

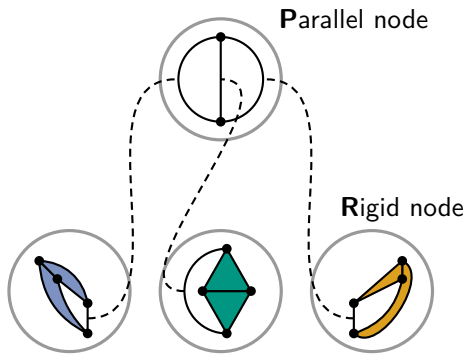
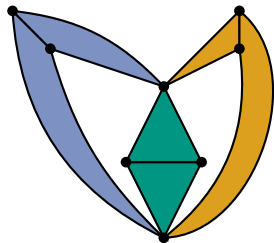
An SPQR-Tree-Like Embedding Representation for Upward Planarity

Guido Brückner¹, Markus Himmel¹, Ignaz Rutter²

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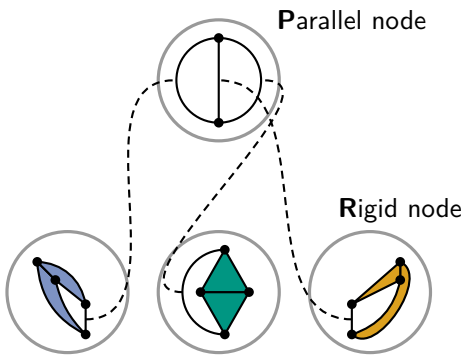
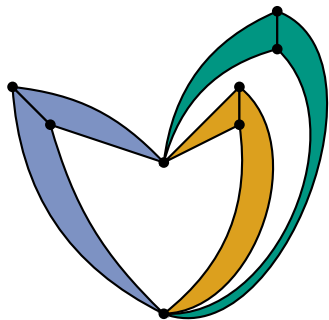


SPQR-Trees [Di Battista, Tamassia '90]



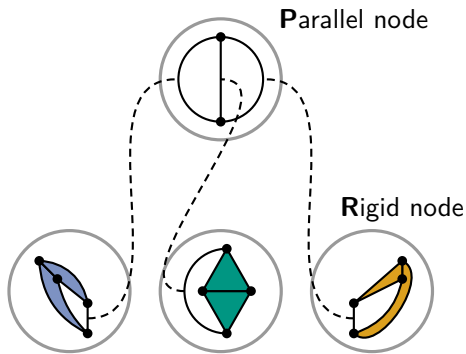
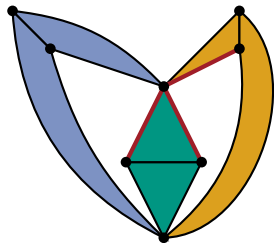
- Compactly represent all planar embeddings of a biconnected planar graph

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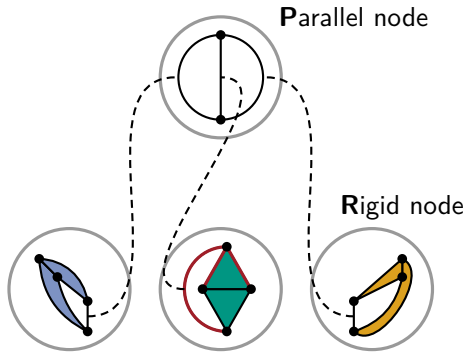
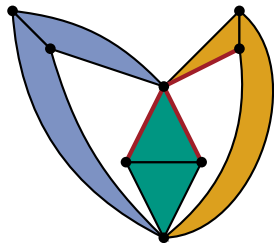
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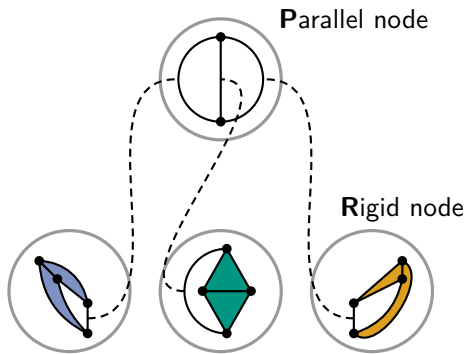
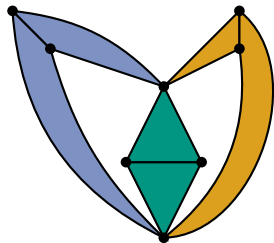
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- Have been used to efficiently solve a range of problems
 - Optimization [e.g. Didimo et al. '18], extension [Angelini et al. '15], ...

