

MATHEMATICAL STRUCTURES (NMAI064)

summer term 2024/25

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**LECTURE 10 (April 23, 2025)**

AXIOM OF CHOICE AND ITS CONSEQUENCES:  
NON-MEASURABLE SETS, THE WELL ORDERING  
THEOREM, THE PROPHET PARADOX

- *The Axiom of Choice (AC)* is the set-theoretic axiom that

$$\forall A: \emptyset \notin A \Rightarrow \exists F: (F: A \rightarrow \bigcup A) \wedge (B \in A \Rightarrow F(B) \in B).$$

As you certainly know, the *sum*  $\bigcup A$  of  $A$ , is the set  $\bigcup A$  such that  $B \in \bigcup A \iff \exists C \in A: B \in C$ . The notation  $F: A \rightarrow B$ , i.e.,  $F$  is a *function (map) from  $A$  to  $B$* , abbreviates the fact that  $F$  is a set of ordered pairs  $(C, D)$  such that always  $C \in A$ ,  $D \in B$ , and for every  $C \in A$  there exists exactly one  $D \in B$  with  $(C, D) \in F$ .

**Exercise 1** Show that the AC is equivalent with the claim that for every surjection  $F: A \rightarrow B$  there is a map  $G: B \rightarrow A$  such that

$$F(G) = F \circ G = \text{id}_B.$$

**Exercise 2** Show that the AC is equivalent with the claim that for every set system  $\{A_i: i \in I\}$ ,  $A_i \neq \emptyset$ , there is a map

$$F: I \rightarrow \bigcup_{i \in I} A_i$$

such that  $F(i) \in A_i$  for every  $i \in I$ .

- *Equivalences and partitions.* First let us review equivalence relations and set partitions.  $R \subseteq A \times A$  is an *equivalence relation* on  $A$  if it is

- reflexive –  $\forall a \in A: aRa$ ,
- symmetric –  $\forall a, b \in A: aRb \Rightarrow bRa$ , and
- transitive –  $\forall a, b, c \in A: aRb \wedge bRc \Rightarrow aRc$

A set *partition* of a set  $A$  is a set  $B$  such that  $\emptyset \notin B$ , the elements of  $B$  are mutually disjoint and  $\bigcup B = A$ . For any equivalence relation  $R$  on a set  $A$  we define the *blocks of  $R$*  to be the sets

$$[a]_R = \{b \in A: aRb\}, a \in A.$$

**Exercise 3** For every set  $A$  and every equivalence relation  $R$  on  $A$ ,

$$A/R := \{[a]_R : a \in A\}$$

is a partition of  $A$ .

**Exercise 4** For every set  $A$  and every partition  $P$  of  $A$ ,

$$R(P) := \{(a, b) \in A^2 : \exists B \in P : a, b \in B\}$$

is an equivalence relation on  $A$ .

**Exercise 5** For every set  $A$ , every equivalence relation  $S$  on  $A$  and every partition  $P$  of  $A$ ,

$$R(A/S) = S \text{ and } A/R(P) = P.$$

**Exercise 6** For  $n \in \mathbb{N} = \{1, 2, \dots\}$  let  $B_n$ , the Bell number<sup>1</sup>, be the number of equivalence relations on an  $n$ -element set  $X$ . Why does  $B_n$  depend only on the cardinality of  $X$  and not on the elements of  $X$ ? Prove that for every  $n$ ,

$$B_n < B_{n+1}.$$

• *Non-measurable sets.* Let

$$S = \{(x, y) \in \mathbb{R}^2 : x^2 + y^2 = 1\}$$

be the *unit circle* in the Euclidean plane  $\mathbb{R}^2$ . For any angle  $\varphi \in [0, 2\pi)$  we denote by

$$F_\varphi : S \rightarrow S, \quad (x, y) \mapsto (?_x, ?_y),$$

the *counter-clockwise rotation around the origin by the angle  $\varphi$* . It is clearly a bijection. An angle  $\varphi \in [0, 2\pi)$  is *rational* if  $\frac{\varphi}{\pi} \in \mathbb{Q}$ . We denote the set of rational angles by  $[0, 2\pi)_{\mathbb{Q}}$ . Obviously,  $[0, 2\pi)_{\mathbb{Q}}$  is a countable set.

**Exercise 7** Define the additive Abelian group

$$([0, 2\pi)_{\mathbb{Q}}, +)$$

of addition modulo  $2\pi$ . Find the above formulas  $?_x$  and  $?_y$  in the definition of  $F_\varphi$  and show that for any  $\varphi, \varphi' \in [0, 2\pi)_{\mathbb{Q}}$ ,

$$F_\varphi \circ F_{\varphi'} = F_{\varphi+\varphi'}.$$

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<sup>1</sup>Named after Eric T. Bell (1883–1960).

