Patrizio Angelini, Henry Förster, Michael Hoffmann, Michael Kaufmann, Stephen Kobourov, Giuseppe Liotta, Maurizio Patrignani

> 27th International Symposium on Graph Drawing and Network Visualization 2019

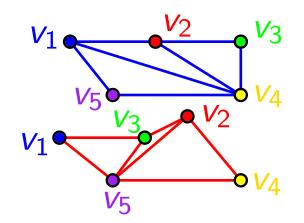
QuaSEFE





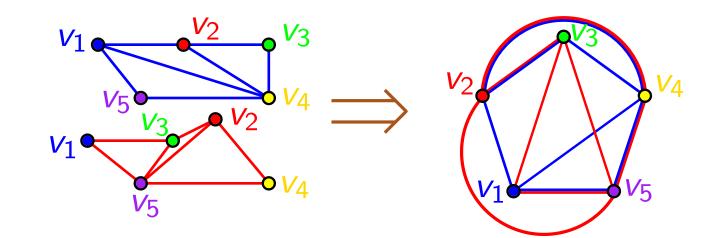


Input: Set of planar graphs with shared vertex set



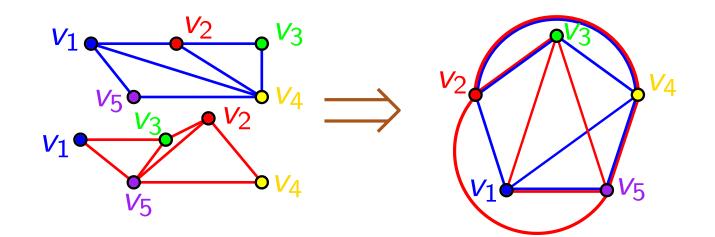


- Input: Set of planar graphs with shared vertex set
- Output: Planar drawings for all graphs such that



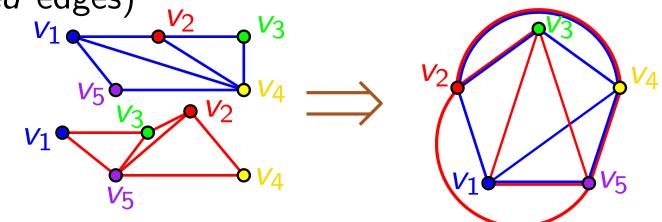
QuaSEFE Simultaneous (Graph) Embedding with Fixed Edges

- Input: Set of planar graphs with shared vertex set
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 - vertices have the same position in all drawings (*simultaneous* drawings)



QuaSEFE Simultaneous (Graph) Embedding with Fixed Edges

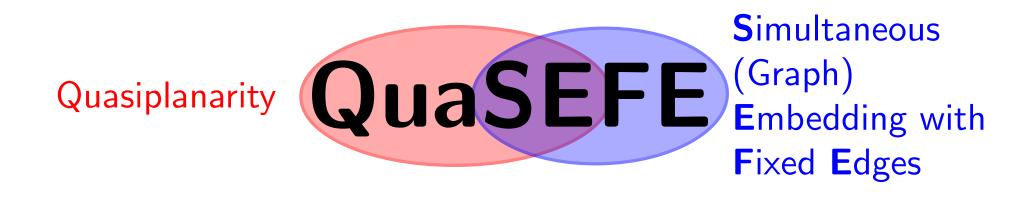
- Input: Set of planar graphs with shared vertex set
- Output: Planar drawings for all graphs such that
 - vertices have the same position in all drawings (*simultaneous* drawings)
 - edges have the same representation in all drawings (fixed edges)

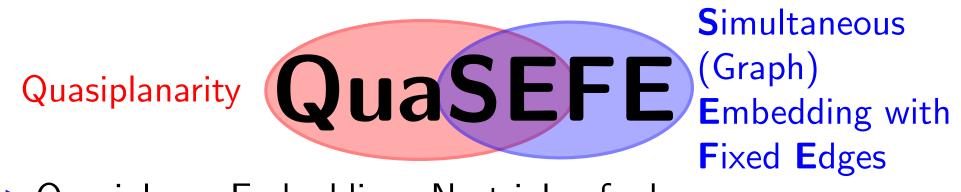


QuaSEFE Simultaneous (Graph) Embedding with Fixed Edges

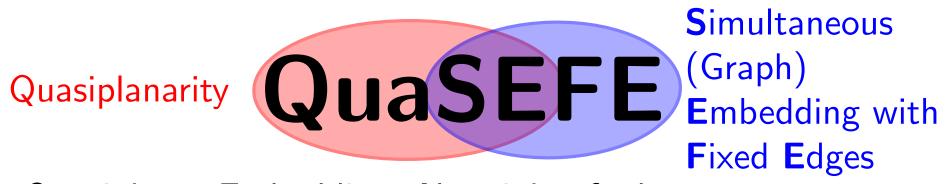




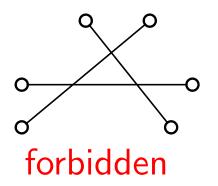




Quasiplanar Embedding: No triple of edges crosses pairwise



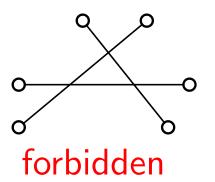
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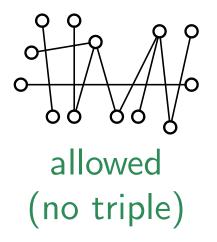