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- **Can we find a similar data structure for the upward planar case?**

Upward planarity testing

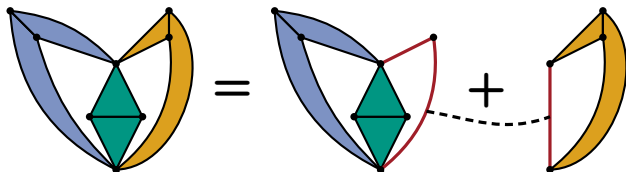
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- Sufficient to only consider biconnected graphs
- Basic idea: Decomposition at 2-vertex cuts
- “Shape” of the rest of the graph \longleftrightarrow suitable **marker graph**

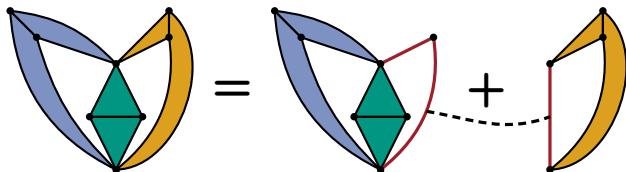


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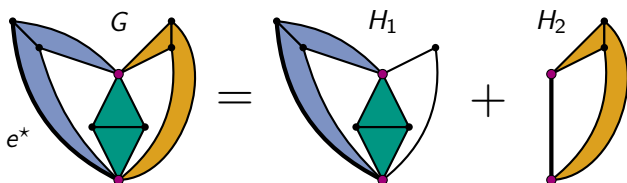
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- Linear-time algorithm by Bertolazzi et al. '98 based on SPQR-trees
- Simpler algorithm by Hutton and Lubiw '96 using general decompositions

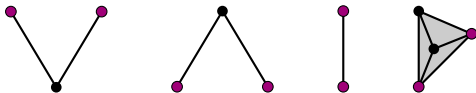
A decomposition result by Hutton and Lubiw

Lemma



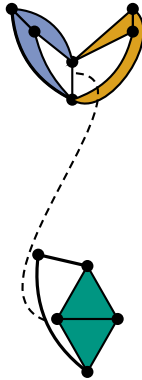
Bijjective correspondence between embeddings of G and combinations of embeddings of H_1 and H_2 where

- *Marker graphs determined by a set of rules*

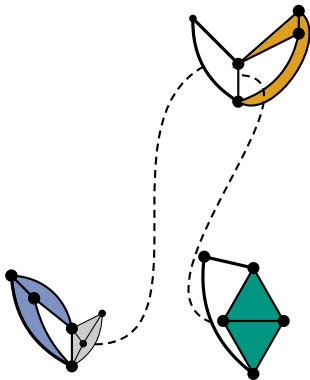


- H_1 or H_2 is single component
- Fixed edge e^* or its marker are leftmost

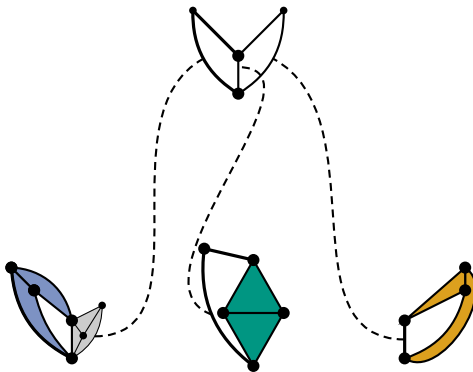
Decomposition Trees for Upward Planar Embeddings



