

# Triangles in random graphs

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## Abstract

We show the number of triangles of  $G_{n, \frac{1}{2}}$  is almost uniformly distributed among residue classes modulo  $q$ , where  $q$  is a prime number bounded by  $\Theta(\log n)$ . This implies a consequence of a conjecture of Bollobás, Pebody and Riordan (that almost every random graph  $G_{n, \frac{1}{2}}$  is uniquely determined by its Tutte polynomial): almost every pair of independently chosen random graphs  $G_{n, \frac{1}{2}}$  has different Tutte polynomials.

## 1 Introduction

The Tutte polynomial of a graph  $G = (V, E)$  is the bivariate polynomial

$$T(G; x, y) = \sum_{A \subseteq E} (x-1)^{r(E)-r(A)} (y-1)^{|A|-r(A)},$$

where  $r(A) = |V| - k(A)$  and  $k(A)$  denotes the number of components of the graph  $(V, A)$ .

It is clear that two isomorphic graphs have the same Tutte polynomial. On the other hand, there are examples of non-isomorphic graphs with the same Tutte polynomial (see [3, 1]). In this context, the question of how much information about the graph its Tutte polynomial contains is natural.

A random graph  $G_{n, \frac{1}{2}}$ , for  $n$  natural number, is a graph on an  $n$  element vertex set such that any pair of vertices forms an edge with probability  $\frac{1}{2}$  independently of other pairs. Let  $\mathcal{G}_{n, \frac{1}{2}}$  denote the corresponding probability

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space. Note that such a random graph dense; the expectation of the number of edges is  $\frac{1}{2}\binom{n}{2}$  and the expectation of the number of triangles is  $\frac{1}{8}\binom{n}{3}$ . The following conjecture was presented in [1].

**Conjecture 1.1** *Almost every graph  $G \in \mathcal{G}_{n, \frac{1}{2}}$  is such that  $T(G; x, y) = T(G'; x, y)$  implies  $G \cong G'$ .*

The main result of this paper stated below implies a weaker result.

**Theorem 1.2** *There are numbers  $q_0$  and  $n_0$  and a function  $f \in \Theta(\log n)$  such that for any  $n > n_0$ , any prime number  $q$ ,  $q_0 < q < f(n)$ , and any choice of  $k \in \{0, 1, \dots, q - 1\}$  the probability that a random graph  $G_{n, \frac{1}{2}}$  contains  $k$  triangles modulo  $q$  is  $q^{-1}(1 + o(1))$ .*

Since the number of triangles of any simple graph can be easily derived from its Tutte polynomial, we can state this easy consequence of Theorem 1.2.

**Corollary 1.3** *The probability that two independently chosen random graphs  $G_{n, \frac{1}{2}}$  have the same Tutte polynomial is of order  $O(1/\log n)$ .*

The another open problem is to find or to prove an existence of a family containing almost every graph such that the Tutte polynomial distinguish any pair of graphs inside this family. Solving this problem affirmatively may give us better understanding and a hope for positive answer to the Conjecture 1.1.

## 2 Counting 2-paths

Let  $G$  be a random graph from  $\mathcal{G}_{n, \frac{1}{2}}$  with the vertex set  $V$ . Let  $V_1$  be a set of  $\lfloor \frac{n}{4} \rfloor$  pairs of vertices  $V_1 = \{u_1, v_1, u_2, v_2, \dots, u_{\lfloor \frac{n}{4} \rfloor}, v_{\lfloor \frac{n}{4} \rfloor}\}$  (suppose  $n$  large enough) and let  $V_2$  be a subset of  $V - V_1$ . Denote the size of  $V_2$  by  $l$ . Fix a pair  $u_i, v_i$  of vertices of  $V_1$ . For  $w \in V_2$ , let  $X_w$  be an indicator of the existence of both edges  $\{u_i, w\}$  and  $\{v_i, w\}$ . Then  $X = \sum_{w \in V_2} X_w$  counts the number of 2-paths connecting vertices  $\{u, v\}$  and having the middle vertex in  $V_2$ . Note that the probability of such a path through a fixed vertex  $w$  is  $(\frac{1}{2})^2$ . Denote by  $p_j(l)$  the probability that  $X = j$ . Since edges from distinct vertices of  $V_1$  going to  $V_2$  are independent, the number

of such 2-paths is independent for distinct pairs of vertices of  $V_1$ , and  $p_j(l)$  is also independent of the exact choice of  $V_2$ .

