-2}{n+1} \cdot C_n$  for any  $n \in \mathbb{N}$ . Hence for  $n \ge 2$ ,

$$C_n = \prod_{k=2}^n \frac{2(2k-3)}{k} = \frac{2^{n-1} \cdot (2n-3)!!}{n!}$$
$$= \frac{2^{n-1}(n-1)! \cdot (2n-3)!!}{(n-1)! \cdot n!} = \frac{1}{n} \binom{2n-2}{n-1}.$$

We have a new simple recurrence for  $C_n$ :  $C_n = 1$  and for  $n \in \mathbb{N}$ ,

$$C_{n+1} = \frac{4n-2}{n+1} \cdot C_n$$

As  $0 < \frac{4n-2}{n+1} < 4$  for  $n \in \mathbb{N}$ , induction gives the exponential upper bound  $C_n < 4^n$ . The exact asymptotics of the Catalan numbers follows from the Stirling formula  $n! = (1 + o(1)) \cdot \sqrt{2\pi n} \cdot (n/e)^n$ . We get that for some constant c > 0,

$$C_n = (c + o(1)) \cdot n^{-3/2} \cdot 4^n \quad (n \to \infty) .$$

In the previous proof we obtained the differential equation

$$(4x-1)C' - 2C + 1 = 0.$$

It means that  $C \in \mathbb{R}[[x]]$  is a D-finite FPS. We define this class of FPS.

**Definition 1** Let  $K = \mathbb{R}$  or  $K = \mathbb{C}$ . We say that a FPS  $f(x) = \sum_{n\geq 0} a_n x^n$  in K[[x]] is D-finite (or holonomic) if there exist polynomials  $p_{-1}, p_0, \ldots, p_k$  in  $K[x], k \in \mathbb{N}_0$ , such that  $p_k \neq 0$  and

$$\sum_{i=-1}^{k} p_i(x) f^{(i)}(x) = 0 \; .$$

Here  $f^{(-1)}(x) := 1$ ,  $f^{(0)}(x) := f(x)$  and  $f^{(i)}(x)$  for  $i \in \mathbb{N}$  is the *i*-th derivative of f,

$$f^{(i)}(x) = \sum_{n=0}^{\infty} a_n n(n-1) \dots (n-i+1) x^{n-i}$$

In words, f(x) satisfies a (non-homogeneous) linear differential equation with polynomial coefficients. As usual, for i = 1, 2, 3 we write for  $f^{(i)}$  synonymously f', f'' and f'''. By repeated differentiation we homogenize the equation and get that  $p_{-1} = 0$ .

C = C(x) satisfies also the equation

$$C^2 - C + x = 0$$

and so is an example of an *algebraic* FPS.

**Definition 2** Let  $K = \mathbb{R}$  or  $K = \mathbb{C}$ . We say that a FPS  $f(x) = \sum_{n\geq 0} a_n x^n$ in K[[x]] is algebraic if there exist polynomials  $p_0, \ldots, p_k$  in  $K[x], k \in \mathbb{N}$ , such that  $p_k \neq 0$  and

$$\sum_{i=0}^{k} p_i(x) f^i(x) = 0 \; .$$

In other words, P(x, f(x)) = 0 for a nonzero polynomial  $P = P(x, y) \in K[x, y]$ .

In the previous proof we deduced D-finiteness of C from its algebraicity. We show that this transition works in general.

**Proposition 3** Let  $K = \mathbb{R}$  or  $K = \mathbb{C}$ . If a FPS  $f \in K[[x]]$  is algebraic then f is D-finite.

*Proof.* Let  $f \in K[[x]]$  be such that  $\sum_{i=0}^{k} p_i f^i = 0$  where  $k \in \mathbb{N}$ ,  $p_i \in K[x]$  and  $p_k \neq 0$ . Thus

$$f^k = \sum_{i=0}^{k-1} q_i f^i \tag{1}$$

where  $q_i \in K(x)$  are rational functions. Differentiating by x we get that

$$kf^{k-1}f' = \sum_{i=0}^{k-1} (q'_i f^i + q_i i f^{i-1}f')$$
 and  $f'\left(kf^{k-1} - \sum_{i=0}^{k-1} q_i i f^{i-1}\right) = \sum_{i=0}^{k-1} q'_i f^i$ .

