3-coloring planar graphs with four triangles

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Definitions (4-critical graphs)

graph G = (V, E)coloring is $\varphi : V \to C$ such that $\varphi(u) \neq \varphi(v)$ if $uv \in E$ G is a *k*-colorable if coloring with |C| = k exists G is a 4-critical graph if G is not 3-colorable but every $H \subset G$ is 3-colorable.



Theorem (Grötzsch '59)

Every planar triangle-free graph is 3-colorable.

More triangles?

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Theorem (Grünbaum '63; Aksenov '74; Borodin '97; Borodin et. al. '12+)

Let G be a planar graph containing at most three triangles. Then G is 3-colorable.



Question: What about four triangles?

Call 4-critical planar graph with four triangles a 4, 4-graph.

3-coloring planar graphs with four triangles?

Havel '69 found



Problem (Sachs '72)

Let G be a 4,4-graph. Can the triangles be partitioned into two pairs so that in each pair the distance between the triangles is less than two?

3-coloring planar graphs with four triangles? Havel '69 found



Problem (Sachs '72)

Let G be a 4, 4-graph. Can the triangles be partitioned into two pairs so that in each pair the distance between the triangles is less than two?

No! Aksenov and Mel'nikov '78,'80, Two infinite series of 4, 4-graphs.



Problem (Erdős '90) Describe 4, 4-graphs.

Borodin '97 - at least 15 infinite families of 4, 4-graphs (all with 4-faces)

Thomas and Walls '04 Infinite family \mathcal{TW} of 4, 4-graphs with no 4-faces. ...

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Some known 4, 4-critical planar graphs



Grötzsch's Theorem proof sketch as inspiration

Theorem (Grötzsch '59)

Every planar triangle-free graph G is 3-colorable.

Proof. CASE1: *G* has no 4-faces CASE2: *G* has a 4-face

OUR PLAN:

- characterize 4, 4-graphs without 4-faces
- describe how 4-faces could look like

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Results

Theorem

All plane 4, 4-graphs with no 4-faces can be obtained from the Thomas-Walls sequence



by replacing dashed edges by edges or



Call them C.

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All plane 4, 4-graphs with no 4-faces are precisely graphs in C.

Theorem

Every 4, 4-graph can be obtained from $G \in C$ by expanding some vertices of degree 3.



Act 1:

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All plane 4, 4-graphs with no 4-faces are precisely graphs in $\mathcal{C}=\{$



Theorem (Kostochka and Yancey; 12+) Let G be a 4-critical graph. Then $3e \ge 5v - 2$. Moreover, 3e = 5v - 2 iff G is 4-Ore.



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G is 4-Ore if $G = K_4$ or *G* is an Ore composition of two 4-Ore graphs.



Not 3-colorable.

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Observation

Let G be a plane 4,4-graph with no 4-faces. Then $3e \le 5v - 2$.

Every 4, 4-graph must be 4-Ore.

Key properties

G is 4, 4-*Ore* if it is 4-Ore and has 4 triangles. (4, 4-graph $\subseteq 4, 4$ -Ore)

• 4,4-Ore is K₄ or Ore composition of two 4-Ore graphs



Key properties

G is 4, 4-*Ore* if it is 4-Ore and has 4 triangles. (4, 4-graph $\subseteq 4, 4$ -Ore)

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Lemma G_a and G_b are 4, 4-Ore.





















And now few more...



















Lemma Every 4, 4-Ore graph is planar. Hence 4, 4-graphs = 4, 4-Ore graphs.

Description of $\ensuremath{\mathcal{C}}$

Theorem

All 4-critical plane graphs with four triangles and no 4-faces can be obtained from the Thomas-Walls sequence



by replacing dashed edges by edges or



Call them C.

Act 2:

Theorem Every 4, 4-graph can be obtained from $G \in C$ by expanding some vertices of degree 3.



(Interior of a 6-cycle is a quadrangulation - only 4-faces)





G - x is 3-colorable since G is 4-critical.



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Theorem (Gimbel and Thomassen '97)

"Quadrangulation is the only possible filling of a 6-cycle in a plane triangle-free 4-critical graph."

Conclusion

Problem (Sachs '72)

Let G be a 4,4-graph. Can the triangles be partitioned into two pairs so that in each pair the distance between the triangles is less than two?

Theorem

All plane 4, 4-graphs with no 4-faces are precisely graphs in C.

Theorem

Every 4, 4-graph with four triangles can be obtained from $G \in C$ by expanding some vertices of degree 3.

Corollary

Let G be a 4, 4-graph. The triangles can be partitioned into two pairs so that in each pair the distance between the triangles is at most two.

Thank you for your attention!