Quasiplanarity QuaSEFE

Simultaneous (Graph) Embedding with Fixed Edges

Quasiplanar Embedding: No triple of edges crosses pairwise

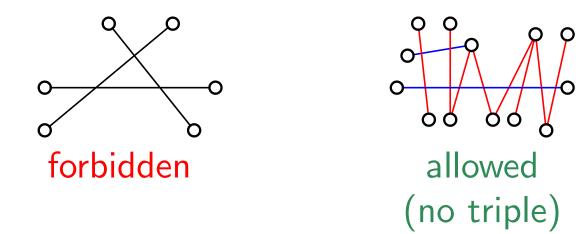




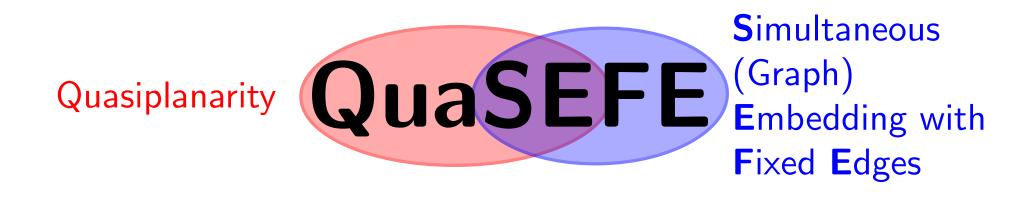
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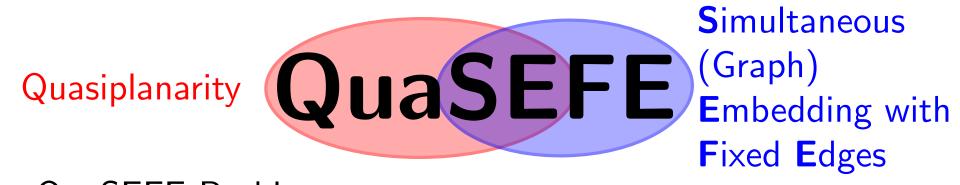
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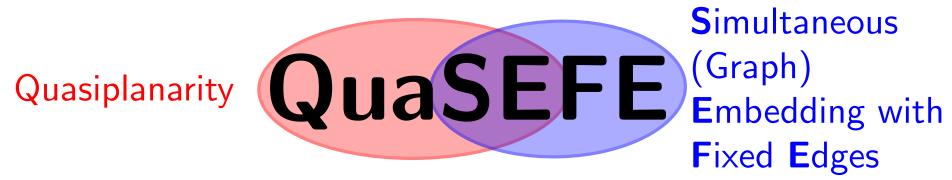


Thickness two drawings (i.e. two-edge colorable drawings) are quasiplanar



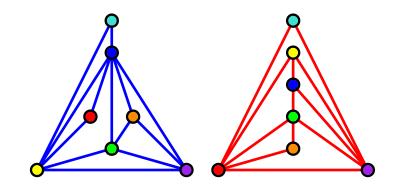


QuaSEFE Problem:



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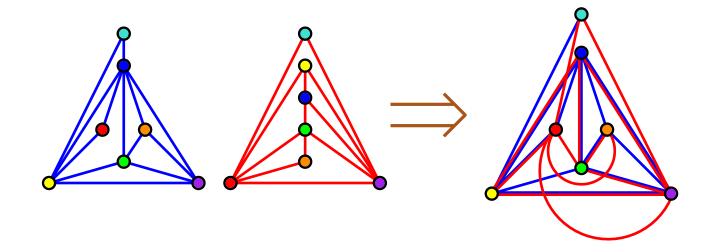
Input: Set of quasiplanar graphs with shared vertex set



Quasiplanarity Quaseful Contraction of the second s

QuaSEFE Problem:

- Input: Set of quasiplanar graphs with shared vertex set
- Output: Simultaneous quasiplanar drawings for all graphs with fixed edges



always positive instances for SEFE

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two caterpillars (in polynomial area) [Brass et al. '06]

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- Variants
 - no fixed mapping between vertices [Brass et al. '06]
 - geometric simultaneous embedding (GSE) [Angelini et al. '11, Di Giacomo et al. '15]

quasiplanar GSE

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- not every two quasiplanar graphs [Di Giacomo et al. '15]

quasiplanar GSE

- ► a tree and a cycle [Didimo et al. '12]
- ► a tree and an outerpillar [Di Giacomo et al. '15]
- ▶ not every two quasiplanar graphs [Di Giacomo et al. '15]
- simultaneous RAC drawings
 - [Argyriou et al. '13, Bekos et al. '16, Evans et al. '16, Grilli '18]

Our Results

always positive instances for QuaSEFE

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- two planar graphs and a tree
- ► a 1-planar graph and a planar graph

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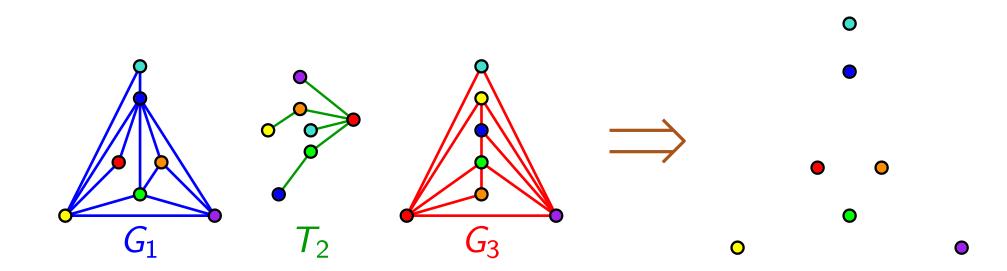
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Our Results

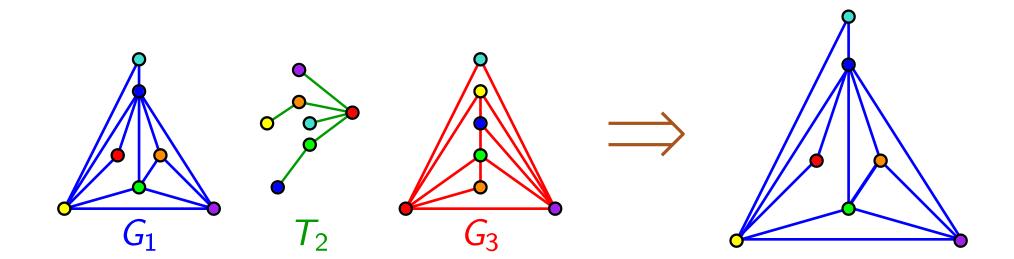
always positive instances for QuaSEFE

- two planar graphs and a tree
- a 1-planar graph and a planar graph
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- counterexamples for QuaSEFE in two special settings

