

# Polynomials

**Definition:** A *polynomial* of degree  $n$  in variable  $x$  over a field  $\mathbb{K}$  is an expression  $p(x) = a_n x^n + a_{n-1} x^{n-1} + \cdots + a_2 x^2 + a_1 x + a_0$ , where  $a_n \neq 0$  and  $a_n, \dots, a_0 \in \mathbb{K}$ . We write  $p \in \mathbb{K}(x)$ .

Operations on polynomials  $p(x) = \sum_{i=0}^n a_i x^i$ ,  $q(x) = \sum_{i=0}^m b_i x^i$ :

- ▶ addition, subtraction:  $(p \pm q)(x) = \sum_{i=0}^{\max n, m} (a_i \pm b_i) x^i$
- ▶ scalar multiple:  $(\alpha p)(x) = \sum_{i=0}^n (\alpha a_i) x^i$
- ▶ product  $(pq)(x) = \sum_{i=0}^{n+m} c_i x^i$ , where  $c_i = \sum_{j=0}^i a_j b_{i-j}$
- ▶ division with a remainder — there are unique polynomials  $r, t \in \mathbb{K}(x)$  such that  $p = qt + r$ , where the degree of  $r$  is less than  $m$ , the degree of  $q$ .

## Example — operations on polynomials over $\mathbb{Z}_5$

Addition:

$$(3x^3 + 2x + 1) + (2x^2 + 3x + 1) = 3x^3 + 2x^2 + 2$$

the degree may decrease:

$$(3x^3 + 2x + 1) + (2x^3 + 3x + 1) = 2$$

Multiple:

$$2 \cdot (3x^3 + 2x + 1) = x^3 + 4x + 2$$

Product:

$$(3x^3 + 2x + 1)(2x^2 + 3x + 1) = x^5 + 4x^4 + 2x^3 + 3x^2 + 1$$

## Example — operations on polynomials over $\mathbb{Z}_5$

Division with the remainder:

$$\begin{array}{r} 4x^5 + 2x^4 + 3x^2 + 3 \\ -4x^5 - 2x^4 - x^3 \\ \hline \phantom{4x^5 + 2x^4} 4x^3 + 3x^2 \\ \phantom{4x^5 + 2x^4} -4x^3 - 2x^2 - x \\ \hline \phantom{4x^5 + 2x^4} \phantom{4x^3 + 3x^2} x^2 + 4x + 3 \\ \phantom{4x^5 + 2x^4} \phantom{4x^3 + 3x^2} - x^2 - 3x - 4 \\ \hline \phantom{4x^5 + 2x^4} \phantom{4x^3 + 3x^2} \phantom{x^2 + 4x + 3} x + 4 \end{array} : 3x^2 + 4x + 2 = 3x^3 + 3x + 2$$

Correctness check  $p = qt + r$ :

$$4x^5 + 2x^4 + 3x^2 + 3 = (3x^2 + 4x + 2)(3x^3 + 3x + 2) + (x + 4)$$

## Fermat's little theorem

Theorem: For any  $x \in \mathbb{Z}_p \setminus \{0\} : x^{p-1} = 1$ .

Proof: The map  $i \rightarrow xi$  is a bijection on  $\{1, \dots, p-1\}$  in  $\mathbb{Z}_p$ .

In  $\prod_{i=1}^{p-1} i = \prod_{i=1}^{p-1} xi = x^{p-1} \prod_{i=1}^{p-1} i$  cancel the nonzero term  $\prod_{i=1}^{p-1} i$ .

Corollary: For any  $x \in \mathbb{Z}_p : x^p - x = 0$ .

Corollary: For any  $q \in \mathbb{Z}_p(x)$  there is  $r \in \mathbb{Z}_p(x)$  of degree at most  $p-1$ , such that  $\forall x \in \mathbb{Z}_p : q(x) = r(x)$ .

Example:

$$4x^5 + 2x^4 + 3x^2 + 3 = 4(x^5 - x) + 2x^4 + 3x^2 + 4x + 3$$

i.e. the polynomial  $q(x) = 4x^5 + 2x^4 + 3x^2 + 3$  yields on  $\mathbb{Z}_5$  the same values as  $r(x) = 2x^4 + 3x^2 + 4x + 3$ .

