

L6 (April 3, 2020) | Let me still refer to ①  
the proof of the equivalence

$(a_n) \subset \mathbb{C} \text{ holon.} \iff \sum_{n=0}^{\infty} a_n x^n \text{ holon.}$  in the previous lecture. In the proof of  $\Leftarrow$  it is simpler to take any  $i_0 \in \{0, 1, \dots, \mathbb{Z}\}$  s.t.  $\exists i_0 \neq \emptyset$  (i.e.  $q_{i_0}(t) \neq 0$ ) and then any  $j_0 \in \mathbb{Z}$ . It is more precise to set  $m = i_0 - j_0$  and to look at the coeff. of  $a_{n+m}$  where  $n$  is a formal variable and  $m \in \mathbb{Z}$  is fixed. This coeff. is a sum of several (and at least one) elements of  $\mathbb{C}[u]$ .  
(nonzero)

Crucially, in this sum the summands,  $\neq 0$  polynomials from  $\mathbb{C}[u]$ , have mutually distinct degrees and therefore their sum is a  $\neq 0$  element of  $\mathbb{C}[u]$ , (the summand with the max. degree cannot be cancelled). ☒

I will pay two debts from L4, the first one being ~~that~~ another equivalence:

a sequence  $(a_n) \subset \mathbb{C}$  is a linear recurrence sequence (LRS)  $\iff$  the OGF  $A(x) = \sum_{n=0}^{\infty} a_n x^n$  is rational,  $= \frac{p(x)}{q(x)}$  with  $p, q \in \mathbb{C}[x]$ ,  $q(0) \neq 0$ .  
(precisely)

First we define LRS. A sequence  $(a_n) =$

$= (a_0, a_1, a_2, \dots) \subset \mathbb{C}$  is a LRS if  $\exists q \in \mathbb{N}_0$  (2)

$\exists c_1, c_2, \dots, c_q \in \mathbb{C}, c_q \neq 0$ , s.t. for every  $n \in \mathbb{N}_0$ ,

$$a_{n+q} = c_1 a_{n+q-1} + c_2 a_{n+q-2} + \dots + c_q a_n. \quad \text{For } q=0 \text{ the R.H.S.}$$

is  $\emptyset$ , we define it as 0 and get the zero sequence with  $a_n = 0$  for every  $n \in \mathbb{N}_0$ . This is the same definition as we gave earlier, only now the recurrence holds for every  $n \in \mathbb{N}_0$  and not ~~just~~ <sup>just</sup> for every  $n > n_0$ . A well known example of a LRS is the Fibonacci numbers  $(F_n)$  with  $F_0 = F_1 = 1$

and  $F_{n+2} = F_{n+1} + F_n$ , so  $(F_n) = (1, 1, 2, 3, 5, 8, 13, 21, 34, 55, 89, 144, \dots)$ .  ~~$(a_n) = (1, 0, 1, 0, 1, 0, \dots)$~~  is also a LRS, given by the recurrence

$$a_{n+2} = a_n \quad (= 0 \cdot a_{n+1} + 1 \cdot a_n). \quad \text{(Exercise for you)}$$

Prove that the eventually constant sequence  $(a_n) = (1, 0, 2, 3, 4, 2020, 2020, 2020, 2020, \dots)$ ,  $a_n = 2020$  for every  $n \geq 5$ , is not a LRS.

The precise version of the above equivalence is as follows.

the Catalan #s

(\*) For the 4 proofs that  $(C_n)$  is not a LRS.

**Theorem**  $(a_n) = (a_0, a_1, \dots) \in \mathbb{C}$  is a LRS (3)



$\exists$  polynomials  $p, q \in \mathbb{C}[x]$  s.t.  $q(0) \neq 0$ ,  $p(x) \neq 0$   
or  $\deg(p) < \deg(q)$ , and  $\sum_{n=0}^{\infty} a_n x^n \stackrel{\text{---}}{=} \frac{p(x)}{q(x)}$ .

Before plunging in the proof we have to explain, of course, what exactly does the equality mean.

It is an equality in the ring  $\mathbb{C}\langle\langle x \rangle\rangle$  of formal power series. In the equality we have three elements of it,  $A(x) = \sum_{n=0}^{\infty} a_n x^n$ ,  $p(x)$  and  $q(x)$ , where  $q(x) \neq 0$ , and the equality claims that

$A(x) = p(x) \cdot q(x)^{-1}$ . Here,  $q(x)^{-1}$  is the multiplicative inverse in  $\mathbb{C}\langle\langle x \rangle\rangle$  to  $q(x)$ . So how are these inverses defined? The elements of  $\mathbb{C}\langle\langle x \rangle\rangle$  are the formal linear combinations  $(*)$

$\sum_{n=0}^{\infty} b_n x^n$  of the powers  $x^n$ ,  $n \in \mathbb{N}_0$ , with complex coefficients  $b_n \in \mathbb{C}$ . Addition and multiplication are defined in  $\mathbb{C}\langle\langle x \rangle\rangle$  as follows.  $x$  is a formal variable.