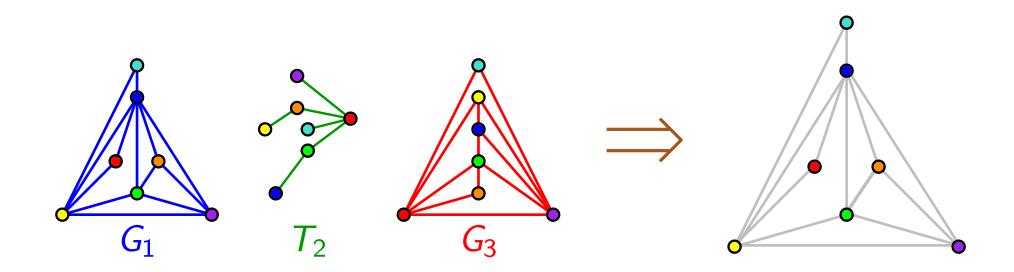
▶ 1. Draw G_1 planar



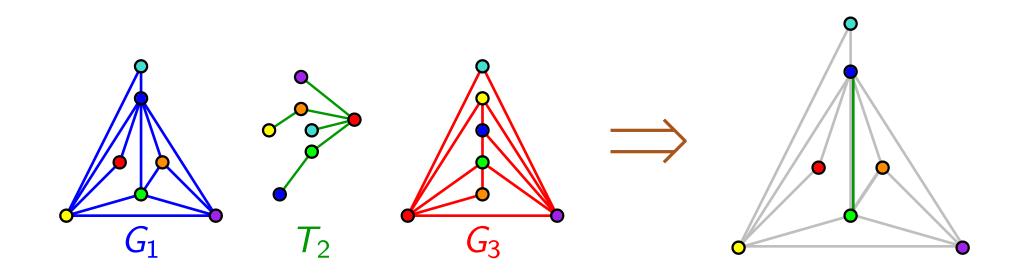
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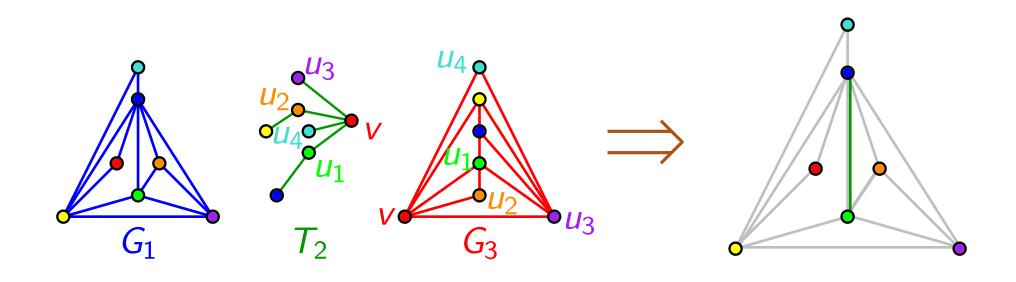
- ▶ 1. Draw G_1 planar
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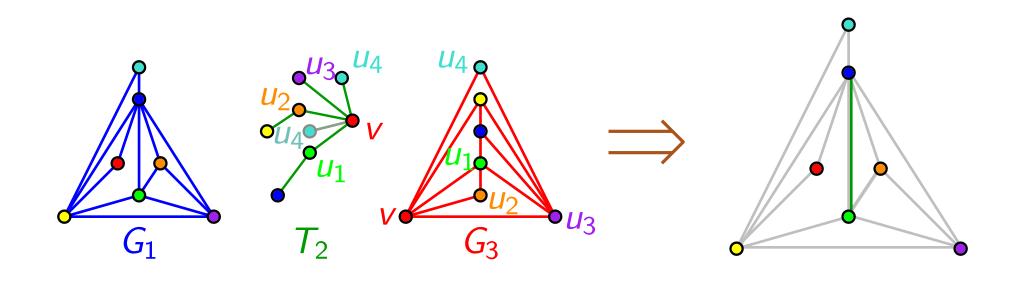
- ▶ 1. Draw G_1 planar
- ▶ 2. Draw T_2 planar
 - ▶ some edges fixed by G_1



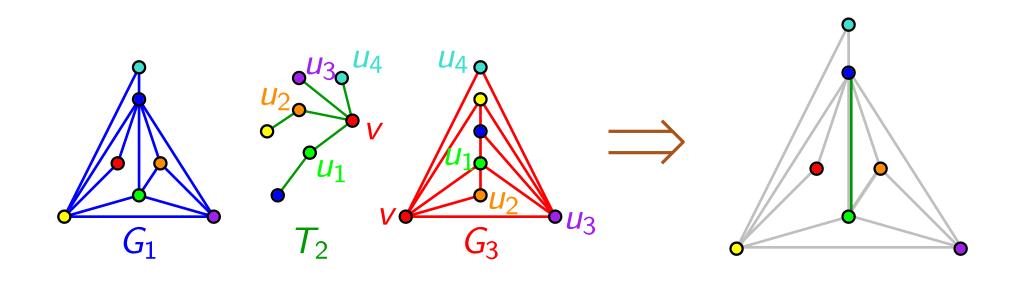
- ▶ 1. Draw G_1 planar
- ► 2. Draw T₂ planar
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 - choose planar rotation system from G_3 for edges in $G_3 \setminus G_1$



