# Algorithmic game theory 

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## 9th lecture

December 2nd 2021


Regret minimization

Example

## Example



| Weather | *** | $\text { F } \begin{aligned} & * \\ & * \\ & * \end{aligned}$ |  | *** | Profit |
| :---: | :---: | :---: | :---: | :---: | :---: |
| Algorithm |  |  | 前 | $J$ | 3 |
| Umbrella |  |  |  | $\sqrt{ }$ | 3 |
| Sunscreen |  |  |  |  | 1 |

Source: No regret algorithms in games (Georgios Piliouras)

Polynomial weights algorithm

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Algorithm 0.4: $\operatorname{Polynomial~weights~algorithm~}(X, T, \eta)$

```
Input: A set of actions }X={1,\ldots,N},T\in\mathbb{N}\mathrm{ , and }\eta\in(0,1/2]
Output: A probability distribution p}\mp@subsup{p}{}{t}\mathrm{ for every time step t.
wi
p
for }t=2,\ldots,
do {l { w wi
```

No-regret dynamics

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Algorithm 0.8: No-Regret dynamics $(G, T, \varepsilon)$
Input: A game $G=(P, A, C)$ of $n$ players, $T \in \mathbb{N}$ and $\varepsilon>0$. Output: A prob. distribution $p_{i}^{t}$ on $A_{i}$ for each $i \in P$ and step $t$. for every step $t=1, \ldots, T$
(Each player $i \in P$ independently chooses a mixed strategy
 Each player $i \in P$ receives a loss vector $\ell_{i}^{t}=\left(\ell_{i}^{t}\left(a_{i}\right)\right)_{a_{i} \in A_{i}}$, where $\ell_{i}^{t}\left(a_{i}\right) \leftarrow \mathbb{E}_{a_{-i}^{t} \sim \prod_{j \neq i} i_{j}^{t}}\left[C_{i}\left(a_{i} ; a_{-i}^{t}\right)\right]$.

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Source: https://img.etimg.com

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## Thank you for your attention.

