## Determining the time complexity 1

The running time of an algorithm usually grows in proportion to the size of the input. To an algorithm, we assign a complexity class which determines the speed of such growth. Let's say that the size of the input is n (e.g. an array consisting of n elements). Then, an algorithm which takes < 4n + 3123 steps is in  $\mathcal{O}(n)$ .

When determining the complexity class, we strip away the (leading) constants. For example, algorithms taking  $\leq 23492n, \leq 31n + 10$ , or  $\leq 2n + 10^{30}$  steps are all in  $\mathcal{O}(n)$ . In general, if the number of steps is less than or equal to an + b, for some constants a and b, then the algorithm is in  $\mathcal{O}(n)$ . Similarly, the algorithms which need  $\leq an^2 + b$  steps belong to  $\mathcal{O}(n^2)$ , or with  $\leq a \log n + b$ steps belong to  $\mathcal{O}(\log n)$ , etc.

1

2

3

4

5

7

8

9

10

2

3

5

6

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9

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12

13

14

15

16

An algorithm which multiplies all elements in the array:

```
int product(int[] arr) {
    int n = arr.length;
2
    int x = 1;
3
4
    int i = 0;
5
    while (i < n) {
6
7
      x = x * arr[i];
8
      i++;
    }
9
10
    return x;
12
```

An algorithm which modifies the last value in the array:

```
void modify(int[] arr) {
      if (\operatorname{arr.length} == 0)
         throw Exception;
      int last = arr[arr.length - 1];
      if (last < 0) \{
         last = -last;
      }
      \operatorname{arr}[\operatorname{arr.length} - 1] = \operatorname{last};
12 }
```

Finding the largest element of the array in three different ways:

```
int largest1(int[] arr) {
2
     int n = arr.length;
     int max = 0;
3
     for (int i=0; i<n; i++) {
5
        bool largest = true;
6
7
        for (int j=0; j<n; j++) {
8
          if \ (\, arr\, [\, i\, ] \ < \ arr\, [\, j\, ]\,)
9
             largest = false;
10
        }
        if (largest)
          \max = \operatorname{arr}[i];
14
     }
16
17
     return max;
18 }
```

```
int largest3(int[] arr) {
    sort(arr);
2
3
    if (arr.length == 0) {
4
5
      return 0;
6
      else {
      int last = arr[arr.length -1];
7
      return last:
8
9
    }
10 }
```

In largest3 assume that sort(arr) is in  $\mathcal{O}(n \log n)$ .

```
1 int largest2(int[] arr) {
     int n = arr.length;
     int max = 0;
     if (arr.length == 0) {
       return 0;
     } else {
       \max = \operatorname{arr}[0];
        for (int i=0; i<n; i++) {
          if (arr[i] > max)
            \max = \operatorname{arr}[i];
        }
        return max;
17
     }
18 }
```

To which complexity classes do those algorithms belong? Possible answers  $\mathcal{O}(1)$ ,  $\mathcal{O}(n), \mathcal{O}(n\log n), \mathcal{O}(n^2).$ 

Procedure	Complexity Class
product	
modify	
largest1	
largest2	
largest3	

Hint: To start with, try what happens if you call product(arr) with arr = [3, 4, 1, 12, 3]. How many steps will it take? And what if arr = [3, 4, 1, 12, 3, 1, 2, 33]? Try the same strategy for the other algorithms as well.

## 2 Comparison of complexity classes

If the number of steps is smaller than  $7n^2 + 2$  then it is also smaller than  $7n^3 + 2$ . This is because, for n getting larger and larger, the number  $n^3$  grows faster than  $n^2$ . Moreover, the constants 7 and 2 can be replaced by any other constants and so any algorithm which is in  $\mathcal{O}(n^2)$  is also in  $\mathcal{O}(n^3)$ . This justifies that we can see  $\mathcal{O}$  to mean "proportional or less".

Example: If an algorithm takes less than  $21n^3 + 6n + 12$  steps then it is in  $\mathcal{O}(n^3)$  because

the number of steps  $\leq 21n^3 + 6n + 12 \leq 21n^3 + 6n^3 + 12 = 27n^3 + 12$ .

Indicate to which complexity class the algorithm belongs based on the number of steps it takes:

number of steps	$\mathcal{O}(1)$	$\mathcal{O}(n)$	$\mathcal{O}(n^2)$	$\mathcal{O}(n^3)$
1	$\checkmark$	$\checkmark$	$\checkmark$	$\checkmark$
$\leq n$	×	~	$\checkmark$	$\checkmark$
$\leq n^2$	×	×	$\checkmark$	$\checkmark$
$\leq 341243$				
$\leq 25n + 72$				
$\leq n^2 + n$				
$\leq n^3 + n^2 + n$				
$\leq 12 + n \times (34 + n)$				
$ \le n^3 + 4 \times (1 + 31n + 4n) $				
$\boxed{ \leq n \times (3n + n \times (83 + n))}$				
$\leq n^5$				
$\leq 2^n$	×	×	×	×

## 3 Exponentials and Logarithms

Recall that, for a number a and a positive integer n, we defined

$$a^n = a \times a \times \dots \times a$$
 (repeated *n* times),  $a^{-n} = \frac{1}{a^n}$  and  $a^{1/n} = \sqrt[n]{a}$ 

We also defined  $\log_a b$  to be the value c such that  $a^c = b$ . Solve the following:

- 1.  $2^2 = ?$ ,  $3^2 = ?$ ,  $4^2 = ?$ ,  $5^2 = ?$ 2.  $2^3 = ?$ ,  $3^3 = ?$ ,  $4^3 = ?$ ,  $5^3 = ?$ 3.  $4^3 = ?$ ,  $4^{-2} = ?$ ,  $4^{1/2} = ?$ ,  $8^{1/3} = ?$
- 4. For which n is  $\frac{16}{2^n} = 1$ ? And when is  $\frac{30}{2^n} \le 1$ ? What is the ceiling  $\lceil \log_2 30 \rceil$ ?
- 5.  $4^{3/2} = ?$ ,  $5^{31}/5^{28} = ?$ ,  $(\sqrt[2]{3})^6 = ?$
- 6.  $\log_3 27 = ?$ ,  $\log_5 \sqrt[3]{5} = ?$ ,  $\log_2 \sqrt[3]{4} = ?$