#### Algorithmic game theory

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- In DSIC mechanisms, maximizing the expected revenue is then the same as maximizing the expected virtual social surplus
   ∑<sup>n</sup><sub>i=1</sub> φ<sub>i</sub>(v<sub>i</sub>)x<sub>i</sub>(v). (Theorem 3.14) where

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- If  $F_1 = \cdots = F_n = F$  is regular, then Vickrey's auction with reserve  $\varphi^{-1}(0)$  maximizes the expected revenue among all single-item auctions.
- What if the seller does not know the distributions  $F_1, \ldots, F_n$ ?

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#### Figure: Jeremy Bulow and Paul Klemperer.

Sources: https://economics.stanford.edu/ and https://www.economics.ox.ac.uk/

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• More competition is better than finding the right auction format.

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Sources: Auctions & Combinatorial auctions (Vincent Conitzer)

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#### Figure: William Vickrey, Edward H. Clarke, and Theodore Groves.

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• We can still do DSIC social surplus maximization.

• The government wants to construct roads connecting diverse cities, and he wants cities to pay for the roads.

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Sources: https://www.science4all.org/article/auction-design/

	Road Network 1	Road Network 2	Road Network 3	Road Network 4
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enteri	5 M\$	8 M\$	4 M\$	12 M\$
¢	2 M\$	1 M\$	20 M\$	4 M\$
<b>W</b>	4 M\$	6 M\$	3 M\$	5 M\$
Lindi	1 M\$	1 M\$	6 M\$	2 M\$
*	1 M\$	2 M\$	2 M\$	3 M\$
Total (social welfare)	19 M\$	32 M\$	37 M\$	42 M\$

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• Cities pay their negative externalities on the collectivity. Other cities would be happier without the biggest city (NYC, say). How much happier they would be is exactly what NYC must pay.

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 If NYC was not there, then road network number 3 (RN3) would have been chosen, as opposed to RN4. The value of RN3 for the other cities would be 35 M\$, as opposed to the 26 M\$ of RN4. Therefore, the negative externality of NYC is 35 - 26 = 9 M\$.





Nash equilibria













Thank you for your attention.