

L11 May 15, 2020 | The rings U and V are non-commutative and have zero divisors (non-zero elements with zero product). Let us denote, for $q \in \mathbb{N}_0$ and x a formal variable, by $[x^q]$ ($a_0 + a_1x + a_2x^2 + \dots$) $:= a_q$, the coeff. of x^q in the expansion following after the symbol $[x^q]$. ①

Proposition (an isomorphism)

Let R be a ring (commutative and with 1_R) and $U = R[[x]]^{n \times n}$ and $V = R^{n \times n}[[x]]$ be the ring of $n \times n$ matrices whose entries are (formal power series) $a_0 + a_1x + a_2x^2 + \dots$ with $a_i \in R$ and the ring of formal power series $A_0 + A_1x + A_2x^2 + \dots$ with $A_i \in R^{n \times n}$ ($n \times n$ matrices with entries in R), respectively. The maps

$$\alpha: U \rightarrow V, M \mapsto \sum_{q=0}^{\infty} ([x^q] M_{i,j})_{i,j=1}^n x^q$$

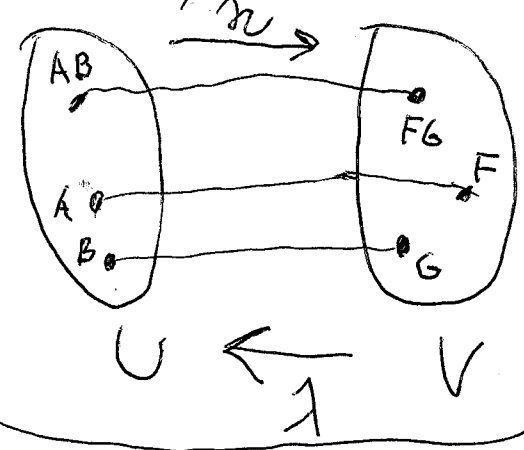
and

$$\lambda: V \rightarrow U, \sum_{q=0}^{\infty} M(q) x^q \mapsto \left(\sum_{q=0}^{\infty} M(q)_{i,j} x^q \right)_{i,j=1}^n$$

are mutually inverse homomorphisms of

rings. Thus α and β are isomorphisms of the rings U and V . ②

Proof. It is clear that α and β send $1_U \mapsto 1_V$ and $1_V \mapsto 1_U$ respectively. Also, α and β preserve subtraction, i.e. for any $A, B \in U$ and $F, G \in V$ we have that $\alpha(A-B) = \alpha(A) - \alpha(B)$ and $\beta(F-G) = \beta(F) - \beta(G)$ (check the definitions of α and β). We show that they also preserve ~~the~~ multiplication; it is only a challenge in notation. Since $\alpha \circ \beta$ and $\beta \circ \alpha$ are identical maps, α and β are mutually inverse bijections and it suffices only to check multiplicativity of, say, the map α :



Let $F, G \in V$ be given ^(formal) power series with

$$F = \sum_{k=0}^{\infty} \eta(k) x^k \text{ and } G = \sum_{k=0}^{\infty} \nu(k) x^k, \text{ where } \eta(k), \nu(k) \in \mathbb{R}^{h \times n}.$$

Then $\alpha(FG) \stackrel{(p.u.)}{=} \sum_{k=0}^{\infty} \left(\sum_{l=0}^k \eta(l) \nu(k-l) \right) x^k =$

be given power series matrix coeff- s , $F =$

$= \lambda \left(\sum_{q=0}^{\infty} P(q) x^q \right)$ where the matrices $P(q)$ ~~are~~ ⁽³⁾ have entries $(c_{ij} \in [n])$

$P(q) \stackrel{(m.m.)}{=} \sum_{l=0}^q \sum_{i_1=1}^n h(l) c_{i_1 i_1}^{N(q-l)} \cdot c_{i_1 i_1}$

Therefore $\lambda(FG) \stackrel{\text{def. of } \lambda}{=} \left(\sum_{q=0}^{\infty} P(q) c_{ij} x^q \right)_{i,j=1}^n$

