

# Packing chromatic number for square and hexagonal lattices

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## Packing Chromatic Number

### Definition

Graph  $G = (V, E)$ ,  $X_d \subseteq V$  is  $d$ -packing if  
 $\forall u, v \in X_d : \text{distance}(u, v) > d$ .

1-packing is an independent set

### Definition

*Packing chromatic number* is the minimum  $k$  such that  
 $V = X_1 \cup X_2 \cup \dots \cup X_k$ ; denoted by  $\chi_\rho(G)$ .

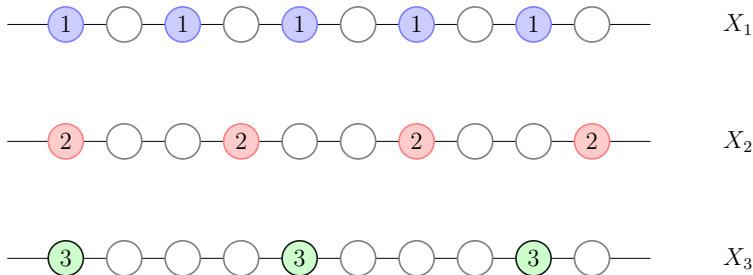
Also known as the *broadcast chromatic number*.



## Example with path $P_\infty$

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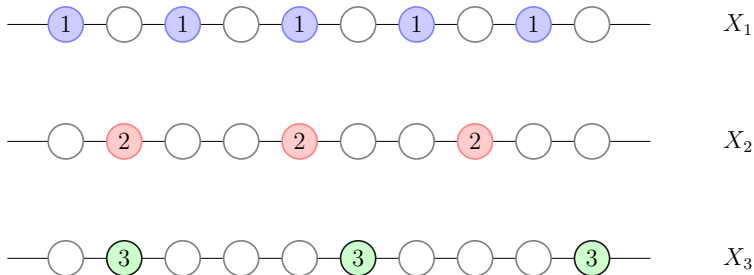
informally, density of  $X_d$  is  $|X_d|/|V|$



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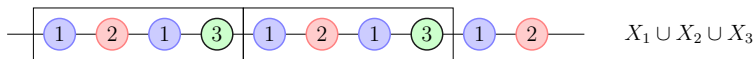
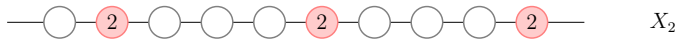
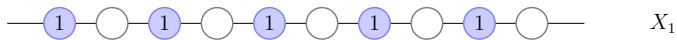


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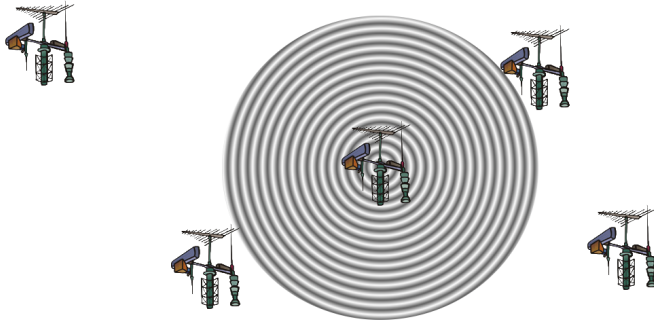
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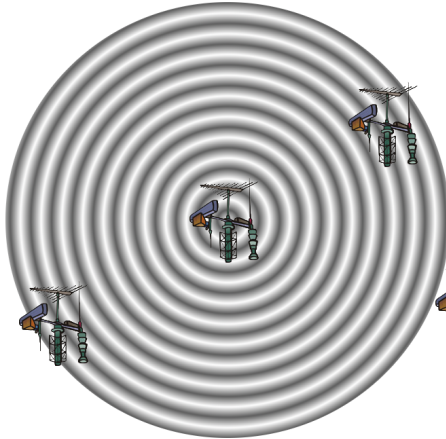
$$\chi_\rho(P_\infty) = 3$$