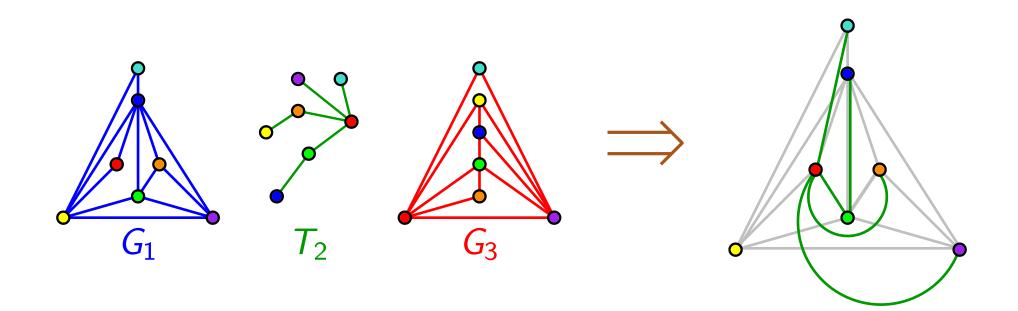
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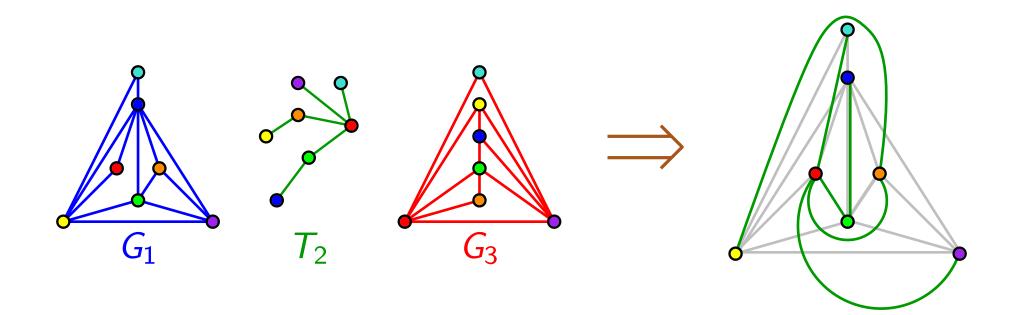
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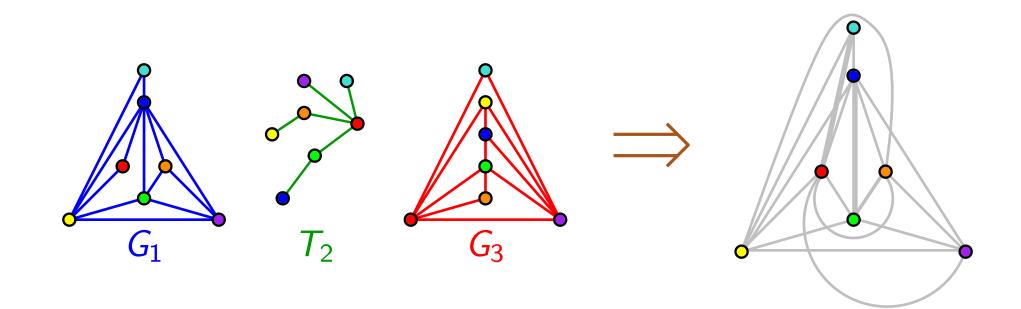
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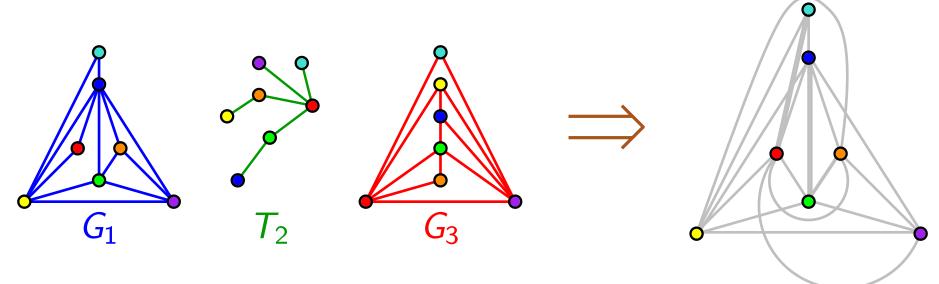
- ▶ 1. Draw G₁ planar
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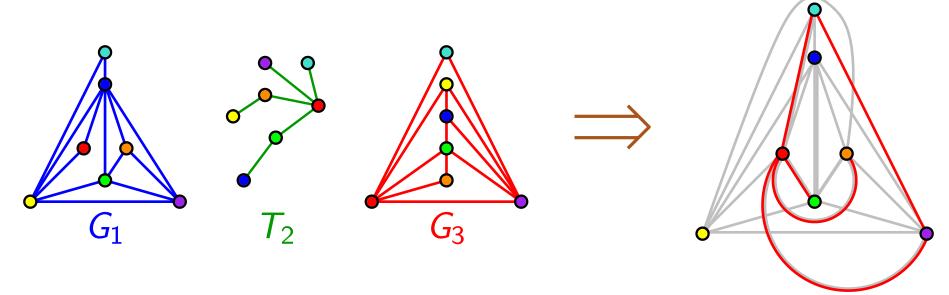
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- ► 3. Draw G₃ quasiplanar



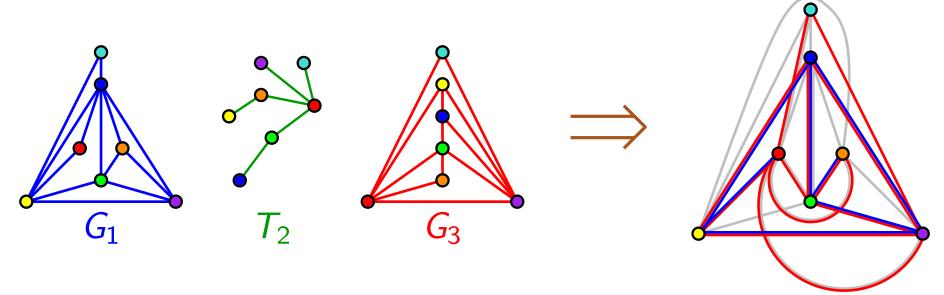
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 - thickness $2 \Rightarrow$ quasiplanar

