

Computing the range of real eigenvalues of an interval matrix

Milan Hladík & David Daney

INRIA, project Coprin
Sophia Antipolis, France



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Notation

An interval matrix

$$\mathbf{A} = [\underline{A}, \overline{A}] = \{A \in \mathbb{R}^{n \times n} \mid \underline{A} \leq A \leq \overline{A}\},$$

Midpoint and radius of \mathbf{A}

$$A_c := \frac{1}{2}(\underline{A} + \overline{A}), \quad A_\Delta := \frac{1}{2}(\overline{A} - \underline{A}).$$

Aim

Determine eigenvalue set

$$\Lambda := \{\lambda \in \mathbb{R}; Ax = \lambda x, x \neq 0, A \in \mathbf{A}\}$$

Complexity

Checking $0 \in \Lambda$ is NP-hard.

Algorithm (Approximation of Λ)

- 1 Compute initial bounds $\underline{\lambda}^0, \bar{\lambda}^0$ such that $\Lambda \subseteq \boldsymbol{\lambda}^0 := [\underline{\lambda}^0, \bar{\lambda}^0]$;
- 2 $L := \{\boldsymbol{\lambda}^0\}$, $L_{\text{inn}} := \emptyset$, $L_{\text{unc}} := \emptyset$;
- 3 while $L \neq \emptyset$ do
 - 1 choose and remove some $\boldsymbol{\lambda}$ from L ;
 - 2 (OUTER TEST)
if $\boldsymbol{\lambda} \cap \Lambda = \emptyset$ then loop;
 - 3 (INNER TEST)
if $\boldsymbol{\lambda} \subseteq \Lambda$ then $L_{\text{inn}} := L_{\text{inn}} \cup \{\boldsymbol{\lambda}\}$ and loop;
 - 4 (PRECISION TEST)
if $\lambda_{\Delta} < \varepsilon$ then $L_{\text{unc}} := L_{\text{unc}} \cup \{\boldsymbol{\lambda}\}$ and loop;
 - 5 (DIVISION)
 $\boldsymbol{\lambda}^1 := [\underline{\lambda}, \lambda_c]$, $\boldsymbol{\lambda}^2 := [\lambda_c, \bar{\lambda}]$, $L := L \cup \{\boldsymbol{\lambda}^1, \boldsymbol{\lambda}^2\}$ and loop;
- 4 return L_{inn} as inner and L_{unc} as unrecognized intervals;

Implementation of the list L

Implemented as a stack (LIFO), useful properties:

- neighbor of an inner box cannot be outer
- neighbor of an outer box cannot be inner

Initial bounds (Rohn, 1998)

Let

$$S_c := \frac{1}{2} (A_c + A_c^T), \quad S_\Delta := \frac{1}{2} (A_\Delta + A_\Delta^T).$$

Then $\Lambda \subseteq \boldsymbol{\lambda}^0 := [\underline{\lambda}^0, \bar{\lambda}^0]$, where

$$\begin{aligned}\underline{\lambda}^0 &= \lambda_{\min}(S_c) - \rho(S_\Delta), \\ \bar{\lambda}^0 &= \lambda_{\max}(S_c) + \rho(S_\Delta).\end{aligned}$$

Remark

λ is eigenvalue of A iff $A - \lambda I$ is singular.

Proposition

If the interval matrix $\mathbf{A} - \lambda I$ is regular, then λ is an outer interval.

Idea

Put $\mathbf{M} := \mathbf{A} - \lambda I$ and check regularity of \mathbf{M} .

Theorem (Sufficient regularity condition)

An interval matrix \mathbf{M} is regular if M_c is nonsingular and $\rho(|M_c^{-1}|M_\Delta) < 1$.

Outer test: Jansson & Rohn method

Theorem (Jansson & Rohn, 1999)

Let $b \in \mathbb{R}^n$ and consider $\mathbf{M}x = b$. The solution set is described by

$$|M_c x - b| \leq M_\Delta |x|.$$

Then \mathbf{M} is regular if and only if its any component is bounded.

Algorithm (Basic scheme)

- 1 Select $b \in \mathbb{R}^n$ and the initial orthant;
- 2 Check for unboundedness. If so, then \mathbf{M} is not regular;
- 3 Repeat for the neighboring orthants, and so on;
- 4 If inspected the whole component, then \mathbf{M} is regular.

Properties

- Not necessarily exponential complexity.