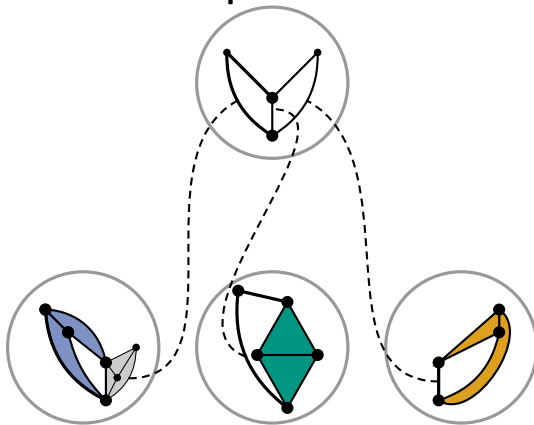
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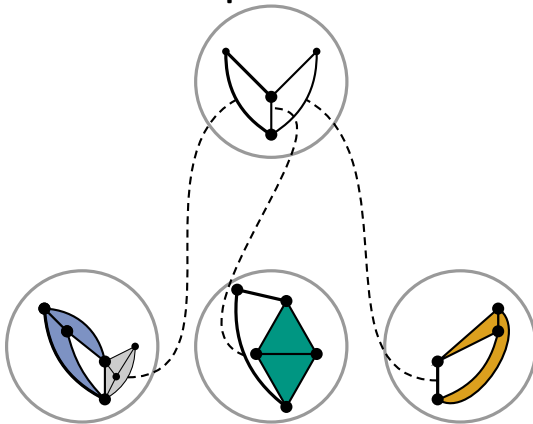


Decomposition Trees for Upward Planar Embeddings



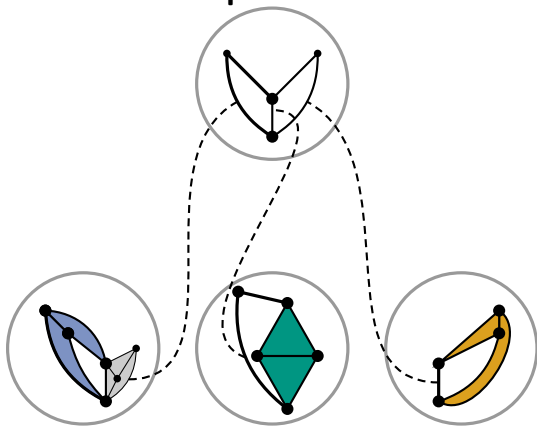
- Upward planar embedding \longleftrightarrow upward planar skeleton embeddings
- Each sequence of decompositions \rightsquigarrow new characterization of upward planar embeddings

Decomposition Trees for Upward Planar Embeddings



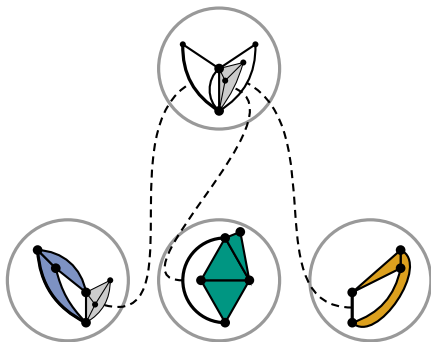
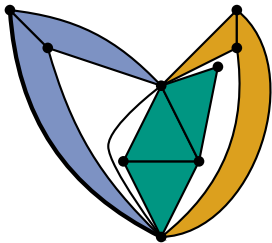
- Upward planar embedding \longleftrightarrow upward planar skeleton embeddings
- Each sequence of decompositions \rightsquigarrow new characterization of upward planar embeddings
 - Actually, order is irrelevant