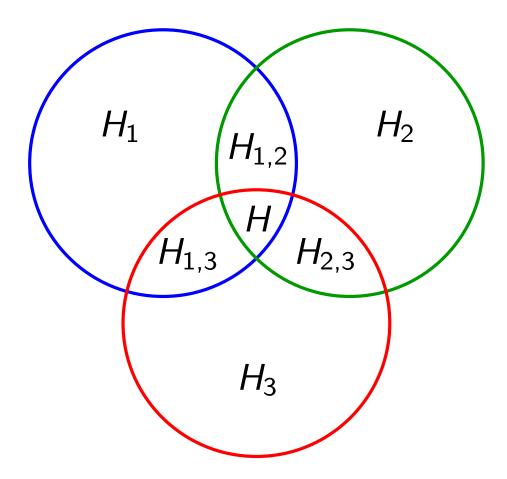
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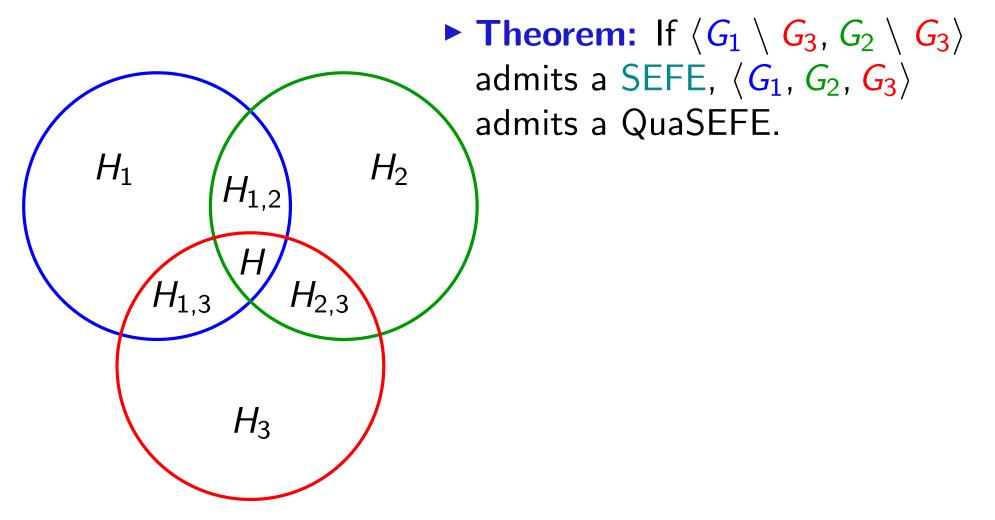
A 1-Planar Graph and a Planar Graph \checkmark

▶ 1. Decompose the 1-planar graph into a planar graph G_1 and a forest T_2 [Ackerman '14]

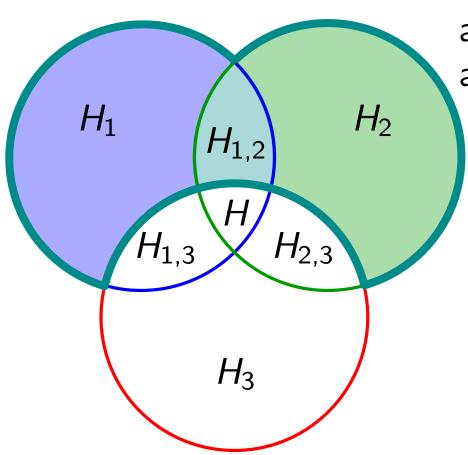
A 1-Planar Graph and a Planar Graph \checkmark

- ▶ 1. Decompose the 1-planar graph into a planar graph G_1 and a forest T_2 [Ackerman '14]
- ▶ 2. Apply the previous result (G_1 and T_2 are planar)

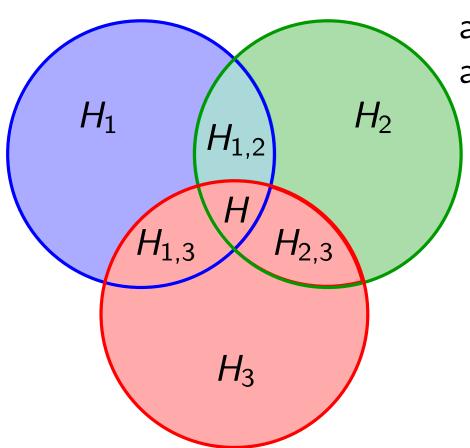




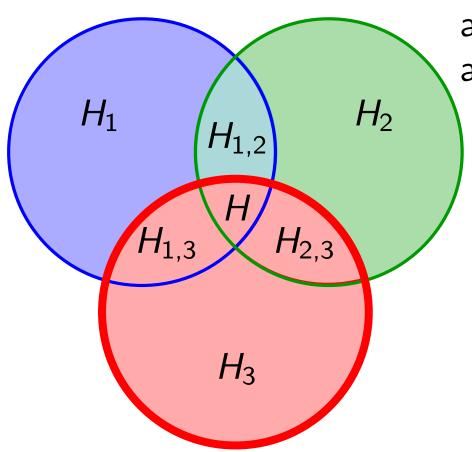
• Let G_1 , G_2 and G_3 planar graphs on V