Outer test: ILS method

Proposition

\mathbf{M} is regular iff there is no solution of the interval system

$$\mathbf{M}x = 0, \|x\|_{\infty} = 1.$$

Algorithm (ILS method)

- ① For $i = 1, \dots, n$ do
 - ① $\mathbf{b} :=$ (the i -th column of \mathbf{M}), $\mathbf{M}' :=$ (\mathbf{M} without the i -th column);
 - ② Solve (approximately) the interval linear system
$$\mathbf{M}'x' = -\mathbf{b}, -e \leq x' \leq e;$$
 - ③ If it has possibly a solution, then stop, \mathbf{M} needn't be regular;
- ② Conclude that \mathbf{M} is regular.

Properties

- We obtain also estimation of eigenvectors.

Algorithm (Inner test for λ)

- 1 Call Jansson & Rohn algorithm with $\mathbf{M} := \mathbf{A} - \lambda_c I$;
- 2 If \mathbf{M} is regular then stop, λ is not inner interval;
- 3 else, let $z \in \{\pm 1\}^n$ be a sign vector for the “unbounded orthant”;
- 4 Solve linear program

$$\begin{aligned} \max \{ & z^T x^1 - z^T x^2; (A_c - A_\Delta \text{diag}(z))(x^1 - x^2) - \underline{\lambda}x^1 + \bar{\lambda}x^2 \leq b, \\ & (A_c + A_\Delta \text{diag}(z))(x^1 - x^2) - \bar{\lambda}x^1 + \underline{\lambda}x^2 \geq b, \\ & \text{diag}(z)(x^1 - x^2) \geq 0, x^1, x^2 \geq 0 \} \end{aligned}$$

- 5 If it is unbounded then λ is an inner interval.

Improvements

- Remember the sign vector z for the following test.

Theorem (Rohn, 1993)

Let $\lambda \in \partial\Lambda$. Then there are $x, p \in \mathbb{R}^n \setminus \{0\}$ and $y, z \in \{\pm 1\}^n$ such that

$$\begin{aligned}(A_c - \text{diag}(y) A_\Delta \text{diag}(z))x &= \lambda x, \\ (A_c^T - \text{diag}(z) A_\Delta^T \text{diag}(y))p &= \lambda p, \\ \text{diag}(z)x &\geq 0, \quad \text{diag}(y)p \geq 0.\end{aligned}$$

Algorithm (Exact bound candidates for λ)

- 1 Call ILS method with the input matrix $\mathbf{A} - \lambda I$ to obtain an outer approximation \mathbf{x} of the right eigenvectors;
- 2 Call ILS method with the input matrix $(\mathbf{A} - \lambda I)^T$ to obtain an outer approximation \mathbf{p} of the left eigenvectors;
- 3 If the number of possible signs of \mathbf{x} and \mathbf{p} is small then enumerate all possibilities.

Numerical experiments

n	ε	R	time
5	0.1	1	2 sec
10	0.1	0.5	9 sec
10	0.1	5	1 min 12 sec
15	0.1	0.1	37 sec
15	0.1	0.5	10 min 29 sec
15	0.1	1	7 min 59 sec
20	0.1	0.1	2 min 16 sec
20	0.1	0.5	21 min 6 sec
25	0.1	0.01	5 min 46 sec
25	0.1	0.05	10 min 39 sec
30	0.01	0.01	14 min 37 sec
30	0.01	0.1	48 min 31 sec

- n =matrix dimension
- ε =accuracy
- random matrices **A**
- entries of A_c in uniform distribution $[-20, 20]$
- entries of A_Δ in uniform distribution $[0, R]$
- exact bounds found
- not rigorous:
GLPK (linear programs)
CLAPACK (linear algebra)

Numerical experiments

- red = eigenvalue set Λ ,
- yellow = sufficient outer test condition,
- green = Jansson & Rohn method



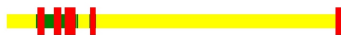
Random matrix,
 $n = 30$, $R = 0.1$,
computing time 48 min 31 sec.



Random symmetric matrix,
 $n = 15$, $R = 0.5$,
computing time 13 min 48 sec.



Random $\mathbf{A}^T \mathbf{A}$ matrix,
 $n = 15$, $R = 0.02$,
computing time 1 min 58 sec.



Random nonnegative matrix,
 $n = 15$, $R = 0.2$,
computing time 2 min 22 sec.

Other approaches

- Interval analysis techniques (filtering, branch & prune, ...) for

$$\mathbf{A}x = \lambda x, \quad \|x\|_{\infty} = 1, \quad \lambda \in \lambda^0.$$

Efficient for $n < 5$.

- Gerschgorin discs, Cassini ovals, ... no sharp results.

Present and future work

- Eigenvalues of symmetric interval matrix;
- Singular values of interval matrix;
- Complex eigenvalues of interval matrix.