# Roots

**Definition:** The *root* of a polynomial  $p \in \mathbb{K}(x)$  is  $r \in \mathbb{K}$  such that  $p(r) = 0$ .

**Observation:** The element  $r \in \mathbb{K}$  is a root of a polynomial  $p$  if and only if the linear polynomial  $x - r$  divides  $p$  without a remainder.

**Definition:** The *multiplicity* of the root  $r$  of  $p \in \mathbb{K}(x)$  is the maximum positive integer  $k$  such that  $(x - r)^k$  divides  $p$ .

**Theorem:** (The fundamental theorem of algebra)  
Every polynomial  $p \in \mathbb{C}(x)$  has at least one root.

**Corollary:** Every polynomial  $p \in \mathbb{C}(x)$  can be factorized into linear factors, i.e. polynomials of degree one.

**Definition:** If every  $p \in \mathbb{K}(x)$  of degree at least 1 has a root, then the field  $\mathbb{K}$  is *algebraically closed*.

## Representations of polynomials of degree $n$

- ▶ by the coefficients  $a_0, \dots, a_n$
- ▶ in algebraically closed fields by  $a_n$  and the  $n$  roots  $r_1, \dots, r_n$
- ▶ by the values of the polynomial in  $n + 1$  distinct points

**Problem:** Given  $n + 1$  pairs  $(x_i, y_i)$  for  $i = 0, \dots, n$ , determine  $p \in \mathbb{K}(x)$  of degree at most  $n$  such that  $p(x_i) = y_i$  for all  $i$ .

**Observation:** Coefficients  $a_0, \dots, a_n$  of  $p$  are solution of the system:

$$\begin{pmatrix} 1 & x_0 & x_0^2 & \dots & x_0^n \\ 1 & x_1 & x_1^2 & \dots & x_1^n \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & x_n & x_n^2 & \dots & x_n^n \end{pmatrix} \begin{pmatrix} a_0 \\ a_1 \\ \vdots \\ a_n \end{pmatrix} = \begin{pmatrix} y_0 \\ y_1 \\ \vdots \\ y_n \end{pmatrix}$$

**Definition:** The matrix of this system is the *Vandermonde matrix*  $\mathbf{V}_{n+1}(x_0, \dots, x_n)$

**Theorem:** The Vandermonde matrix  $\mathbf{V}_{n+1}(x_0, \dots, x_n)$  is regular if and only if  $x_0, \dots, x_n$  are distinct.

## Proof of the regularity of the Vandermonde matrix

$$\mathbf{V}_{n+1}(x_0, \dots, x_n) = \begin{pmatrix} 1 & x_0 & x_0^2 & \dots & x_0^n \\ 1 & x_1 & x_1^2 & \dots & x_1^n \\ 1 & x_2 & x_2^2 & \dots & x_2^n \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & x_n & x_n^2 & \dots & x_n^n \end{pmatrix}$$

Deduct the first row from others.  
Then factor out  $x_i - x_0$  from the  $i$ -th row for each  $i = 1, \dots, n$ .  
In the first column are  $n$  zeros so we can expand:  $\det(\mathbf{V}_{n+1}) =$

$$= \prod_{i=1}^n (x_i - x_0) \cdot \begin{vmatrix} 1 & x_1 + x_0 & x_1^2 + x_1 x_0 + x_0^2 & \dots & x_1^{n-1} + x_1^{n-2} x_0 + \dots + x_0^{n-1} \\ 1 & x_2 + x_0 & x_2^2 + x_2 x_0 + x_0^2 & \dots & x_2^{n-1} + x_2^{n-2} x_0 + \dots + x_0^{n-1} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & x_n + x_0 & x_n^2 + x_n x_0 + x_0^2 & \dots & x_n^{n-1} + x_n^{n-2} x_0 + \dots + x_0^{n-1} \end{vmatrix}$$

Now *backward* subtract from every column the  $x_0$ -multiple of the previous one.  
By this we eliminate all terms containing  $x_0$ .