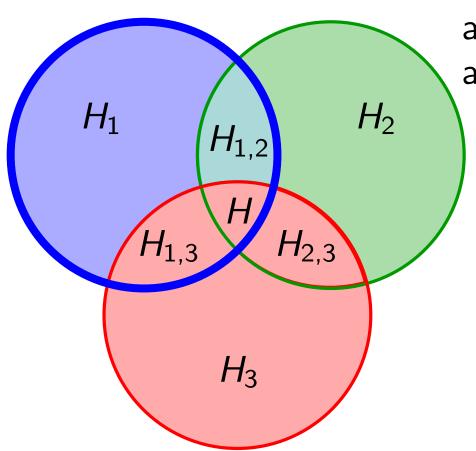
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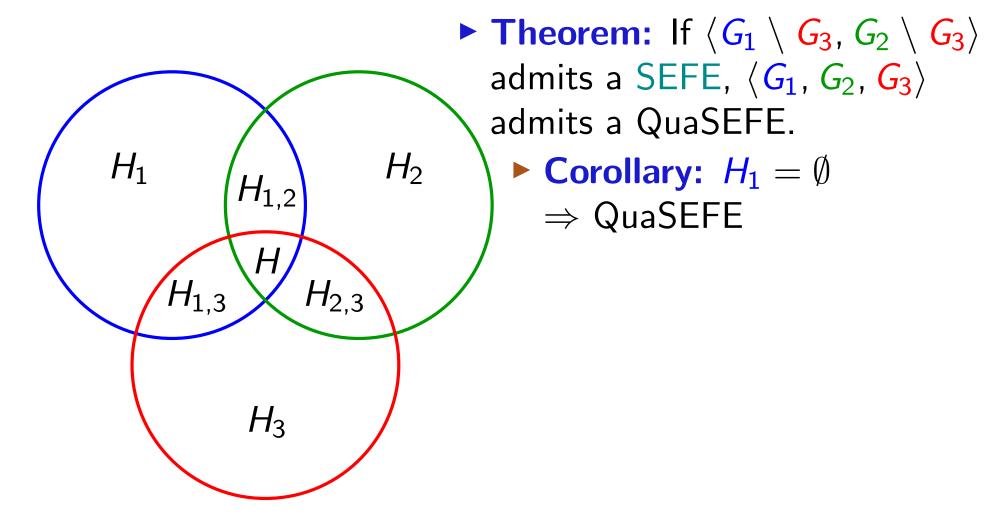


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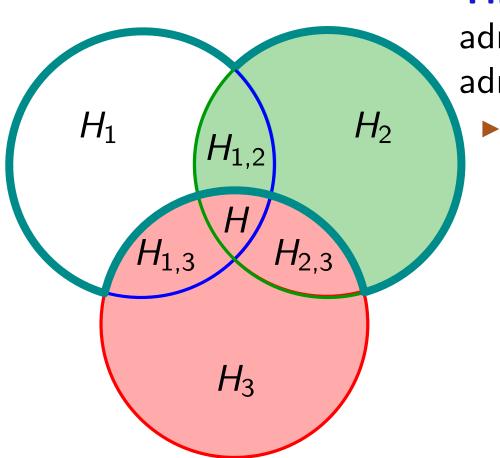


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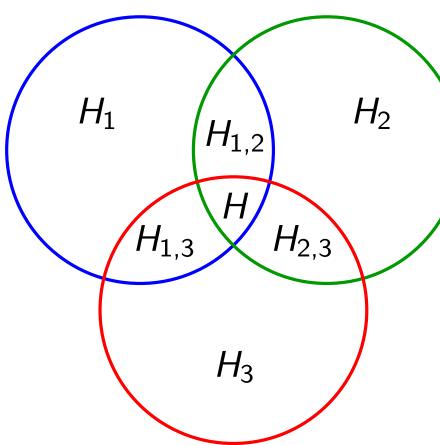
• Let G_1 , G_2 and G_3 planar graphs on V



• Theorem: If $\langle G_1 \setminus G_3, G_2 \setminus G_3 \rangle$ admits a SEFE, $\langle G_1, G_2, G_3 \rangle$ admits a QuaSEFE.

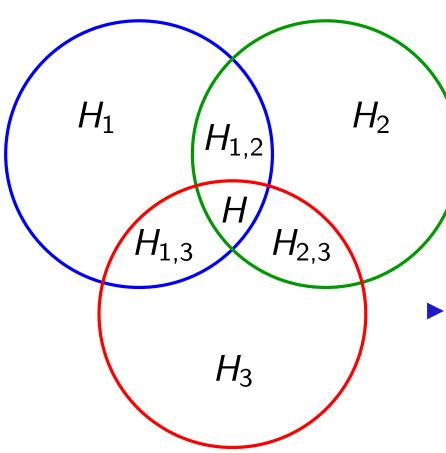
• **Corollary:** $H_1 = \emptyset$ \Rightarrow QuaSEFE

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 - Corollary: $H_{1,2}$ is forest of paths \Rightarrow QuaSEFE
- Theorem: If H is a forest of paths, $\langle G_1, G_2, G_3 \rangle$ admits a QuaSEFE.