Assume  $q$  to be an arbitrary fixed integer greater than 1 and smaller than  $n$ . Denote by  $s_i(l) = \sum_{j \equiv i \pmod{q}} p_j(l)$  the probability that the number of connecting 2-path is equal to  $i$  modulo  $q$ .

**Lemma 2.1** *Let  $1 \leq q < n$ . There exists a number  $c = c(q) > 0$  such that  $s_i(l) \geq c$  for  $i = 0, 1, \dots, q-1$  and  $l \geq q-1$ .*

**Proof:** Using the recursion  $p_i(l) = \frac{1}{4}p_{i-1}(l-1) + \frac{3}{4}p_i(l-1)$  one can derive  $s_i(l) = \frac{1}{4}s_{i-1}(l-1) + \frac{3}{4}s_i(l-1)$  (where  $s_{-1}(l)$  means  $s_{q-1}(l)$ ).

It is easy to observe that if  $s_i(l_0) \geq c$  for all possible  $i$ 's and some  $l_0$  and  $c$ , then also  $s_i(l) \geq c$  for all  $l \geq l_0$ . Note that  $s_i(q-1) = p_i(q-1)$ , and clearly, each  $p_i(q-1) \geq (\frac{1}{4})^{q-1}$ . Hence for any  $l \geq q-1$  and any  $i = 1, 2, \dots, q-1$  we have  $s_i(l) \geq (\frac{1}{4})^{q-1}$ . By setting  $c(q) = (\frac{1}{4})^{q-1}$  we complete the proof. □

**Lemma 2.2** *Let  $n$  be sufficiently large natural number. Then with probability at least  $1 - \frac{1}{15} \cdot \frac{\log_2 n}{n^{1/3}}$  a graph  $G \in \mathcal{G}_{n, \frac{1}{2}}$  has, for any  $q \leq q_0 = \lfloor \log_2 n^{1/3} \rfloor$ , disjoint pairs  $u_k, v_k$  of vertices of  $G$  ( $k = 1, 2, \dots, q-1$ ) where the number of 2-paths connecting  $u_k$  and  $v_k$  is  $k$  modulo  $q$ .*

**Proof:** Let  $V_1$  be the set of  $\lfloor \frac{n}{4} \rfloor$  pair of vertices as described above and let  $V_2 = V - V_1$ . We want to chose pairs  $u_k, v_k$  from the pairs of vertices forming  $V_1$ . By Lemma 2.1 the probability that the number of 2-paths connecting  $u_i$  and  $v_i$  is equal to  $k$  modulo  $q$  is at least  $c(q)$  for any  $k = 0, 1, \dots, q-1$  (and hence at most  $1 - (q-1)c(q)$ ). Moreover, these probabilities are independent for distinct pairs  $u_i, v_i$  and  $u_j, v_j$ .

Denote the event that there is a pair  $u_i, v_i$  (among the pairs from  $V_1$ ) connected by exactly  $k$  2-paths (counted modulo  $q$ ) by  $Y_k$  ( $Y_k = 1$  if such a pair exists,  $Y_k = 0$  otherwise). Then the probability  $P[Y_1, Y_2, \dots, Y_{q-1}]$  is equal to  $1 - P[\exists k : \overline{Y_k}]$ . Assume the edges inside  $V_1$  are fixed. The probability that vertices  $u_i, v_i$  are connected by  $k$  2-paths is  $s_{i'}(|V_2|)$ , where  $i'$  is such that the number of 2-paths with the middle vertex inside  $V_1$  plus  $i'$  is equal to  $k$  modulo  $q$ . Hence  $P[\overline{Y_k}] = \prod_{i=1}^{\lfloor \frac{n}{4} \rfloor} (1 - s_{i'}(|V_2|)) \leq (1 - c(q))^{\lfloor \frac{n}{4} \rfloor}$ .

$$P[\exists k : \overline{Y_k}] \leq \sum_{k=1}^{q-1} P[\overline{Y_k}] \leq \sum_{k=1}^{q-1} (1 - c(q))^{\lfloor \frac{n}{4} \rfloor} = (q-1)(1 - c(q))^{\lfloor \frac{n}{4} \rfloor}.$$

