

Universal structures II

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Apr 29 2020

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Definition (Age of structure L -structure \mathbf{A})

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Definition (Amalgamation class)

Class \mathcal{C} is an **amalgamation class** if it is

- ① closed under isomorphism,
- ② closed under taking induced substructures,
- ③ contains only countably many non-isomorphic structures, and,
- ④ has the amalgamation property.

Definition (Homogeneous structure)

Structure is **homogeneous** if every isomorphism of its two finite substructures extends to an automorphism.

Theorem (Fraïssé 1953)

- A class \mathcal{C} of finite (relational) structures is the age of a countable homogeneous structure \mathbf{H} if and only if \mathcal{C} is an amalgamation class.*
- If conditions of (a) are satisfied then the structure \mathbf{H} is unique up to isomorphism.*

Definition

Let \mathcal{K} be a class of structures. We say that \mathbf{U} is \mathcal{K} -universal if for every $\mathbf{A} \in \mathcal{K}$ there is an embedding $\mathbf{A} \rightarrow \mathbf{U}$.

Example

classes containing universal structure:

- ① Class of all countable graphs, partial orders, rational metric spaces. . .
- ② For every $k \geq 3$ the class of all graphs not containing clique of size k .
- ③ For every $k \geq 3$ the class of all graphs not containing independent of size k .
- ④ Class of all countable graphs with maximal degree 2.

Definition

Given a finite family of structure \mathcal{F} , denote by $\text{Forb}_e(\mathcal{F})$ / $\text{Forb}_m(\mathcal{F})$ / $\text{Forb}_h(\mathcal{F})$ the class of all finite or countable structures \mathbf{A} such that for every $\mathbf{F} \in \mathcal{F}$ there is no embedding / monomorphism / homomorphism $\mathbf{A} \rightarrow \mathbf{F}$ (respectively).

Question

Does every class $\text{Forb}_e(\mathcal{F})$, for \mathcal{F} finite contain an universal structure?

Theorem (Cherlin, Shelah, Shi, 1999; Covington 1990)

Let \mathcal{F} be a finite family of connected structures. Then $\text{Forb}_h(\mathcal{F})$ contains an universal structure.

H.N. 2005: In fact this can be generalized to so called **regular** families \mathcal{F} .

Definition (Minimal separating cut)

A vertex cut R of graph G is **minimal separating** if there exists components C_1 and C_2 of $G \setminus R$ such that $C = N_G(C_1) \cap N_G(C_2)$.

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Definition (Piece of a graph, Hubička, Nešetřil 2005)

Let A be a connected graph and R a minimal separating cut for a component A_1 in A .

A **piece** of A is then a rooted graph

$\mathfrak{P} = (\mathcal{P}, \vec{R})$, where the tuple \vec{R} consists of the vertices of the cut R in a (fixed) linear order and \mathcal{P} is a structure induced by A on $A_1 \cup R$.

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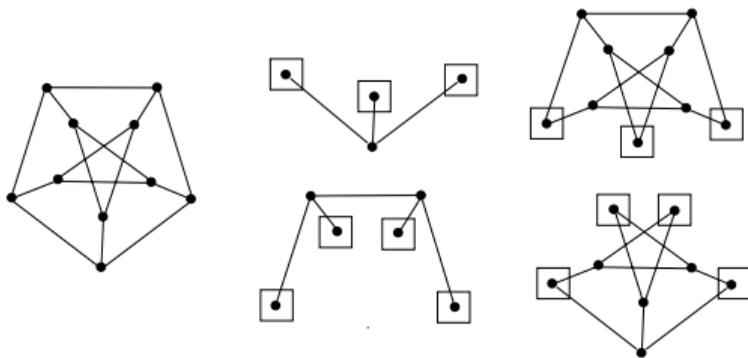
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Basic idea of the proof

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Definition (Language $L_{\mathcal{F}}^+$)

Given finite family \mathcal{F} of finite connected graphs. Enumerate by $\mathfrak{P}_1, \mathfrak{P}_1, \dots, \mathfrak{P}_n$ all pieces of structures in \mathcal{F} . Then language $L_{\mathcal{F}}^+$ extends language of graphs by relations R_1, R_2, \dots, R_n where arity of R_i is the width of \mathfrak{P}_i .

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Definition (Canonical \mathcal{F} -lift)

Given graph $\mathbf{G} \in \text{Forb}_h(\mathcal{F})$ we denote by $L(\mathbf{G})$ the **canonical lift** of \mathbf{G} with is an $L_{\mathcal{F}}^+$ -structure on same vertex set and with same edges. We put $(v_1, v_2, \dots, v_w) \in R_i$ if and only if there exists homomorphism from \mathfrak{P}_i to \mathbf{G} sending roots of \mathfrak{P}_i to v_1, v_2, \dots, v_w .

The proof

Definition (Maximal \mathcal{F} -lift)

The canonical \mathcal{F} -lift $L_{\mathcal{F}}(\mathbf{A})$ is **maximal** on $B \subseteq A$ if for every $\mathbf{C} \in \text{Forb}_h(\mathcal{F})$ such that \mathbf{C} contains \mathbf{A} as substructure, the \mathcal{F} -lift induced on \mathbf{B} by $L_{\mathcal{F}}(\mathbf{A})$ is the same as the \mathcal{F} -lift induced on \mathbf{B} by $L_{\mathcal{F}}(\mathbf{C})$.