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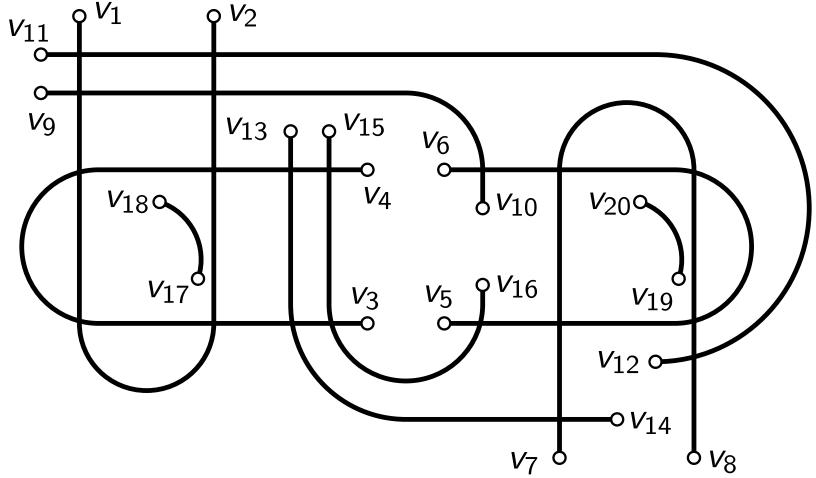
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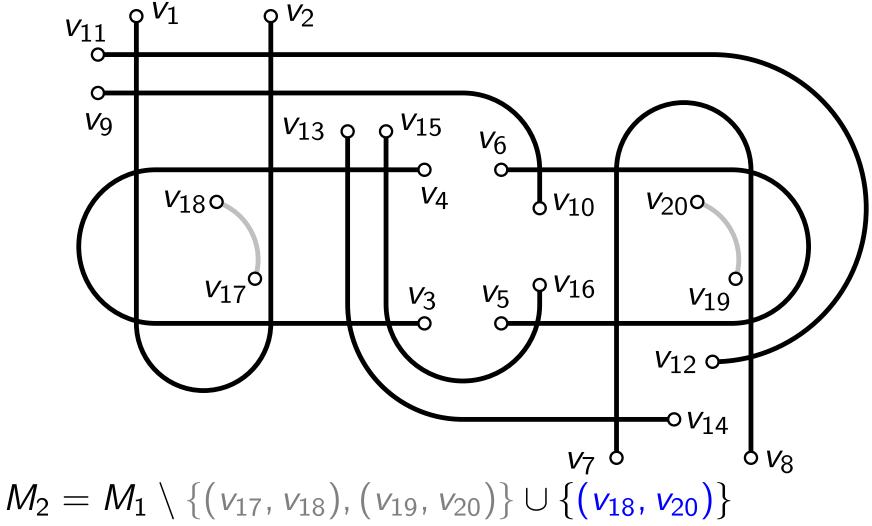
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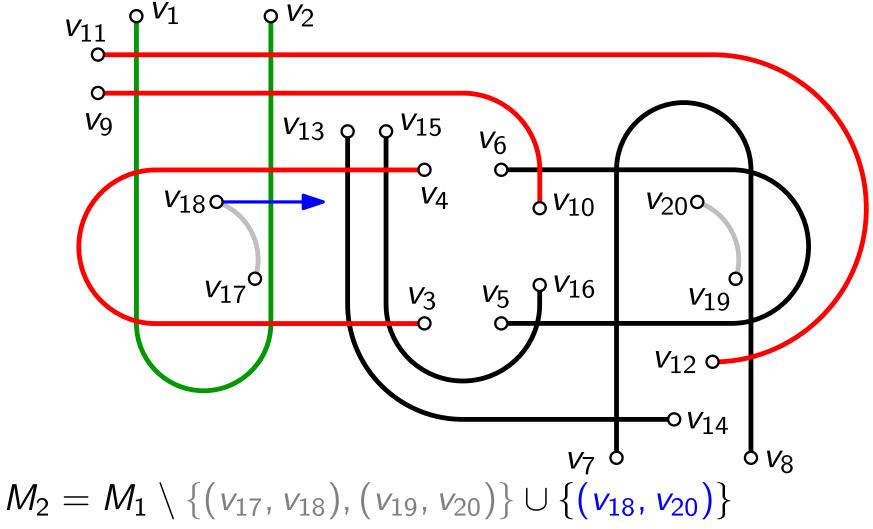
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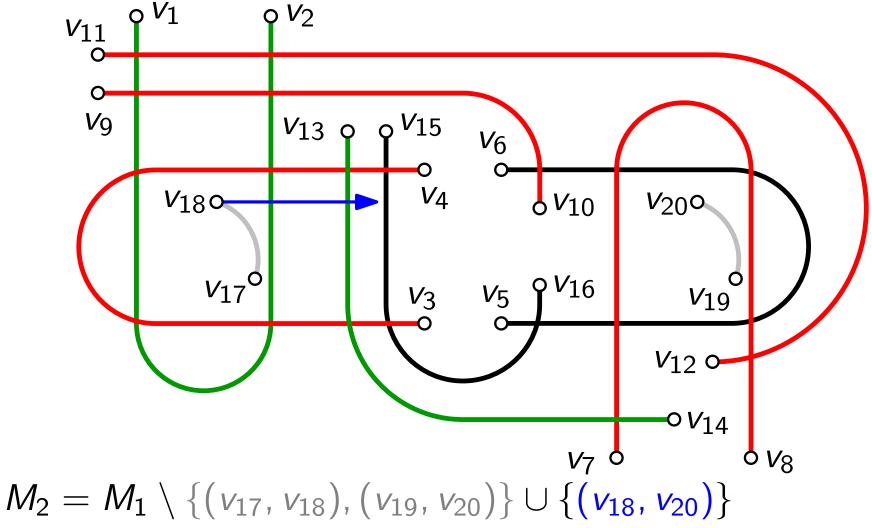
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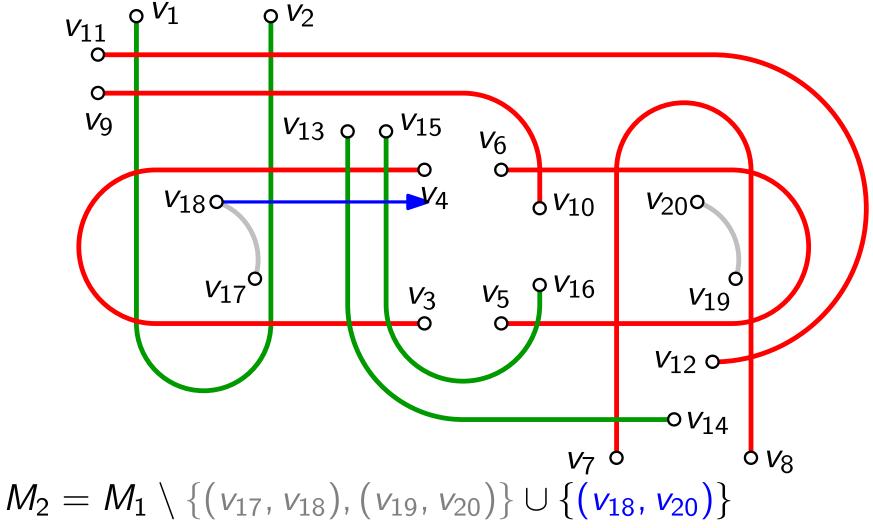
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 - each G_i is drawn with thickness 2