Thus

$$f' = \frac{P_1(f)}{Q_1(f)} \in K((x))$$
(2)

(the field of formal Laurent series in x with coefficients in K) where  $P_1(y)$ and  $Q_1(y) \neq 0$  are polynomials in K(x)[y].  $Q_1(y) \neq 0$  because it has leading coefficient k. We show by induction on  $j \in \mathbb{N}$  that there exist polynomials  $P_j(y)$ and  $Q_j(y) \neq 0$  in K(x)[y] such that

$$f^{(j)} = \frac{P_j(f)}{Q_j(f)}$$

For j = 1 we proved it above. We differentiate this equation by x and get that

$$f^{(j+1)} = \frac{(P_j(f))_x \cdot Q_j(f) - P_j(f) \cdot (Q_j(f))_x}{Q_j(f)^2}$$

We have that  $(P_j(f))_x = R_j(f) + S_j(f)f'$  and  $(Q_j(f))_x = T_j(f) + U_j(f)f'$ where  $R_j$ ,  $S_j$ ,  $T_j$  and  $U_j$  are in K(x)[y]. Replacing f' by equation (2) we get the required expression

$$f^{(j+1)} = \frac{P_{j+1}(f)}{Q_{j+1}(f)}$$
 with  $Q_{j+1} = Q_1 \cdot Q_j^2$ 

We bring the k + 1 fractions  $f^{(0)} = \frac{f}{1}$ ,  $f^{(1)} = \frac{P_1(f)}{Q_1(f)}$ , ...,  $f^{(k)} = \frac{P_k(f)}{Q_k(f)}$  to a common denominator, reduce powers of f in the numerators by equation (1) and get the expressions

$$f^{(0)} = \frac{A_0(f)}{B(f)}, f^{(1)} = \frac{A_1(f)}{B(f)}, \dots, f^{(k)} = \frac{A_k(f)}{B(f)}$$

where  $B(y) \neq 0$  and  $A_i(y)$  are in K(x)[y] and every  $A_i(y)$  has in y degree at most k-1. It follows that the k+1 numerators can be non-trivially linearly combined to 0 by some coefficients in K(x). Hence

$$\sum_{i=0}^{k} c_i f^{(i)} = 0$$

for some  $c_i \in K(x)$ , not all of them 0. Thus f is D-finite.

In the initial proof we derived for  $(C_n) = (C_1, C_2, ...)$  the recurrence that for every  $n \in \mathbb{N}$ ,

$$(n+1)C_{n+1} + (2-4n)C_n = 0$$

Recurrences of this form are characteristic for sequences of coefficients of D-finite formal power series.

**Definition 4** Let  $K = \mathbb{R}$  or  $K = \mathbb{C}$ . We call a sequence  $(a_n) = (a_0, a_1, ...) \subset K$  P-recurrent (over K) if there exist polynomials  $p_i \in K[x]$ , where i = 0, 1, ..., k with  $k \in \mathbb{N}_0$  and  $p_k \neq 0$ , such that for every  $n \in \mathbb{N}_0$ ,

$$\sum_{i=0}^{k} p_i(n) \cdot a_{n+i} = 0$$

It is easy to see that this is equivalent with the modified definition when the last displayed equality holds only for every  $n > n_0$ . To get  $(C_n)$  P-recurrent exactly according to Definition 4 we extend it with  $C_0 := -\frac{1}{2}$ .

**Proposition 5** Let  $K = \mathbb{R}$  or  $K = \mathbb{C}$  and  $f = f(x) = \sum_{n=0}^{\infty} a_n x^n \in K[[x]]$ . Then f is D-finite  $\iff (a_n)$  is P-recurrent.