( $(a_n)$  is now a generic sequence, not the one ~~from~~ <sup>(4)</sup> the above theorem)  $\sum_{n \geq 0} a_n x^n + \sum_{n \geq 0} b_n x^n =$

$$:= \sum_{n \geq 0} (a_n + b_n) x^n$$

addition in  $\mathbb{C}$

and

addition in  $\mathbb{C}[[x]]$ , defined by this equality

$$\sum_{n \geq 0} a_n x^n \cdot \sum_{n \geq 0} b_n x^n :=$$

$$\sum_{n \geq 0} \left( \sum_{q=0}^n a_q b_{n-q} \right) x^n$$

multiplication in  $\mathbb{C}[[x]]$

multiplication and addition in  $\mathbb{C}$ .

This is also sometimes called the Cauchy product of (formal) power series. Exercise

Show that  $(\mathbb{C}[[x]], +, \cdot)$ , with the neutral elements  $0 = 0 \cdot x^0 + 0 \cdot x^1 + \dots$  and  $1 = 1 \cdot x^0 + 0 \cdot x^1 + 0 \cdot x^2 + \dots$ , is a ring (i.e. satisfies ~~the~~ <sup>all</sup> axioms of a ring)

Proposition  $\mathbb{C}[[x]]$  is an integral domain (\*)

Proof: Let  $\sum_{n \geq 0} a_n x^n, \sum_{n \geq 0} b_n x^n$  be  $\neq 0$  elements of  $\mathbb{C}[[x]]$ . Let  $n_0$  be the minimum  $n$  s.t.  $a_n \neq 0$  and  $n_0$  be  $-1$  —  $b_n \neq 0$ . Then  $\sum_{n \geq 0} a_n x^n \cdot \sum_{n \geq 0} b_n x^n =$

(\*)  $R$  is an i.d.;  $a, b \in R, ab = 0_R \Rightarrow a = 0_R \vee b = 0_R$

$= a_{u_0} b_{u_0} x^{u_0 + u_0} + \dots \neq 0$  in  $\mathbb{C}[x]$ . (5)

$\neq 0$  Thus  $\mathbb{C}[x]$  is an integral domain.

the minimum  $n \in \mathbb{N}_0$  s.t.  $a_n \neq 0$  in a  $\sum_{u \geq 0} a_u x^u \in \mathbb{C}[x]$

$0 \neq A(x) \in \mathbb{C}[x]$

is in fact called the order of  $A(x)$  and denoted  $\text{ord}(A(x))$ . We set  $\text{ord}(0) = -\infty$ . (Exercise)

Prove that for every  $A, B \in \mathbb{C}[x]$ ,  $\text{ord}(A \cdot B) = \text{ord}(A) + \text{ord}(B)$ .

Now we come to the question of units in the ring  $\mathbb{C}[x]$  (i.e. invertible elements) which has to be clarified before we <sup>can</sup> prove the above theorem.

Proposition  $A(x) \in \mathbb{C}[x]$  is a unit if and only if  $\text{ord}(A(x)) = 0$ , i.e.  $A(x) = a_0 + a_1 x + a_2 x^2 + \dots$  with  $a_0 \neq 0$ .

Proof. Suppose that  $A(x)$  is a unit, there is a  $B(x) \in \mathbb{C}[x]$  s.t.  $A(x)B(x) = 1 (= 1x^0 + 0x^1 + \dots)$ . By the exercise  $\uparrow$ ,  $\text{ord}(A) + \text{ord}(B) = \text{ord}(1) = 0$ , which implies that  $\text{ord}(A) + \text{ord}(B) = 0$  (because the values of  $\text{ord}(\cdot)$  are  $\mathbb{N}_0 \cup \{-\infty\}$ ). Suppose that  $\text{ord}(A) = 0$ . We are looking for

a  $B(x) = b_0 + b_1x + \dots \in \mathbb{C}[x] \cap \mathbb{H}$  s.t.  $A \cdot B = 1$ . (6)

$B''$  This is equivalent with the infinite linear non-homogeneous system of equations (\*)

$$a_0 b_0 = 1, a_0 b_1 + a_1 b_0 = 0, a_0 b_2 + a_1 b_1 + a_2 b_0 = 0, \dots$$

where  $a_0, a_1, a_2, \dots \in \mathbb{C}$  are given and  $b_0, b_1, b_2, \dots$  are the unknowns. Since  $a_0 \neq 0$ , we see that the system has a unique solution  $b_0, b_1, \dots \in \mathbb{C}$ :

$$b_0 = \frac{1}{a_0}, b_1 = \frac{1}{a_0} (-a_1 b_0), b_2 = \frac{1}{a_0} (-a_1 b_1 -$$

$-a_2 b_0),$  and so on. Thus  $A$  has an inverse  $B$  and is a unit.  $\square$

Now the condition  $q(0) \neq 0$  in the theorem above is clear. Of course, we take  $\mathbb{C}[x]$  as contained in  $\mathbb{C}[x] \cap \mathbb{H}$ . Exercise In  $\mathbb{C}[x] \cap \mathbb{H}$ ,  $\frac{1}{1-x} = ?$

And  $\frac{1}{1+x} = ?$

Now we are theoretically prepared <sup>to prove</sup> ~~for the proof of~~ the above theorem on LRS but of course do not have time for it. ~~and~~ <sup>I will</sup> present the proof in the next lecture. See you next week.

$(\Leftrightarrow)$

(\*) in better English ~~is~~: in f. system of non-h. lin. eq-s,