

An interval linear programming contractor

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Motivation

Interval data are used to model:

- real life uncertainties
- measurement errors
- sensitivity analysis

Notation

An interval matrix

$$\mathbf{A} := [\underline{A}, \overline{A}] = \{A \in \mathbb{R}^{m \times n} \mid \underline{A} \leq A \leq \overline{A}\}.$$

The center and radius matrices

$$A_c := \frac{1}{2}(\overline{A} + \underline{A}), \quad A_\Delta := \frac{1}{2}(\overline{A} - \underline{A}).$$

Interval linear programming

Consider a linear programming problem

$$\min c^T x \text{ subject to } Ax = b, x \geq 0,$$

where $A \in \mathbf{A}$, $b \in \mathbf{b}$, $c \in \mathbf{c}$.

State of the art

- optimal value range (Chinneck & Ramadan, 2000, Hladík, 2009, Jansson, 2004, Mráz, 1998, Rohn, 2006, etc.)
- duality (Gabrel & Murat, 2010, Rohn, 1980, Serafini, 2005)
- basis stability (Beeck, 1978, Koníčková, 2001, Hladík, 2010, Rohn, 1993)
- optimal solution set (Beeck, 1978, Jansson, 1988, Machost, 1970)

Problem statement

The optimal solution set

Denote by $\mathcal{S}(A, b, c)$ the set of optimal solutions to

$$\min c^T x \quad \text{subject to} \quad Ax = b, \quad x \geq 0,$$

Then the optimal solution set is defined

$$\mathcal{S} := \bigcup_{A \in \mathbf{A}, b \in \mathbf{b}, c \in \mathbf{c}} \mathcal{S}(A, b, c).$$

Goal

Find a tight enclosure to \mathcal{S} .

Approaches

- Interval arithmetic (conservative)

Characterization

By duality theory, we have that $x \in \mathcal{S}$ if and only if there is some $y \in \mathbb{R}^m$, $A \in \mathbf{A}$, $b \in \mathbf{b}$, and $c \in \mathbf{c}$ such that

$$Ax = b, x \geq 0, A^T y \leq c, c^T x = b^T y$$

where $A \in \mathbf{A}$, $b \in \mathbf{b}$, $c \in \mathbf{c}$.

Relaxation

$$Ax = b, x \geq 0, A'^T y \leq c, c'^T x = b'^T y,$$

where $A, A' \in \mathbf{A}$, $b, b' \in \mathbf{b}$, $c, c' \in \mathbf{c}$.

Our approach

Description of the relaxed problem

$$\begin{aligned}\underline{A}x &\leq \overline{b}, \\ -\overline{A}x &\leq -\underline{b}, \\ x &\geq 0, \\ A_c^T y - A_\Delta^T |y| &\leq \overline{c}, \\ |c_c^T x - b_c^T y| &\leq c_\Delta^T x + b_\Delta^T |y|.\end{aligned}$$

Properties

- The solution set is non-convex in general
- It is linear at any orthant
- NP-hard to obtain exact bounds

Idea

Linearize $|y|$.

Theorem (Beaumont, 1998)

For every $y \in \mathbf{y} \subset \mathbb{R}$ with $\underline{y} < \bar{y}$ one has

$$|y| \leq \alpha y + \beta, \quad (1)$$

where

$$\alpha = \frac{|\bar{y}| - |\underline{y}|}{\bar{y} - \underline{y}} \quad \text{and} \quad \beta = \frac{\bar{y}|\underline{y}| - \underline{y}|\bar{y}|}{\bar{y} - \underline{y}}.$$

Moreover, if $\underline{y} \geq 0$ or $\bar{y} \leq 0$ then (1) holds as equation.

Linearization of $|y|$

Now, the linearization reads

$$\begin{aligned}\underline{A}x &\leq \underline{b}, \quad -\overline{A}x \leq -\underline{b}, \quad x \geq 0, \\ (A_c^T - A_\Delta^T \text{diag}(\alpha))y &\leq \overline{c} + A_\Delta^T \beta, \\ \underline{c}^T x + (-b_c^T - b_\Delta^T \text{diag}(\alpha))y &\leq b_\Delta^T \beta, \\ -\overline{c}^T x + (b_c^T - b_\Delta^T \text{diag}(\alpha))y &\leq b_\Delta^T \beta,\end{aligned}$$

where

$$\alpha_i := \begin{cases} \frac{|\overline{y}_i| - |\underline{y}_i|}{\overline{y}_i - \underline{y}_i} & \text{if } \underline{y}_i < \overline{y}_i, \\ \text{sgn}(\overline{y}_i) & \text{if } \underline{y}_i = \overline{y}_i, \end{cases}$$
$$\beta_i := \begin{cases} \frac{\overline{y}_i |\underline{y}_i| - \underline{y}_i |\overline{y}_i|}{\overline{y}_i - \underline{y}_i} & \text{if } \underline{y}_i < \overline{y}_i, \\ 0 & \text{if } \underline{y}_i = \overline{y}_i. \end{cases}$$

Algorithm (Optimal solution set contractor)

- ❶ Compute an initial interval enclosure $\mathbf{x}^0, \mathbf{y}^0$
- ❷ $i := 0$;
- ❸ **repeat**
 - ❶ compute the interval hull $\mathbf{x}^i, \mathbf{y}^i$ of the linearized system;
 - ❷ $i := i + 1$;
- ❹ **until** improvement is nonsignificant;
- ❺ **return** \mathbf{x}^i ;

Properties

- Each iteration requires solving an interval hull ($2n$ linear programs).
- In practice, it converges quickly, but not to \mathcal{S} in general.

