Trees in Mem (continue)

In Java, in order to represent a binary tree which stores integers, one can define a class

```
1 class Node {
2 Node left;
3 int value;
4 Node right;
5 }
```

In the last tutorial, we said that we can represent trees in Mem. Each node then takes 3 locations in memory. In some sense, we mimicked Java's representation because whenever you write

```
1 \text{ node} = \mathbf{new} \text{ Node}();
```

this corresponds to our:

```
1 node = allocate_memory (3); // then node stores just an *address*, e.g. number 693
```

Furthermore,

1	node.left = \mathbf{null} ; // $empty$	corresponds to	1	Mem[node] = END; // empty
2	node.value = $123;$		2	$\operatorname{Mem}[\operatorname{node}+1] = 123;$
3	node.right = node2;		3	$\operatorname{Mem}[\operatorname{node}+2] = \operatorname{node}2;$
4			4	
5	// assume that node2 has		5	// assume that node2 has
6	// been created earlier		6	// been created earlier

The concrete representation in Mem, might look like as follows. For example, if node = 693 and the second node stores 555 and has no children, we might have:

l	Mem[1]
693	END
694	123
695	698
696	949193
697	419399
698	END
699	555
700	END
701	149939

Exercises:

Solve the following in the Mem representation. You can first write the solution to the exercises in Java and then translate them the same way as above. *Hint: Use recursion, stacks or queues.*

- $5-\frac{1}{2}$. Write a function int size(int root) which computes the number of nodes in the tree.
 - 5. Write a function **int sum(int root)** which computes the sum of all numbers stored in the nodes of the tree.
 - 6. What is the time complexity of your function sum from (5)? Express the time complexity with respect to n = the size of the tree.

Does it make sense to express the time complexity in terms of the tree's height?

7. Bonus: Write a function int maxLessThan(int root, int x) which finds the largest value stored in the binary search tree which is $\leq x$.