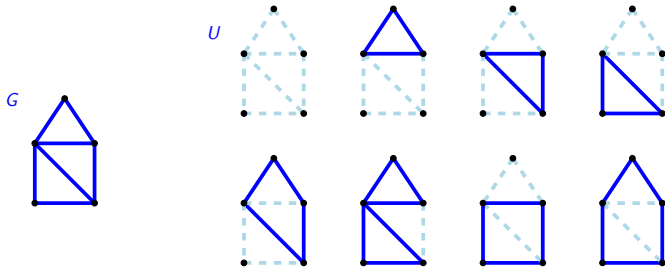


Even subgraphs

Let G be a connected graph and U contain edge sets A such that every vertex of G belongs to an even number of edges in A . Such sets A yield the so called *even subgraphs* of G .

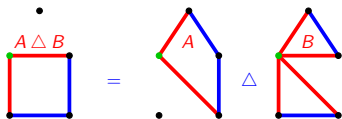


Problem: How many even subgraphs does G contain?

Vector space of even subgraphs

The symmetric difference \triangle preserves even degrees, because the symmetric difference of two sets of **even** cardinality, namely the **edges** incident to **a vertex**, has also an **even** cardinality.

$$|A \triangle A| = |A| + |A| - 2|A \cap A|$$



Hence $(\mathcal{A}, \triangle, \cdot)$ forms a vector space over \mathbb{Z}_2 .

Vector spaces of finite cardinality satisfy $|\mathcal{A}| = |\mathbb{F}|^{\dim(\mathcal{A})}$.

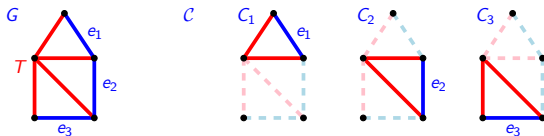
Equivalent problem: Construct a basis of \mathcal{A} .

Construction of a basis \mathcal{C}

For any spanning tree T of graph G , let $\{e_1, \dots, e_k\} = E_G \setminus E_T$.

For each $i \in \{1, \dots, k\}$ let C_i be the unique cycle in $T \cup e_i$.

The set $\mathcal{C} = \{C_1, \dots, C_k\}$ is *linearly independent*, since for every i , the edge e_i is in C_i , but cannot be eliminated by the symmetric difference of C_i with other cycles from \mathcal{C} , as they do not have e_i .



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For any even subgraph A , let $\{e_{i_1}, \dots, e_{i_l}\} = A \setminus E_T$, i.e., $i_1, \dots, i_l \in \{1, \dots, k\}$ are indices of edges from A but not in T .

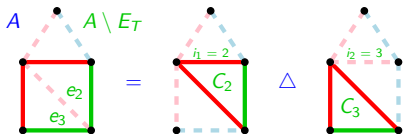
The graph $A \triangle C_{i_1} \triangle C_{i_2} \triangle \dots \triangle C_{i_l}$ is an even subgraph of G , but it is also a subgraph of T , since it has no edges from $E_G \setminus E_T$.

The tree has no cycles, so $A \triangle C_{i_1} \triangle C_{i_2} \triangle \dots \triangle C_{i_l} = \mathbf{0}$.

From the equality follows

$$A = C_{i_1} \triangle C_{i_2} \triangle \dots \triangle C_{i_l}.$$

Therefore, \mathcal{C} *generates* \mathcal{A} .



Answer: Each connected graph G has $2^{|E_G| - |V_G| + 1}$ even subgraphs.