## Exercise Sheet, Week 8

**Question 1.** Use *merge sort* to sort the following sequences:

 $\langle 2, 7, 2, 3, 8, 5, 4, 1 \rangle$   $\langle 1, 2, 3, 4, 5, 6, 7 \rangle$ 

**Question 2.** Use *quicksort* to sort the first sequence of Question 1 with the following pivot-selection strategies:

- (a) The leftmost element.
- (b) The middle element.
- (c) A random element given by the following "random" sequence of positions:

 $6, 0, 2, 0, 1, 4, 1, 3, 1, 2, 4, 1, 0, 2, 3, 5, 1, 2, 0, 0, 0, 0, 0, 0, \dots$ 

(Use mod if the position is out of scope of the part of the array which you are sorting.)

**Question 3.** What would be the time complexity of the modified *merge sort* which calls itself recursively only as long as the size is n or  $\frac{n}{2}$  and, for smaller inputs, it calls *selection sort* instead?

**Question 4.** Instead of recursion, use stacks to implement *quicksort*. You can use function **partition** as in the lecture. *Hint*: It might help to use the following class to remember which partitions are to be sorted:

```
class Interval {
1
       int left;
2
       int right;
3
4 }
void quicksort(int[] arr) {
       s = new Stack < Interval > ();
       s.push(new Interval(0, arr.length -1));
3
4
       while (s.notEmpty()) {
6
7
9
10
11
12
13
       }
14
15 }
```

Question 5. Write code for int partition(arr, left, right) as we described in in the lecture, that is, partition rearranges the array so that

- the small entries are stored on positions left, left+1, left+2, ..., pivot\_index-1,
- pivot is stored on position pivot\_index and
- the large entries are stored on pivot\_index+1, pivot\_index+2, ..., right.

The return value is the position of the pivot in **arr**, that is, **pivot\_index**.

```
1 int partition(int[] arr, int left, int right) {
        old_pivot_index = choosePivot(arr, left, right);
2
         \begin{array}{ll} tmpLE = new & int [right - left + 1]; \\ tmpG & = new & int [right - left + 1]; \end{array} 
3
 4
 5
        pivot_index = -1;
 6
        a = 0;
 \overline{7}
        b = 0;
 8
9
        for (int i=0; i<right-left+1; i++) {
10
              if (arr[left+i] <= arr[old_pivot_index]) {</pre>
11
12
13
14
15
16
17
18
              } else {
19
20
              }
21
        }
22
23
        // Copy tmpLE and tmpG into arr:
24
25
26
27
^{28}
29
30
        // Return the location of pivot in arr:
31
        return ???;
32
33 }
```

(Bonus) Question 6. Instead of recursion, use stacks/queues to implement *merge sort*. (You can use function **merge** as in the lecture.)

(Bonus) Question 7. Write partitioning for quicksort which "in-place", that is, it does not need to allocate any extra memory. *Remark:* This can break stability of quicksort.