Let  $q_0 = \lfloor \log_2 n^{1/3} \rfloor$ . Since  $q \leq q_0$  and  $c(q) \geq c(q_0)$  we can bound

$$\begin{aligned} P[\exists k : \overline{Y}_k] &\leq q_0 (1 - c(q_0))^{\lfloor \frac{n}{4} \rfloor} = q_0 \frac{(1 - c(q_0)^2)^{\lfloor \frac{n}{4} \rfloor}}{(1 + c(q_0))^{\lfloor \frac{n}{4} \rfloor}} \leq \\ &\leq \frac{q_0}{(1 + c(q_0))^{\lfloor \frac{n}{4} \rfloor}} \leq \frac{q_0}{1 + \lfloor \frac{n}{4} \rfloor c(q_0)}. \end{aligned}$$

From Lemma 2.1 we have  $c(q_0) = (\frac{1}{4})^{q_0-1} \geq (\frac{1}{4})^{\log n^{1/3}} = 2^{-2 \log n^{1/3}} = n^{-2/3}$ . Hence

$$P[\exists k : \overline{Y}_k] \leq \frac{\log n^{1/3}}{1 + \lfloor \frac{n}{4} \rfloor n^{-2/3}} \leq \frac{\log n^{1/3}}{\frac{n}{5} n^{-2/3}} = \frac{1}{15} \frac{\log n}{n^{1/3}}.$$

□

### 3 Proof of the Theorem

Using the results of [2] we can state the following useful proposition.

**Proposition 3.1** *Let  $q$  be a sufficiently large prime number; then the number of subsets  $S \subseteq \{1, 2, \dots, q-1\}$  such that the sum  $\sum_{i \in S} i \equiv k \pmod{q}$  is  $2^{q-1} q^{-1} (1 + o(1))$  for any  $k = 0, 1, \dots, q-1$ .*

Combining the result of Lemma 2.2 with the previous proposition we can prove Theorem 1.2.

**Proof of Theorem 1.2:** Let us consider a graph  $G \in \mathcal{G}_{n, \frac{1}{2}}$ . It is possible to find disjoint pairs of vertices  $u_i, v_i$  such that the number of connecting 2-paths between  $u_i$  and  $v_i$  is equal to  $i$  modulo  $q$  for all  $i = 1, 2, \dots, q-1$  with probability at least  $1 - \frac{1}{15} \frac{\log n}{n^{1/3}}$ . Let us define graphs  $G_S$  for any  $S \subseteq \{1, 2, \dots, q-1\}$  obtained from  $G$  by adding all edges  $\{u_i, v_i\}$  with  $i \in S$  (if they are not in  $G$ ) and by deleting all edges  $\{u_i, v_i\}$  with  $i \notin S$  (if they are in  $G$ ). The probability of all  $G_S$  in the probability space  $\mathcal{G}_{n, \frac{1}{2}}$  is the same.

From now on, all computation will be done modulo  $q$ . Let  $t_0$  be the number of triangles of  $G_\emptyset$ . Since the pairs  $u_i, v_i$  are disjoint, edges  $\{u_i, v_i\}$  do not affect the number of 2-paths between these pairs. Then the number of triangles of  $G_S$  is  $t_0 + \sum_{i \in S} i$ . By Proposition 3.1, the probability that

the number of triangles of  $G_S$  is equal to  $k$  modulo  $q$  is  $q^{-1}(1 + o(1))$ , where  $S$  is chosen randomly so that each element is chosen independently with probability  $\frac{1}{2}$ .

Now, consider an arbitrary ordering of pairs of vertices and to each graph  $G \in \mathcal{G}_{n, \frac{1}{2}}$  assign (if possible) the pairs  $u_i, v_i$ ,  $i = 1, 2, \dots, q - 1$ , described above which are minimum in the fixed ordering. Note that graphs  $G_1$  and  $G_2$  are assigned the same pairs  $u_i, v_i$  and are identical on all other pairs if and only if they are  $G_{S_1}$  and  $G_{S_2}$  for some  $G$  and  $S_1, S_2 \subseteq \{1, 2, \dots, q - 1\}$ . Hence they have the same probability in  $\mathcal{G}_{n, \frac{1}{2}}$ .

Up to a  $\left(\frac{1}{15} \frac{\log n}{n^{1/3}}\right)$ -fraction, we have partitioned the graphs of  $\mathcal{G}_{n, \frac{1}{2}}$  into classes of equal size such that in each class the number of triangles computed modulo  $q$  is almost uniformly distributed. Since  $\frac{\log n}{n^{1/3}}$  tends to 0 for  $n$  going to infinity, we can conclude that the probability that  $G_{n, \frac{1}{2}}$  contains exactly  $k$  triangles modulo  $q$  is  $q^{-1}(1 + o(1))$  for any  $k = 0, 1, \dots, q - 1$ . □

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## References

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