Consequently, we get a recurrence that could be expanded as follows:

$$\begin{aligned} \det(\mathbf{V}_{n+1}(x_0, \dots, x_n)) &= \left( \prod_{i=1}^n (x_i - x_0) \right) \det(\mathbf{V}_n(x_1, \dots, x_n)) \\ &= \prod_{i < j} (x_j - x_i) \end{aligned}$$

## Example for $n = 3$

$$\det(\mathbf{V}_4(x_0, \dots, x_3)) =$$

$$= \begin{vmatrix} 1 & x_0 & x_0^2 & x_0^3 \\ 1 & x_1 & x_1^2 & x_1^3 \\ 1 & x_2 & x_2^2 & x_2^3 \\ 1 & x_3 & x_3^2 & x_3^3 \end{vmatrix} \begin{matrix} \text{I} \\ \text{II} \\ \text{III} \\ \text{IV} \end{matrix} = \begin{vmatrix} 1 & x_0 & x_0^2 & x_0^3 \\ 0 & x_1 - x_0 & x_1^2 - x_0^2 & x_1^3 - x_0^3 \\ 0 & x_2 - x_0 & x_2^2 - x_0^2 & x_2^3 - x_0^3 \\ 0 & x_3 - x_0 & x_3^2 - x_0^2 & x_3^3 - x_0^3 \end{vmatrix} \begin{matrix} \text{I} \\ \text{II} - \text{I} \\ \text{III} - \text{I} \\ \text{IV} - \text{I} \end{matrix}$$

$$= \begin{vmatrix} x_1 - x_0 & x_1^2 - x_0^2 & x_1^3 - x_0^3 \\ x_2 - x_0 & x_2^2 - x_0^2 & x_2^3 - x_0^3 \\ x_3 - x_0 & x_3^2 - x_0^2 & x_3^3 - x_0^3 \end{vmatrix} \begin{matrix} : (x_1 - x_0) \\ : (x_2 - x_0) \\ : (x_3 - x_0) \end{matrix}$$

$$= (x_1 - x_0)(x_2 - x_0)(x_3 - x_0) \begin{vmatrix} 1 & x_1 + x_0 & x_1^2 + x_1x_0 + x_0^2 \\ 1 & x_2 + x_0 & x_2^2 + x_2x_0 + x_0^2 \\ 1 & x_3 + x_0 & x_3^2 + x_3x_0 + x_0^2 \end{vmatrix} \begin{matrix} \text{I} & \text{II} & \text{III} \end{matrix}$$

$$= \prod_{i=1}^3 (x_i - x_0) \begin{vmatrix} 1 & x_1 & x_1^2 \\ 1 & x_2 & x_2^2 \\ 1 & x_3 & x_3^2 \end{vmatrix} \begin{matrix} \text{I} & \text{II} - x_0\text{I} & \text{III} - x_0\text{II} \end{matrix} = \left( \prod_{i=1}^3 (x_i - x_0) \right) \det(\mathbf{V}_3(x_1, x_2, x_3))$$



# Lagrange interpolation

... an alternative way to interpolate a polynomial  $p \in \mathbb{K}(x)$  of degree  $n$  through  $n + 1$  points  $(x_i, y_i)$  for  $i = 1, \dots, n + 1$ .

1. determine  $n + 1$  auxiliary polynomials of degree  $n$

$$p_i(x) = \frac{\prod_{j \neq i} (x - x_j)}{\prod_{j \neq i} (x_i - x_j)} = \frac{(x - x_1) \dots (x - x_{i-1})(x - x_{i+1}) \dots (x - x_{n+1})}{(x_i - x_1) \dots (x_i - x_{i-1})(x_i - x_{i+1}) \dots (x_i - x_{n+1})}$$

Observe that  $p_i(x_i) = 1$  and  $p_i(x_j) = 0$  for  $i \neq j$ .

2. compose  $p(x)$  as the linear combination  $p(x) = \sum_{i=1}^{n+1} y_i p_i(x)$ .

Then  $p(x_i) = y_i p_i(x_i) = y_i$  as in all the other terms  $p_j(x_i) = 0$ .

## Example of Lagrange interpolation

Goal: interpolate a polynomial

$p(x) = a_4x^4 + a_3x^3 + a_2x^2 + a_1x + a_0$  over  $\mathbb{Z}_{11}$  through points  $(1, 5)$ ,  $(2, 1)$ ,  $(3, 3)$ ,  $(4, 4)$ ,  $(5, 3)$ ,  $(6, 5)$  and  $(7, 10)$ .