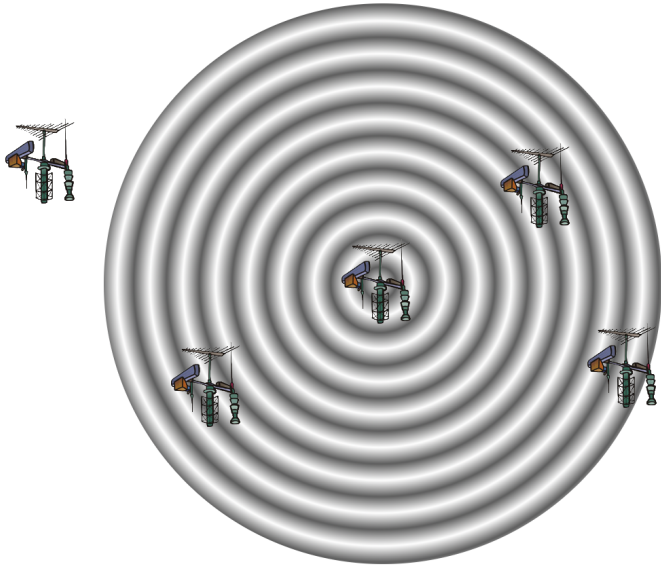
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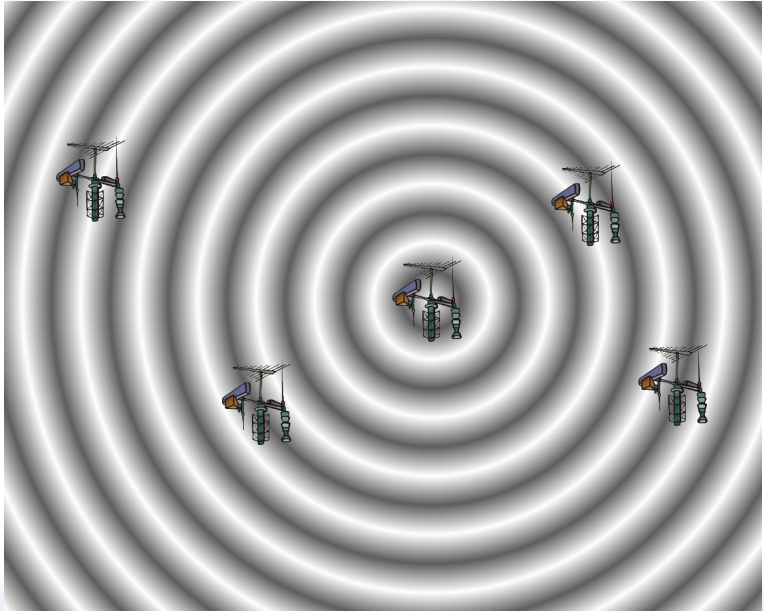


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## Complexity of $\chi_\rho$

### Theorem (Goddard et al. '08)

Let  $G$  be a graph.

- Decide if  $\chi_\rho(G) \leq k$  is  $\mathcal{NP}$ -complete ( $k$  on input).
- Decide if  $\chi_\rho(G) \leq 3$  is in  $\mathcal{P}$ .
- Decide if  $\chi_\rho(G) \leq 4$  is  $\mathcal{NP}$ -complete.

### Theorem (Fiala and Golovach '09)

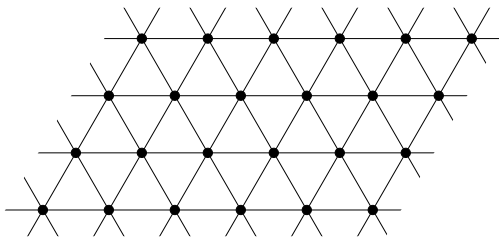
Decide if  $\chi_\rho(G) \leq k$  for trees is  $\mathcal{NP}$ -complete ( $k$  on input).



## Triangular lattice $\mathcal{T}$

### Theorem (Finbow and Rall '07)

*Infinite triangular lattice  $\mathcal{T}$  cannot be colored by a finite number of colors.*



We use notation  $\chi_\rho(\mathcal{T}) = \infty$ .



## Hexagonal Lattice

Theorem (Brešar, Klavžar and Rall '07)

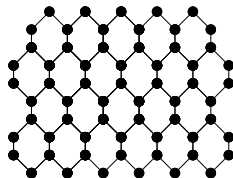
For hexagonal lattice  $\mathcal{H}$ :  $6 \leq \chi_\rho(\mathcal{H}) \leq 8$

Theorem (Vesel '07)

$7 \leq \chi_\rho(\mathcal{H})$

Theorem (Fiala, Klavžar, L. '09)

$\chi_\rho(\mathcal{H}) \leq 7$



## Square lattice

Theorem (Goddard et al. '02)

*For infinite planar square lattice  $\mathbb{Z}^2$ :*

$$9 \leq \chi_\rho(\mathbb{Z}^2) \leq 23$$

Theorem (Schwenk '02)

$$\chi_\rho(\mathbb{Z}^2) \leq 22$$

Theorem (Fiala, Klavžar, L. '09)

