## Small Complete Minors Above the Extremal Edge Density by Asaf Shapira and Benny Sudakov

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**Definition.** G graph, subgraph H of G is a *minor* of G if it can be obtained from G by a sequence of edge and vertex deletions and contractions.

 $K_t$ -minor of a graph G: connected subgraphs  $S_1, \ldots, S_t$  of G, such that there exist internally disjoint paths  $P_{i,j}$  joining each  $S_i, S_j$ .

**Definition.**  $c(t) = min\{c : d(t) \ge c \implies Ghas \ a \ K_t minor\}$ 

THE RESULT:

**Theorem. 1.** For all  $\epsilon > 0$ , integer t,  $\exists n_0(\epsilon, t)$  such that each graph G fulfilling  $|V(G)| \ge n_0$  and  $d(G) \ge c(t) + \epsilon$  has a  $K_t$ -minor of order at most  $O(\frac{c(t)t^2}{\epsilon} \log(n) \log\log(n))$ .

WE WILL PROVE:

**Lemma. 3.1.**  $\forall \epsilon > 0, t \ge 3, n \ge n_1(\epsilon, t)$ each graph *G* of order *n* and with  $d(G) \ge c(t) = \epsilon$  has a  $K_t$ -minor of order at most  $O(\frac{c(t)t^2}{\epsilon}log(n)(loglog(n))^3)$ .

**Definition.**  $\delta$ -expander Graph H, |V(H)| = m is a  $\delta$ -expander if  $\forall d : 0 \leq d \leq loglog(m) - 1$ ,  $\forall S : S \subseteq V(H), |S| \leq \frac{m}{2^{2^d}} \Longrightarrow$ 

$$|N(S)| \ge \frac{\delta 2^d}{\log(m) \log\log(m)^2} |S|.$$

**Lemma. KEY LEMMA 1.2.** Graph G, d(G) = c, then  $\forall \delta$ :  $0 < \delta < \frac{1}{256} \exists H$  subgraph of G such that:  $d(H) \ge (1 - \delta)c$  and H is a  $\delta$ -expander.

**Claim. 2.1.** Graph G on n vertices, d(G) = c, subset  $S \subset V(G)$  such that  $|N(S)| < \gamma |S| \implies$  one of the following is fulfilled:

- 1.  $d(G[V \setminus S])| \ge c$
- 2.  $d(G[S \cup N(S)])| \ge (1 \gamma)c.$

Proof. (KEY LEMMA)

- sequence of graphs:  $G = G_0, \ldots$
- end when  $|G_t| \leq 256$  or  $G_t$  is a  $\gamma$ -expander
- use claim 2.1. and define  $G_{t+1}$ :
- case1  $G_{t+1} = G[V(G_t) \setminus S_t]$
- case2  $G_{t+1} = G[S_t \cup N(S_t)]$
- find the maximal edge loss when t passing through intervals  $[2^{2^{k-1}}, 2^{2^k}]$

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**Lemma. 3.2.**  $\forall \delta > 0, t, m \ge m_0(\delta, t)$ each graph H is a  $\delta$ -expander of order m then H has a  $K_t$ -minor of order at most  $O(\frac{t^2}{\delta}log(n)(loglog(n))^3)$ . **Definition.** Ball with radius k and center v is a set  $B_k(v)$  of all vertices whose distance from v is  $\leq k$ .

**Claim. 3.3.** Let U, V be subsets of V(G) fulfilling:  $\forall d: 0 \leq d \leq loglog(m) - 1$ , whenever  $|B_k(U)| \leq \frac{m}{22^d}$  then

$$|N(B_{k+1}(U))| \ge \frac{\delta 2^d}{10\log(m)(\log\log(m))^2} |B_k(U)|,$$

and the same holds for the set V.

Then there exists a path from U to V of length at most  $\frac{20}{\delta} log(m)(loglog(m))^3$ .

Proof. (claim 3.3)

• sufficient:  $|B_k(U)| > \frac{m}{2}$  for some  $k \ge \frac{20}{\delta} log(m) (loglog(m))^3$ 

**Definition.**  $\delta$ -expanding ball Ball  $B_k(v)$  is  $\delta$ -expanding if  $\forall i : 1 \le i \le k-1$ :  $|B_{i+1}(v)| \ge |B_i(v)|(1+\gamma)$ .

**Claim. 3.5.**  $\forall \delta > 0, t, m \ge m_0(\delta, t)$  each  $\delta$ -expander G of order m fulfills one of the following:

- 1. G has t vertices of order > t
- 2. *G* contains *t* disjoint sets  $S_1, \ldots S_t$  such that  $G[S_i]$  is  $\gamma$ -expanding, and  $m^{\frac{1}{5}} \leq |S_i| \leq m^{\frac{1}{4}} \, \forall i$  and  $\forall v \in G[S_i] \deg(v) \leq \log^4(m)$ , where  $\gamma = \frac{\delta}{5(\log\log(m))^2}$ .

Proof. (claim 3.5.)

- if not ad 1. then  $T = \{v, deg(v) > log^4(m)\} < t$
- "nice" set S: G[S] is a  $\gamma\text{-expanding ball and }\frac{1}{5} < |S| < \frac{1}{4}$
- each  $W \supseteq T$ ,  $|W| \le m^{\frac{1}{3}}$  contains a "nice" subset one can iteratively pick  $S_i$  from  $G \setminus W_i$ , where  $W_i = (\bigcup_{j \le i} S_j) \cup$
- proof goes on by contradiction: suppose there is some W of the given properties, that does not contain a "nice" subset
- set  $G_0 = G \setminus W$ , take  $k_i$  the smallest that violates the expander condition for some fixed  $v_1$
- set  $T_1 = B_{k_1}(v_1)$ , then  $G_1 = G_0 \setminus T_1, \ldots$  until the first time  $\bigcup_i T_i > \frac{\sqrt{(m)}}{2}$
- yield a contradiction

**Claim. 3.6.**  $\forall \delta > 0, t, m \ge m_0(\delta, t)$  and G a  $\delta$ -expander of order m, which has t vertices fulfilling ad 1. from Claim 3.5.

 $\implies$  G has a  $K_t$ -minor of order at most  $O(\frac{t^2}{\delta}log(n)(loglog(n))^3)$ .

Claim. 3.6.  $\forall \delta > 0, t, m \ge m_0(\delta, t)$  and G a  $\delta$ -expander of order m, which has t subsets fulfilling ad 2. from Claim 3.5.  $\implies G$  has a  $K_t$ -minor of order at most  $O(\frac{t^2}{\delta}log(n)(loglog(n))^3)$ .

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