$\stackrel{(m.m.)}{=} \stackrel{(p.m.)}{=} \left(\sum_{q=0}^{\infty} \left(\sum_{l=0}^q \sum_{i_1=1}^n h(l) c_{i_1 i_1}^{N(q-l)} c_{i_1 i_1} \right) x^q \right)_{i,j=1}^n$ • On

the other hand, $\lambda(F) \lambda(G) \stackrel{\text{def. of } \lambda}{=} \left(\sum_{q=0}^{\infty} h(q) c_{ij} x^q \right)_{i,j=1}^n \left(\sum_{q=0}^{\infty} N(q) c_{ij} x^q \right)_{i,j=1}^n$

$\stackrel{(p.m.)}{=} \left(\sum_{i_1=1}^n \left(\sum_{q=0}^{\infty} h(q) c_{i_1 i_1} x^q \right) \left(\sum_{q=0}^{\infty} N(q) c_{i_1 i_1} x^q \right) \right)_{i,j=1}^n$

$\stackrel{(p.m.)}{=} \left(\sum_{i_1=1}^n \sum_{q=0}^{\infty} \left(\sum_{l=0}^q h(l) c_{i_1 i_1}^{N(q-l)} c_{i_1 i_1} \right) x^q \right)_{i,j=1}^n$

$\stackrel{(p.m.)}{=} \left(\sum_{i_1=1}^n \sum_{q=0}^{\infty} \left(\sum_{l=0}^q h(l) c_{i_1 i_1}^{N(q-l)} c_{i_1 i_1} \right) x^q \right)_{i,j=1}^n$

Above (λ) , $(m.m.)$ and $(p.m.)$ mean that the def. of λ , matrix multiplication and formal power series multiplication was applied, respectively.

In the last \square we change order of summation, (4)
~~move~~ move the finite outer sum $\sum_{i_1=1}^n \dots$ completely in-
 side and get the previous \square . Hence $\lambda(F)\lambda(G) =$
 $= \lambda(FG)$. Why could we move the $\sum_{i_1=1}^n \dots$ comple-
 tely inside? The first exchange

$$\sum_{i_1=1}^n \sum_{z=0}^{\infty} \rightarrow \sum_{z=0}^{\infty} \sum_{i_1=1}^n$$

is in fact an application of the ~~coefficient-wise~~ addi-
 tion of f . power series

definition of

as ~~coefficient-wise~~ coefficient-wise addi-
 tion. The next exchange $\sum_{i_1=1}^n \sum_{l=0}^{\infty} \rightarrow \sum_{l=0}^{\infty} \sum_{i_1=1}^n$ is
 the combinatorial

double counting trick, which

we can have state abstractly ~~and generally~~ as follows.

Lemma (double counting) Let $S = (S, +)$ be an asso-
 ciative and commutative groupoid (i.e. $+$ is a binary
 operation on S which is a. and c.). If $a_{i,j} \in S$
 for $i \in [m]$ and $j \in [n]$ ($m, n \in \mathbb{N}$) then

$$\sum_{i=1}^m \sum_{j=1}^n a_{i,j} = \sum_{j=1}^n \sum_{i=1}^m a_{i,j}$$

Proof
 An exercise for
 you. \square

This finishes the proof of the Proposition \square

Let us revisit in the perspective of the isomorphisms α and λ the derivation of the TMM formula. For the given $A \in R^{k \times n}$ we take the matrix

$$M := \left(\sum_{q=0}^{\infty} (A^q)_{i,j} x^q \right)_{i,j=1}^n \in U. \text{ We apply } \alpha \text{ and}$$

get the f. power series $F := \alpha(M) = \sum_{q=0}^{\infty} A^q x^q \in V.$

$$In V, \left(\sum_{q=0}^{\infty} A^q x^q \right) (\cancel{I} x^0 - Ax) = \sum_{q=0}^{\infty} A^q x^q - \sum_{q=0}^{\infty} A^{q+1} x^q =$$

$$= I x^0. \text{ So in } V \text{ we have that } F = (I x^0 - Ax)^{-1} \text{ (Since } V \text{ is a ring of f. power series with non-commutative coeff-s, one has to check carefully applicability of the geometric series formula.)}$$

