

Universal structures

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By $\text{Age}(\mathbf{A})$ we denote the class of all L -structures isomorphic to some (induced) substructure of \mathbf{A} .

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Class \mathcal{C} is an **amalgamation class** if it is

- ① closed under isomorphism,
- ② closed under taking induced substructures,
- ③ contains only countably many non-isomorphic structures, and,
- ④ has the amalgamation property.

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- 4 has the amalgamation property.

Definition (Homogeneous structure)

Structure is **homogeneous** if every isomorphism of its two finite substructures extends to an automorphism.

Theorem (Fraïssé 1953)

- (a) A class \mathcal{C} of finite (relational) structures is the age of a countable homogeneous structure \mathbf{H} if and only if \mathcal{C} is an amalgamation class.
- (b) If conditions of (a) are satisfied then the structure \mathbf{H} is unique up to isomorphism.

Definition

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Definition

Given a finite family of structure \mathcal{F} , denote by $\text{Forb}_e(\mathcal{F})$ / $\text{Forb}_m(\mathcal{F})$ / $\text{Forb}_h(\mathcal{F})$ the class of all finite or countable structures \mathbf{A} such that for every $\mathbf{F} \in \mathcal{F}$ there is no embedding / monomorphism / homomorphism $\mathbf{A} \rightarrow \mathbf{F}$ (respectively).

Question

Does every class $\text{Forb}_e(\mathcal{F})$, for \mathcal{F} finite contain an universal structure?

Non-existence of universal graphs

Theorem (Füredi, Komjáth, 1997)

If \mathbf{G} is a 2-connected graph that is not complete, then $\text{Forb}_m(\{\mathbf{G}\})$ contains no universal graph.

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Lemma

For every k and g there exist a $t = t(k, g)$ and a countable k -uniform hypergraph \mathbf{H}_g^k of girth g , vertex set ω and edges E_1, E_2, \dots such that $(i + t) \in E_i \subseteq \{i - t, \dots, i + t\}$.

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Put $k = 2|G| - 1, g = |G| + 1$ and let \mathbf{H} be given by the lemma.

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If copies of graphs $\mathbf{G}(\epsilon)$ and $\mathbf{G}(\epsilon')$ overlap by vertices $1, 2, \dots, t$ then they overlap completely.

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Theorem (Cherlin, 1996)

Let \mathcal{C} be family of cycles, $\text{Forb}_m(\mathcal{C})$ contains an universal graph if and only if \mathcal{C} is a set of odd cycles of length $1, 3, 5, \dots, k$ for some k .

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A vertex cut R of graph G is **minimal separating** if there exists components C_1 and C_2 of $G \setminus R$ such that $C = N_G(C_1) \cap N_G(C_2)$.

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Let \mathbf{A} be a connected graph and R a minimal separating cut for a component \mathbf{A}_1 in \mathbf{A} .

A **piece** of \mathbf{A} is then a rooted graph

$\mathfrak{P} = (\mathbf{P}, \vec{R})$, where the tuple \vec{R} consists of the vertices of the cut R in a (fixed) linear order and \mathbf{P} is a structure induced by \mathbf{A} on $\mathbf{A}_1 \cup R$.

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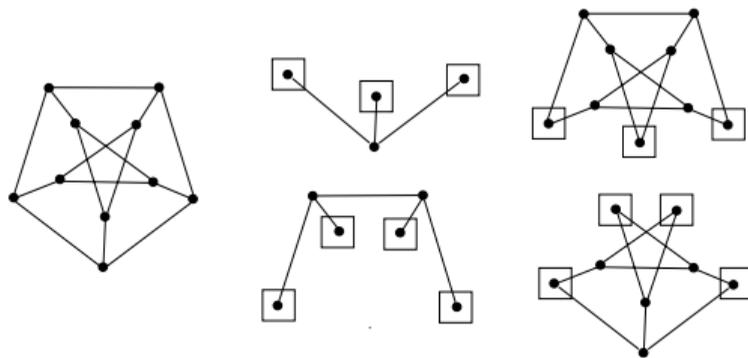
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Definition (Language $L_{\mathcal{F}}^+$)

Given finite family \mathcal{F} of finite connected graphs. Enumerate by $\mathfrak{P}_1, \mathfrak{P}_1, \dots, \mathfrak{P}_n$ all pieces of structures in \mathcal{F} . Then language $L_{\mathcal{F}}^+$ extends language of graphs by relations R_1, R_2, \dots, R_n where arity of R_i is the width of \mathfrak{P}_i .

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Definition (Canonical \mathcal{F} -lift)

Given graph $\mathbf{G} \in \text{Forb}_h(\mathcal{F})$ we denote by $L(\mathbf{G})$ the **canonical lift** of \mathbf{G} with is an $L_{\mathcal{F}}^+$ -structure on same vertex set and with same edges. We put $(v_1, v_2, \dots, v_w) \in R_i$ if and only if there exists homomorphism from \mathfrak{P}_i to \mathbf{G} sending roots of \mathfrak{P}_i to v_1, v_2, \dots, v_w .

The proof

Definition (Maximal \mathcal{F} -lift)

The canonical \mathcal{F} -lift $L_{\mathcal{F}}(\mathbf{A})$ is **maximal** on $B \subseteq A$ if for every $\mathbf{C} \in \text{Forb}_h(\mathcal{F})$ such that \mathbf{C} contains \mathbf{A} as substructure, the \mathcal{F} -lift induced on \mathbf{B} by $L_{\mathcal{F}}(\mathbf{A})$ is the same as the \mathcal{F} -lift induced on \mathbf{B} by $L_{\mathcal{F}}(\mathbf{C})$.

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\mathcal{F} -lift \mathbf{X} is **maximal** if there exists $\mathbf{A} \in \text{Forb}_h(\mathcal{F})$ such that \mathbf{X} is induced on \mathbf{X} by $L_{\mathcal{F}}(\mathbf{A})$ and the canonical \mathcal{F} -lift $L_{\mathcal{F}}(\mathbf{A})$ of \mathbf{A} is maximal on \mathbf{X} .

Theorem (Hubička, Nešetřil 2005)

For every finite family \mathcal{F} of connected structures, the class of finite maximal \mathcal{F} -lifts is an amalgamation class.

The proof

Definition

Finite relational structure **D** is *dual* of finite relational structure **F** iff for every relational structure **A** holds:

$$\mathbf{F} \not\rightarrow \mathbf{A} \iff \mathbf{A} \rightarrow \mathbf{D}.$$

Theorem (Nešetřil, Tardif (2000))

*Directed graph core **G** has dual if and only if **G** is an oriented tree.*

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Theorem (P. L. Erdős, Pálvölgyi, Tardif, Tardos, 2017)

Family \mathcal{F} of trees has dual if and only if \mathcal{F} is regular.