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\mathcal{F} -lift \mathbf{X} is **maximal** if there exists $\mathbf{A} \in \text{Forb}_h(\mathcal{F})$ such that \mathbf{X} is induced on \mathbf{X} by $L_{\mathcal{F}}(\mathbf{A})$ and the canonical \mathcal{F} -lift $L_{\mathcal{F}}(\mathbf{A})$ of \mathbf{A} is maximal on \mathbf{X} .

Theorem (Hubička, Nešetřil 2005)

For every finite family \mathcal{F} of connected structures, the class of finite maximal \mathcal{F} -lifts is an amalgamation class.

Recall



$\text{Forb}_H(\mathcal{F})$ -universality



Infinite families



Bounding arities



The proof

Homomorphism universality

Definition

Finite relational structure \mathbf{D} is *dual* of finite relational structure \mathbf{F} iff for every relational structure \mathbf{A} holds:

$$\mathbf{F} \not\rightarrow \mathbf{A} \iff \mathbf{A} \rightarrow \mathbf{D}.$$

In other words, \mathbf{D} is **homomorphism-universal** for the class $\text{Forb}_h(\{\mathbf{F}\})$: For every $\mathbf{A} \in \text{Forb}_h(\{\mathbf{F}\})$ there exists a homomorphism to \mathbf{D} .

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Theorem (Nešetřil, Tardif (2000))

*Directed graph core **G** has dual if and only if **G** is an oriented tree.*

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Theorem (Nešetřil, Tardif (2000))

*Directed graph core **G** has dual if and only if **G** is an oriented tree.*

Finite family \mathcal{F} of finite cores has dual if and only if \mathcal{F} is a family of trees.

Infinite families

Theorem (P. L. Erdős, Pálvölgyi, Tardif, Tardos, 2017)

Family \mathcal{F} of trees has dual if and only if \mathcal{F} is regular.

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Theorem (Myhill-Nerode)

Let L be a set of words over alphabet Σ . Let X and Y be strings. We put $X \sim_L Y$ if and only if for every string Z

$$XZ \in L \iff YZ \in L.$$

L is **regular** (recognizable by a finite automaton) if and only if \sim_L has finitely many equivalence classes.

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Let \mathcal{F} be a family of tree cores. Let \mathfrak{P}_1 and \mathfrak{P}_2 be pieces of \mathcal{F} . We say that $\mathfrak{P}_1 \sim_{\mathcal{F}} \mathfrak{P}_2$ if and only if for every rooted tree \mathfrak{T} it holds

$$\mathfrak{P}_1 \oplus \mathfrak{T} \in \text{Forb}_h(\mathcal{F}) \iff \mathfrak{P}_2 \oplus \mathfrak{T} \in \text{Forb}_h(\mathcal{F}).$$

\mathcal{F} is **regular** if and only if $\sim_{\mathcal{F}}$ has finitely many equivalence classes.

Regular families of structures

Let \mathcal{F} be a family of structures. A piece \mathfrak{P} is **incompatible** with rooted structure \mathfrak{A} if $\mathfrak{P} \oplus \mathfrak{A}$ is isomorphic to some structure in \mathcal{F} .

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Denote by $I_{\mathcal{P}}$ the set of all incompatible structures of \mathfrak{P} and put $\mathfrak{P}_1 \sim_{\mathcal{F}} \mathfrak{P}_2$ if $I_{\mathfrak{P}_1} = I_{\mathfrak{P}_2}$.

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Family of structures \mathcal{F} is **regular** iff for every $n \geq 1$ $\sim_{\mathcal{F}}$ has only finitely many equivalence classes of pieces of \mathcal{F} of width n .

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Family of structures \mathcal{F} is **regular** iff for every $n \geq 1$ $\sim_{\mathcal{F}}$ has only finitely many equivalence classes of pieces of \mathcal{F} of width n .

Family \mathcal{F} is upwards closed if for every $\mathbf{F} \in \mathcal{F}$ we also have $\mathbf{F}' \in \mathcal{F}$ provided that \mathcal{F} is connected and there is a homomorphism $\mathbf{F} \rightarrow \mathbf{F}'$.

Theorem (H., Nešetřil 2005)

Let L be a finite relational language. Let \mathcal{F} be a upwards closed family of finite connected L -structures. Then the following conditions are equivalent:

- ① \mathcal{F} is a regular family of connected structures.
- ② $\text{Forb}_h(\mathcal{F})$ contains an ω -categorical universal structure.

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Bounding arities

Theorem (H., Nešetřil 2005)

Let \mathcal{F} be a finite set of cores such that there is no $\mathbf{F}_1 \neq \mathbf{F}_2 \in \mathcal{F}$ with a homomorphism $\mathbf{F}_1 \rightarrow \mathbf{F}_2$. Then every homogenization of $\text{Forb}_h(\mathcal{F})$ in finite language uses relations of arity n where n is the maximal number of roots of a piece of structure in \mathcal{F} .

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