We seek  $a_4, a_3, a_2, a_1$  and  $a_0$ , that satisfy (over  $\mathbb{Z}_{11}$  !!)

$$\begin{array}{rcccccc} a_4 & + & a_3 & + & a_2 & + & a_1 & + & a_0 & = & 5 \\ 5a_4 & + & 8a_3 & + & 4a_2 & + & 2a_1 & + & a_0 & = & 1 \\ 4a_4 & + & 5a_3 & + & 9a_2 & + & 3a_1 & + & a_0 & = & 3 \\ 3a_4 & + & 9a_3 & + & 5a_2 & + & 4a_1 & + & a_0 & = & 4 \\ 9a_4 & + & 4a_3 & + & 3a_2 & + & 5a_1 & + & a_0 & = & 3 \\ 9a_4 & + & 7a_3 & + & 3a_2 & + & 6a_1 & + & a_0 & = & 5 \\ 3a_4 & + & 2a_3 & + & 5a_2 & + & 7a_1 & + & a_0 & = & 10 \end{array}$$

In fact, 5 points suffices. We may restrict ourselves to the first 5 equations (and the first 5 points).

We first calculate partial polynomials  $p_1, \dots, p_5$ .

These polynomials satisfy:  $p_i(x_i) = 1$  and also  $j \neq i : p_i(x_j) = 0$ .

$$p_1(x) = \frac{(x-2)(x-3)(x-4)(x-5)}{(1-2)(1-3)(1-4)(1-5)} = \frac{x^4+8x^3+5x^2+10}{2} = 6x^4 + 4x^3 + 8x^2 + 5$$

$$p_2(x) = \frac{(x-1)(x-3)(x-4)(x-5)}{(2-1)(2-3)(2-4)(2-5)} = \frac{x^4+9x^3+4x^2+3x+5}{5} = 9x^4 + 4x^3 + 3x^2 + 5x + 1$$

$$p_3(x) = \frac{(x-1)(x-2)(x-4)(x-5)}{(3-1)(3-2)(3-4)(3-5)} = \frac{x^4+10x^3+5x^2+10x+7}{4} = 3x^4 + 8x^3 + 4x^2 + 8x + 10$$

$$p_4(x) = \frac{(x-1)(x-2)(x-3)(x-5)}{(4-1)(4-2)(4-3)(4-5)} = \frac{x^4+8x^2+5x+8}{5} = 9x^4 + 6x^2 + x + 6$$

$$p_5(x) = \frac{(x-1)(x-2)(x-3)(x-4)}{(5-1)(5-2)(5-3)(5-4)} = \frac{x^4+x^3+2x^2+5x+2}{2} = 6x^4 + 6x^3 + x^2 + 8x + 1$$

The desired polynomial is combined from the partial polynomials and from the values in the given points  $(i, p(i))$  as:

$$\begin{aligned} p(x) &= \sum_{i=1}^5 y_i p_i(x) = 5p_1(x) + p_2(x) + 3p_3(x) + 4p_4(x) + 3p_5(x) \\ &= 3x^4 + 5x^2 + 2x + 6 \end{aligned}$$

We may check, whether the other points  $(6, 5)$ ,  $(7, 10)$  lie on  $p(x)$

$$p(6) = 3 \cdot 6^4 + 5 \cdot 6^2 + 2 \cdot 6 + 6 = 3 \cdot 9 + 5 \cdot 3 + 2 \cdot 6 + 6 = 5$$

$$p(7) = 3 \cdot 7^4 + 5 \cdot 7^2 + 2 \cdot 7 + 6 = 3 \cdot 3 + 5 \cdot 5 + 2 \cdot 7 + 6 = 10$$

## Applications

**Problem:** Given numbers  $m$  and  $n$ , design  $m$  keys so that:

- ▶ It is possible to reconstruct a given secret from any combination of  $n$  keys, but
- ▶ it is impossible to reconstruct a given secret from any combination of less than  $n$  keys.

Assume that the way the keys are constructed is publicly known.

**Solution:** Construct a polynomial of degree  $n - 1$  and distribute  $m$  distinct pairs  $(x_i, p(x_i))$  as keys. The secret is the polynomial.

The field could be e.g.  $\mathbb{R}$  or  $\mathbb{Z}_p$  with  $p > m$ .

---

**Problem:** Can two integers of  $n$  digits be multiplied in  $o(n^2)$  time?

**Solution:**

- ▶ Interpret these integers as polynomials  $p, q$  of degree  $n - 1$ ,
- ▶ choose  $2n$  pairs  $(i, p(i)), (i, q(i))$  and compute  $(i, p(i)q(i))$ ,
- ▶ then find the coefficients of the product  $pq$  in time  $O(n \log n)$ .

The choice of a suitable field and the recurrence behind is the principle of the so called *fast Fourier transform*.