Problems

- How to determine an initial enclosure $\mathbf{x}^0, \mathbf{y}^0$?

Example

Consider an interval linear program

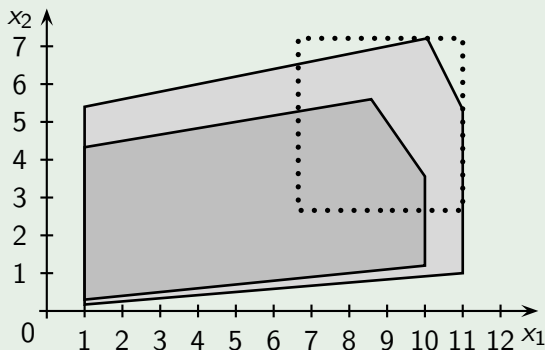
$$\begin{aligned} \min & -[15, 16]x_1 - [17, 18]x_2 \quad \text{subject to} \\ & x_1 \leq [10, 11], \\ & -x_1 + [5, 6]x_2 \leq [25, 26], \\ & [6, 6.5]x_1 + [3, 4.5]x_2 \leq [81, 82], \\ & -x_1 \leq -1, \\ & x_1 - [10, 12]x_2 \leq -[1, 2]. \end{aligned}$$

Take the initial enclosure

$$\begin{aligned} \mathbf{x}^0 &= 1000 \cdot ([-1, 1], [-1, 1])^T, \\ \mathbf{y}^0 &= 1000 \cdot ([0, 1], [0, 1], [0, 1], [0, 1], [0, 1])^T. \end{aligned}$$

Example

Example (cont.)



- In grey the largest and the smallest feasible area.
- The final enclosure of the optimal solution set \mathcal{S} is dotted.

Application: portfolio selection

Given:

- J possible investments;
- T time periods;
- r_{jt} , return on investment j in time period t ;
- μ , risk aversion parameter (upper bound for risk).

Then:

- Estimated reward on investment j : $R_j := \frac{1}{T} \sum_{t=1}^T r_{jt}$;
- Risk measure of investment j : $\frac{1}{T} \sum_{t=1}^T |r_{jt} - R_j|$;
- Maximal allowed risk: $\frac{1}{T} \sum_{t=1}^T \left| \sum_{j=1}^J (r_{jt} - R_j) x_j \right| \leq \mu$.

Application: portfolio selection

Portfolio selection problem formulation

$$\begin{aligned} \max \quad & \sum_{j=1}^J R_j x_j \\ \text{subject to} \quad & -y_j \leq \sum_{j=1}^J (r_{jt} - R_j) x_j \leq y_t, \quad \forall t = 1, \dots, T, \\ & \sum_{j=1}^J x_j = 1, \quad \frac{1}{T} \sum_{t=1}^T y_t \leq \mu, \\ & x_j \geq 0, \quad \forall j = 1, \dots, J, \end{aligned}$$

where

$$R_j := \frac{1}{T} \sum_{t=1}^T r_{jt}.$$

Example

Example

$J = 4$ investments, $T = 5$ time periods, $\mu = 2$ risk aversion parameter.

The returns:

time period t	return on investment			
	1	2	3	4
1	10	20	9	11
2	12	25	11	14
3	9	17	12	12
4	11	21	11	14
5	11	19	13	16

- The optimal return: 20.08.
- The optimal solution: $x^* = (0, 0.9643, 0.0357, 0)^T$

Example

Suppose that the returns are not known precisely:

- 1 (1% tolerance) extending r_{tj} to intervals $[0.99r_{tj}, 1.01r_{tj}]$, we calculate:

$$\mathbf{x}^{(1)} = ([0, 0.1699], [0.7621, 1], [0, 0.181], [0, 0.2379])^T.$$

- 2 (5% tolerance) replace r_{jt} by intervals $[0.95r_{tj}, 1.05r_{tj}]$, $j \neq 2$

$$\mathbf{x}^{(2)} = ([0, 0.0495], [0.9276, 0.9712], [0, 0.0531], [0, 0.0724])^T.$$

Conclusion

- Effective contractor for the optimal solution set \mathcal{S} .
- Each iteration requires solving $2n$ linear programs.
- In practice, it converges quickly.

Future work

- Initial enclosure of the optimal solution set \mathcal{S} .