Show that for any fixed  $x \in S$ , the function  $F_\varphi(x)$  is injective in the variable  $\varphi \in [0, 2\pi)$ .

For the unit circle  $S$  we denote by  $\mathcal{P}(S)$  the set of subsets of  $S$ . For a subset  $X \subseteq \mathcal{P}(S)$  with  $S \in X$ , we say that a map

$$\lambda: X \rightarrow [0, +\infty)$$

is an *arc length on  $X$*  if the following three conditions hold.

1.  $\lambda(S) > 0$  — the whole unit circle has positive arc length.
2. For every pairwise disjoint sets  $A_n \in X$ ,  $n \in \mathbb{N}$ , with  $\bigcup_{n=1}^{\infty} A_n \in X$  one has that

$$\lambda\left(\bigcup_{n=1}^{\infty} A_n\right) = \sum_{n=1}^{\infty} \lambda(A_n).$$

We say that the arc length is  $\sigma$ -additive.

3. For every  $\varphi \in [0, 2\pi)$  and every  $A \in X$ , if  $F_\varphi[A] \in X$  then

$$\lambda(F_\varphi[A]) = \lambda(A).$$

We say that the arc length is invariant under rotations.

**Theorem 8 (a troublesome set)** *There exists a set  $X \subseteq S$  such that the set*

$$\{F_\varphi[X]: \varphi \in [0, 2\pi)_{\mathbb{Q}}\}$$

*is a partition of  $S$ .*

**Proof.** By Exercise 9, the relation  $\sim$  on  $S$ , defined by

$$a \sim b \iff \exists \varphi \in [0, 2\pi)_{\mathbb{Q}}: F_\varphi(a) = b,$$

is an equivalence relation. We define  $X \subseteq S$  by means of the AC by taking one representative element from each block of  $\sim$ . We show that for  $\varphi$  running in  $[0, 2\pi)_{\mathbb{Q}}$  the sets  $F_\varphi[X]$  are disjoint and form a partition of  $S$ . Their union is  $S$  because each  $s \in S$  lies in a block  $B$  of  $\sim$  and thus  $F_\varphi(r) = s$  for some  $\varphi \in [0, 2\pi)_{\mathbb{Q}}$  for the representative  $r \in X$  of  $B$ . If  $F_\varphi[X] \cap F_{\varphi'}[X] \neq \emptyset$  for two distinct rational angles  $\varphi$  and  $\varphi'$ , then

$$F_\varphi(r) = F_{\varphi'}(r') \text{ for some } r, r' \in X.$$

Then  $r \neq r'$  by the injectivity of  $F_\varphi(x)$  in  $\varphi$  for fixed  $x$  (Exercise 7). Also,

$$F_{\varphi-\varphi'}(r) = r' \quad \text{for } \varphi - \varphi' \in [0, 2\pi)_{\mathbb{Q}}$$

(again by Exercise 7) and therefore  $r \sim r'$ . This is impossible for two distinct elements of  $X$ . It is clear that always  $F_\varphi[X] \neq \emptyset$ .  $\square$

**Exercise 9** Prove that the relation  $\sim$  on  $S$  defined in the previous proof is an equivalence relation.

**Corollary 10 (impossible arc length)** There is no arc length  $\lambda$  on the whole power set  $\mathcal{P}(S)$ .

**Proof.** Indeed, suppose in the way of contradiction that

$$\lambda: \mathcal{P}(S) \rightarrow [0, +\infty)$$

is an arc length and consider the set  $X \subseteq S$  of the previous theorem. Then we get by the theorem and by the three properties of any arc length that

$$\lambda(S) = \sum_{\varphi \in [0, 2\pi)_{\mathbb{Q}}} \lambda(F_\varphi[X]) = \sum_{\varphi \in [0, 2\pi)_{\mathbb{Q}}} \lambda(X) = 0 \quad \text{or} \quad +\infty.$$

But this is a contradiction because  $\lambda(S) \in (0, +\infty)$ .  $\square$

**Exercise 11** Show that if property 1 of arc length is not required, then the previous corollary does not hold.

• *Well orderings.* Let  $X$  be a set. A relation

$$\leq_X \subseteq X^2$$

is a *linear order* on  $X$  if it is reflexive, transitive, weakly asymmetric ( $a \leq_X b \wedge b \leq_X a \Rightarrow a = b$ ) and total ( $\forall a, b \in X: a \leq_X b \vee b \leq_X a$ ). We say that a linear order  $\leq_X$  on  $X$  is a *well ordering* if every nonempty set  $Y \subseteq X$  has a minimum element  $y \in Y$ : for every  $z \in Y$  we have  $y \leq_X z$ .

**Exercise 12** Prove that minimum elements are unique.

**Exercise 13** A linear order  $(X, \leq_X)$  is a well ordering if and only if there is no infinite strictly descending chain

$$x_1 >_X x_2 >_X \dots, \quad x_n \in X.$$

Here  $x >_X y$  means that  $y \leq_X x$  and  $y \neq x$ .