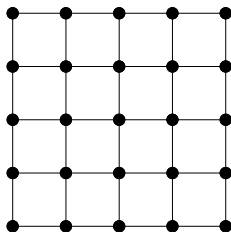
$$10 \leq \chi_\rho(\mathbb{Z}^2)$$

Theorem (Holub and Soukal '09)

$$\chi_\rho(\mathbb{Z}^2) \leq 17$$

Theorem

$$12 \leq \chi_\rho(\mathbb{Z}^2)$$



## How we did the lower bound 12

### Wish (Conjecture)

*If  $\chi_\rho(\mathbb{Z}^2) = k$  then exist  $X_1, \dots, X_k$  such that  $\forall i$   $X_i$  has maximum possible density after fixing  $\bigcup_{1 \leq j < i} X_j$ .*

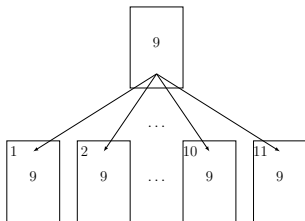
Wish implies the lower bound 12.

No wish implies brute force computer search (backtracking).

Find a (small) part of  $\mathbb{Z}^2$  that cannot be colored by 11 colors.



Grid  $15 \times 9$  cannot be colored using 11 colors when  $[5, 5]$  is precolored by 9.



- *easy* to compute in parallel
- two implementations
- 115 days of CPU time (in 2009)
- 43112312093324 configurations tested



## Layers of a lattice

Theorem (Finbow and Rall '07)

$$\chi_\rho(\mathbb{Z}^3) = \infty$$

Theorem (Fiala, Klavžar, L. '09)

$$\chi_\rho(P_2 \square \mathbb{Z}^2) = \infty$$

Theorem (Fiala, Klavžar, L. '09)

$$\chi_\rho(P_6 \square \mathcal{H}) = \infty$$

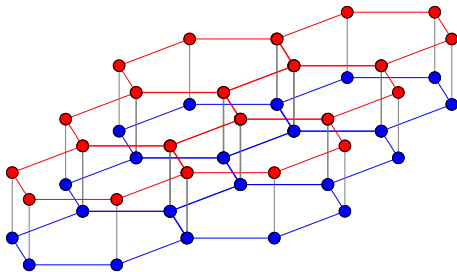
Theorem

$$\chi_\rho(P_2 \square \mathcal{H}) \leq 526 \text{ (large but finite)}$$

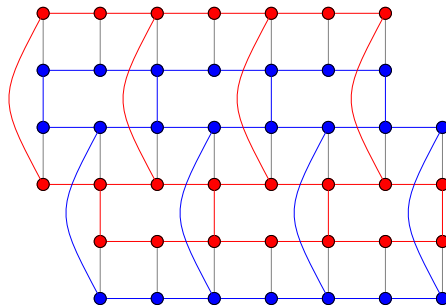
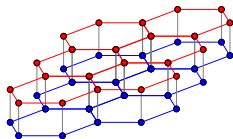




$$P_2 \square \mathcal{H}$$



$$P_2 \square \mathcal{H}$$



$$P_2 \square \mathcal{H}$$

## Observations

- large colors have large period
- use small colors as much as possible (the Wish)
- there might be a lot of locally good patterns
- do not try to fill large area of zeros

## Ideas

- glue small patterns with small colors to large pattern for large colors
- use some random while constructing pattern
- repeat search lot of times with crossed fingers



$$P_2 \square \mathcal{H}$$

## Result

- pattern ( $2 \times$ )  $768 \times 768$
- extra program for checking correctness of the result
- 526 colors (probably not optimal)
- lower bound for  $\chi_\rho(P_2 \square \mathcal{H})$  around 20  
(more that  $\chi_\rho(\mathbb{Z}^2)$ , huge gap)



## Open problems

- Is  $\chi_\rho(\mathcal{H} \square P_3)$  finite?
- What is  $\chi_\rho(\mathbb{Z}^2)$  for the infinite planar square lattice  $\mathbb{Z}^2$ ?
- Is there  $c$  such that every cubic graph  $G$  has  $\chi_\rho(G) \leq c$ ?
  - if  $G$  is planar?
  - if  $G$  has large girth?



## Open problems

- Is there  $c$  such that every cubic graph  $G$  has  $\chi_\rho(G) \leq c$ ?
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Theorem (Sloper '02)

*3-regular infinite tree  $T_3$ :  $\chi_\rho(T_3) = 7$*

Theorem (Sloper '02)

*4-regular infinite tree  $T_4$ :  $\chi_\rho(T_4) = \infty$*



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