*Proof.* Suppose that f is D-finite. For  $i \in \mathbb{N}_0$  we introduce notation  $(x)_0 = 0$  and  $(x)_i = x(x-1)\dots(x-i+1)$  for i > 0. Thus f satisfies the differential equation

$$\sum_{i=0}^{k} p_i(x) f^{(i)}(x) = \sum_{(h,i) \in P} b_{h,i} x^h f^{(i)}(x) = 0$$
(3)

where  $k \in \mathbb{N}_0$ ,  $p_i \in K[x]$ ,  $p_k \neq 0$ ,  $P \subset \mathbb{N}_0^2$  is a nonempty set and every coefficient  $b_{h,i} \in K$  is nonzero. Let  $H \in \mathbb{N}_0$  be the maximum value of h in the pairs  $(h, i) \in P$ . Since  $(i \in \mathbb{N}_0)$ 

$$f^{(i)}(x) = \sum_{n=0}^{\infty} (n+i)_i a_{n+i} x^n ,$$

by setting the coefficient of  $x^n$ ,  $n \in \mathbb{N}_0$ , in equation (3) to zero we get for every  $n \ge H$  that

$$\sum_{(h,i)\in P} b_{h,i} \cdot (n-h+i)_i \cdot a_{n-h+i} = 0 .$$

Grouping together summands sharing the same coefficient  $a_{n-h+i}$ , we deduce that the sequence  $(a_n)$  satisfies for  $n > n_0$  a (nontrivial) P-recurrence as in Definition 4.

Suppose that the sequence  $(a_n)$  of coefficients of f(x) is P-recurrent. Thus for every  $n \in \mathbb{N}_0$  we have the equality

$$\sum_{i=0}^{k} p_i(n) \cdot a_{n+i} = 0$$

for some  $p_i \in K[x]$  with  $p_k \neq 0$ . We switch from the basis  $\{x^i \mid i \in \mathbb{N}_0\}$  of the *K*-vector space K[x] to the basis  $\{(x+i)_i \mid i \in \mathbb{N}_0\}$  and express every  $p_i(x)$  as the linear combination

$$p_i(x) = \sum_{j=0}^{d_i} c_{i,j} \cdot (x+j)_j$$

where  $d_i = \deg p_i \in \mathbb{N}_0$ ,  $c_{i,j} \in K$  and  $c_{k,d_k} \neq 0_K$  (if  $p_i = 0$  then this sum is empty). Thus for every  $n \in \mathbb{N}_0$ ,

$$\sum_{i=0}^{k} \left( \sum_{j=0}^{d_i} c_{i,j} \cdot (n+j)_j \right) a_{n+i} = 0 \; .$$

By grouping the terms according to the pairs (i, j) we get that for every  $n \in \mathbb{N}_0$ the equality

$$\sum_{(i,j)\in Q} c_{i,j} \cdot (n+j)_j \cdot a_{n+i} = 0$$

holds, where  $Q \subset \mathbb{N}_0^2$  is nonempty and (as we know) not all  $c_{i,j} \in K$  are zero. It follows that

$$\sum_{(i,j)\in Q} c_{i,j} \cdot \left(x^{j-i} \cdot f(x)\right)^{(j)} = 0$$

because for every  $n \in \mathbb{N}_0$  the coefficient of  $x^n$  in the FPS on the left side is 0. It is clear that this relation can be converted to a linear differential equation with polynomial coefficients, not all of them zero, for f(x). We see that f(x) is D-finite.

D-finite formal power series were introduced by R. P. Stanley in [1]. For more information on uses of algebraic and D-finite generating functions in enumerative combinatorics see his book [2].

## References

- R. P. Stanly, Differentiably finite power series, European J. Combinatorics 1 (1980), 175–188
- [2] R. P. Stanly, *Enumerative Combinatorics. Volume 2*, Cambridge University Press, Cambridge, UK 1999