Decomposition Trees for Upward Planar Embeddings

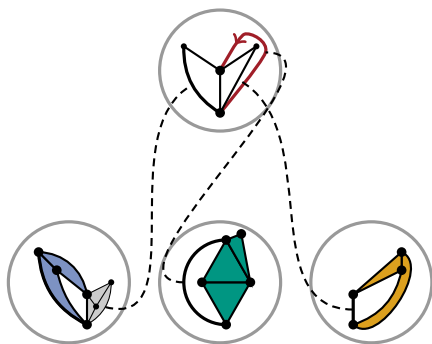
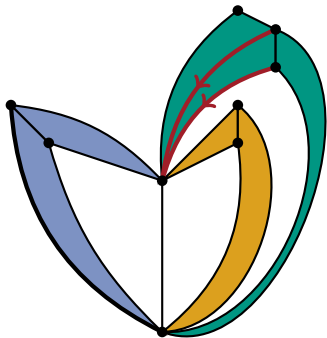


- Important question: Which decomposition tree should we use?
- SPQR-tree is nice for planar embeddings, but offers too many choices
- Idea: Modify SPQR-tree to have upward planar skeletons

P-Node Splits

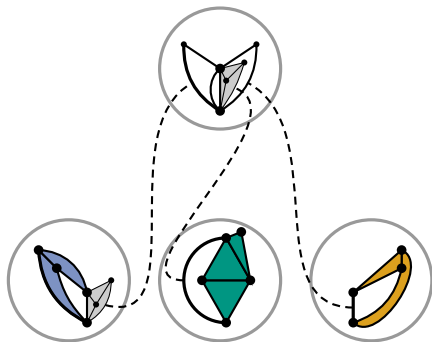
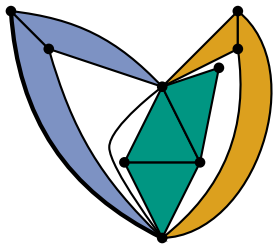


P-Node Splits



- Problem: Permuting edges at P-nodes \rightsquigarrow non-upward-planar skeleton

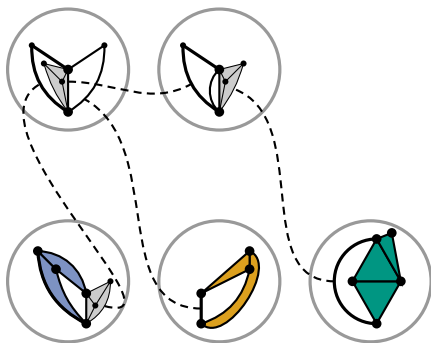
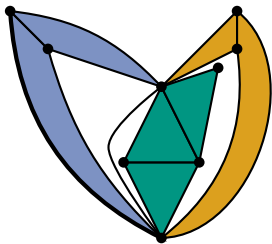
P-Node Splits



- Problem: Permuting edges at P-nodes \rightsquigarrow non-upward-planar skeleton
- Solution: Split P-nodes by marker type

■ Relevant here:  and .

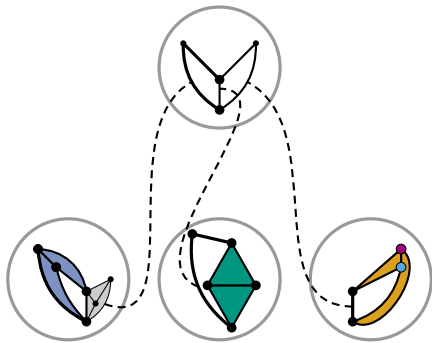
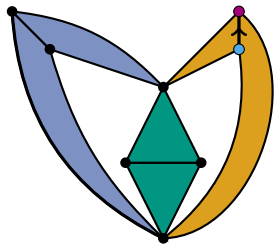
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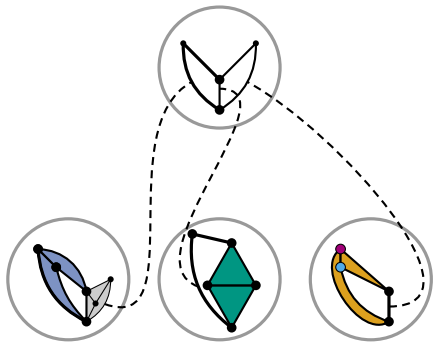
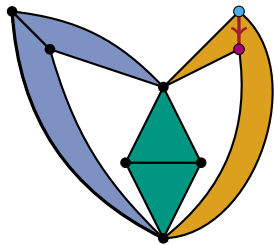
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Arc Contractions