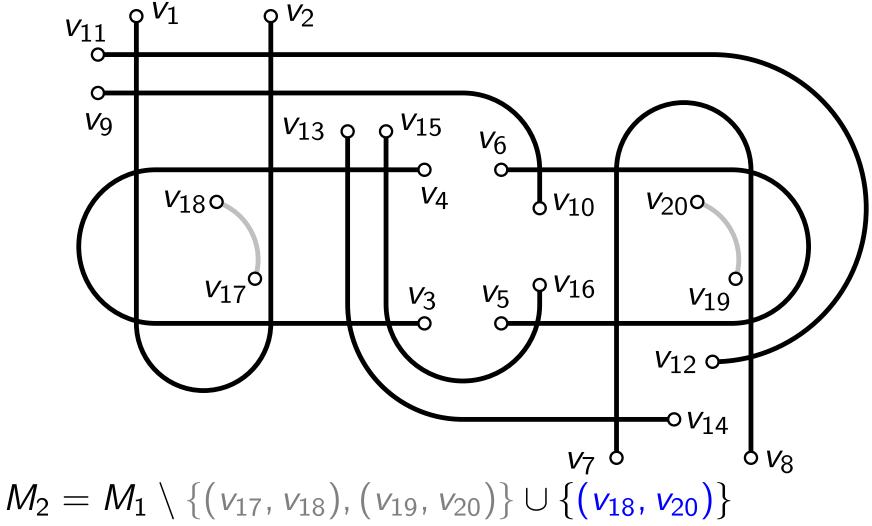


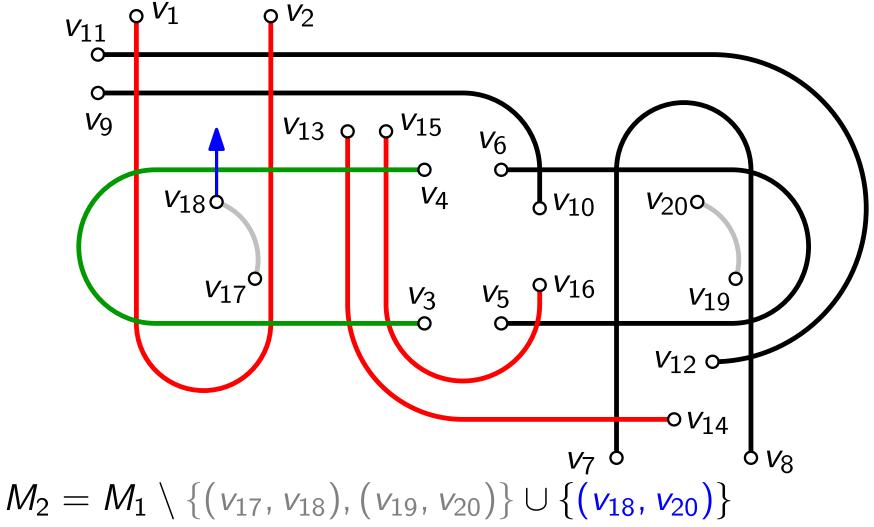


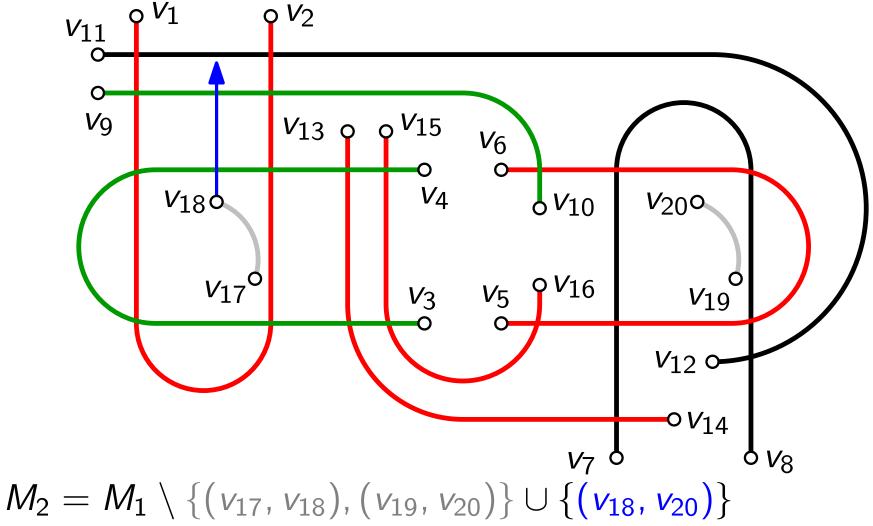


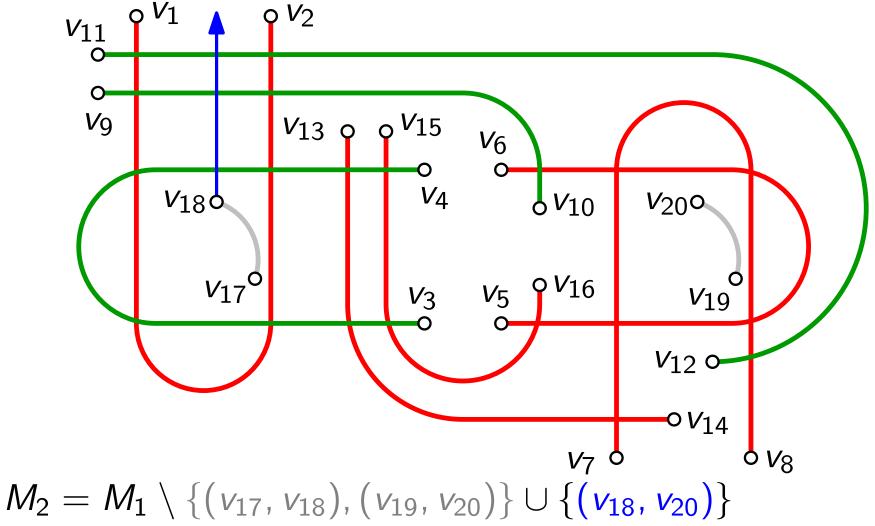












Do the following always admit a QuaSEFE?

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- What is the computational complexity of QuaSEFE?
- Extend to other beyond planar graph classes such as k-planar graphs.

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Thank you for your attention!