Some isomorphisms of rings preserve inverse elements, by applying λ we therefore get that $\lambda(F) \in U$ actually is $\lambda(F) = (I - xA)^{-1}.$ Applying the (Cramer's) formula for inverse matrix we obtain that

$$N_{ij} = \frac{(-1)^{i+j} \det((I-xA) [i,j])}{\det(I-xA)} \Big)_{i,j=1}^n$$

Hence

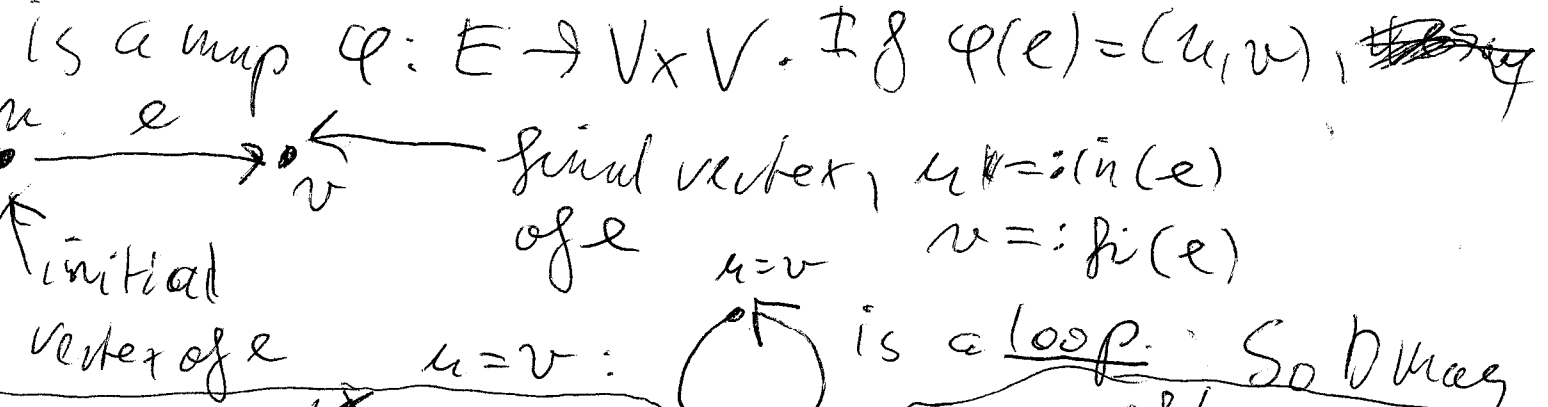
(*) Formalized proof of the d. counting lemma could be of interest.

$$N = \lambda(F) = \lambda(\alpha(M)) = (\lambda \circ \alpha)(M) \quad (6)$$

thus we get the TOM formula that $\forall i, j \in [n]$

$$M_{ij} = N_{ij} \cdot \left[\text{Let us now finally turn to the combinatorial side of the TOM formula. We have a directed multigraph (shortly digraph) } D = (V, E, \varphi) \text{ where } V = \{v_1, v_2, \dots, v_p\} \text{ are the vertices, } E \text{ are the edges and } \varphi \text{ is a map } \varphi: E \rightarrow V \times V. \text{ If } \varphi(e) = (u, v), \text{ then } u \text{ is the initial vertex of } e \text{ and } v \text{ is the final vertex of } e. \text{ If } u = v, \text{ then } e \text{ is a loop. So } D \text{ may look like: } E = \{e_1, \dots, e_6\} \text{ or } V = \{a, b, c\} \text{ such that } \dots \text{ and } \varphi_i(e_i) = \text{in}(e_{i+1}) \text{ for } i = 1, 2, \dots, n-1. \text{ It is closed if } u = v. \text{ We consider a weight function } w: E \rightarrow R, \text{ where } R \text{ is any ring (comm.) with } 1. \text{ To be continued in the last lecture} \right]$$

Let us now finally turn to the combinatorial side of the TOM formula. We have a directed multigraph (shortly digraph) $D = (V, E, \varphi)$ where $V = \{v_1, v_2, \dots, v_p\}$ are the vertices, E are the edges and φ is a map $\varphi: E \rightarrow V \times V$. If $\varphi(e) = (u, v)$, then u is the initial vertex of e and v is the final vertex of e . If $u = v$, then e is a loop. So D may look like:



sequence such that $\varphi_i(e_i) = \text{in}(e_{i+1})$ for $i = 1, 2, \dots, n-1$. It is closed if $u = v$. We consider a weight function $w: E \rightarrow R$, where R is any ring (comm.) with 1 .

To be continued in the last lecture