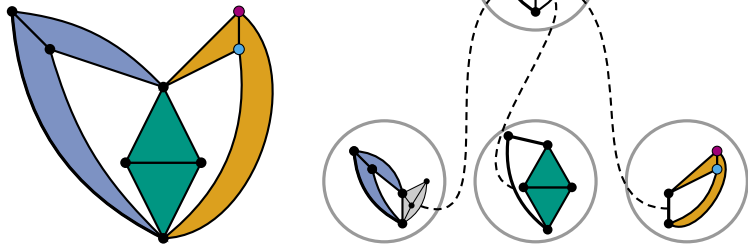


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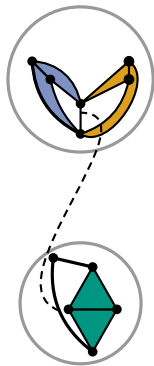
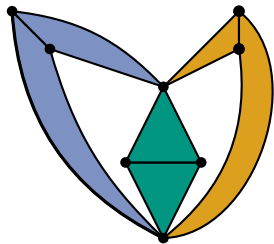
- Problem: Flipping operation at R-nodes \rightsquigarrow non-upward-planar skeleton

Arc Contractions



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 - Upward planarity test for embedded single-source graphs [Bertolazzi et al. '98]

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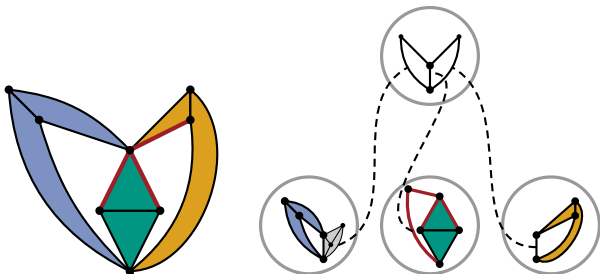
The UP-Tree

- SPQR-tree + P-node splits + arc contractions =: UP-tree

Theorem

For each biconnected single-source DAG G and e^* incident to s there is a decomposition tree \mathcal{T} computable in linear time that

- represents the upward planar embeddings of G in which e^* is leftmost
- does so using P-nodes and R-nodes
- Example: Partial upward embedding problem solvable in quadratic time



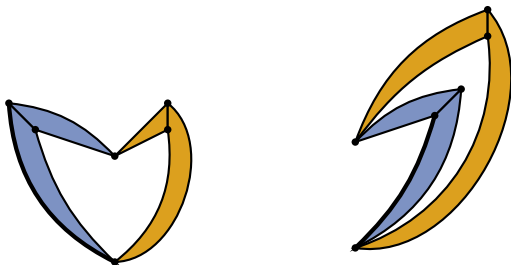
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- NB: Dependency on e^* is necessary:

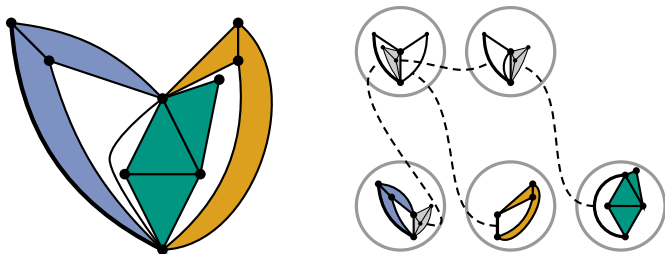


Conclusion

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Future work: Survey more algorithms that use SPQR-trees and adapt to upward planar case