**Exercise 14** Assume that there is a well ordering on every set and deduce from this the AC.

**Theorem 15 (Zermelo)** The axiom of choice holds if and only if every set has a well ordering.

**Proof.** The “if” part is proven in Exercise 14. We prove the other implication: if AC holds then every set has a well ordering. Let  $X \neq \emptyset$  and  $f: \mathcal{P}(X) \setminus \{\emptyset\} \rightarrow X$  be a *selector on X*, i.e., a function satisfying  $f(A) \in A$  (it is guaranteed by AC). We consider the set

$$L = \{R: R \subseteq D(R)^2, D(R) \subseteq X, R \text{ is a linear order on } D(R)\}$$

of linear orders  $R$  on sets  $D(R) \subseteq X$ . For any  $R \in L$  we set

$$D_R = \{A \subseteq D(R): x, y \in D(R), y \in A, xRy \Rightarrow x \in A\}.$$

So  $D_R$  is the set of *downsets* in the linear order  $R$ . Let further

$$C = \{R \in L: A \in D_R, A \neq D(R) \Rightarrow f(X \setminus A) = \min_R(D(R) \setminus A)\}$$

be those linear orders  $R$  on subsets  $D(R)$  of  $X$ , for which for every proper downset  $A$  in  $R$  the selector  $f$  chooses from its complement to  $X$  an element that is also the minimum element of the complement of  $A$  to  $D(R)$ . We show that  $C$  contains (as an element) a well ordering on  $X$ . The set  $C \neq \emptyset$ , for example  $\{(f(X), f(X))\} \in C$ .

Firstly we show that every  $R \in C$  is a well ordering on  $D(R)$ . Let  $R \in C$ . For any nonempty  $B \subseteq D(R)$  we set

$$A = \{y \in D(R) \setminus B: x \in B \Rightarrow yRx\}.$$

The set  $D(R) \setminus A$  contains  $B$  and is therefore nonempty. Clearly,  $A$  is a downset in  $R$ . Thus

$$y = f(X \setminus A) = \min_R(D(R) \setminus A).$$

From the facts that  $D(R) \setminus A \supset B$  and that  $y$  is the minimum element in  $D(R) \setminus A$  we get that  $yRx$  for every  $x \in B$ . If  $y \notin B$ , we would have  $y \in A$  by the definition of  $A$ , which is impossible. Hence  $y$  is in  $B$  and is the minimum element of  $B$ , even of the superset  $D(R) \setminus A$ .

Secondly we show that for every two linear orders  $R, S \in C$  one of them extends the other:  $D(R) \in D_S \wedge R \subseteq S$  or  $D(S) \in D_R \wedge S \subseteq R$ . Let  $R, S \in C$  be given; we set

$$A = \{x \in D(R) \cap D(S) \mid Rx = Sx \wedge R \cap (Rx \times Rx) = S \cap (Sx \times Sx)\}$$

(here  $Rx = \{y \in D(R) \mid yRx\}$  and similarly for  $Sx$ ). The set  $A$  consists exactly of the elements that determine the same downset in  $R$  and in  $S$ , that is moreover ordered in  $R$  and in  $S$  in the same way. We claim that  $A \in D_R \cap D_S$  —  $A$  is a downset both in  $R$  and in  $S$ ). Let

$$z, y, x \in X \text{ with } x \in A \text{ and } yRx.$$

Then  $ySx$  because  $Rx = Sx$ . If  $zRy$  then  $zSy$  and vice-versa (in both cases  $y, z \in Rx = Sx$  and this set is ordered in the same way in  $R$  and in  $S$ ). Thus  $Ry = Sy$ . This set is contained in  $Rx = Sx$ , and therefore it is ordered in the same way both in  $R$  and in  $S$ . Hence  $y \in A$  and  $A$  is a downset in  $R$ . One shows in the same way that  $A$  is a downset in  $S$ .

Now if both  $D(R) \setminus A$  and  $D(S) \setminus A$  are nonempty,  $y = f(X \setminus A)$  is the minimum element of  $D(R) \setminus A$  with respect to  $R$  and it is also the minimum element of  $D(S) \setminus A$  with respect to  $S$ , and so  $Ry = A \cup \{y\} = Sy$ . It is also clear that  $R$  and  $S$  give  $A \cup \{y\}$  the same order (they add a new element  $y$  at the end), and so  $y \in A$ , which is a contradiction. Thus for example  $A = D(R)$ ,  $R \subseteq S$  and  $S$  extends  $R$ .

Thirdly we show that

$$T = \bigcup C \in C,$$

and therefore  $C$  has (unique) inclusion-wise maximum element. By the previous paragraph,  $T$  is a linear order on  $D(T) = \bigcup_{R \in C} D(R)$  and for  $x, y \in D(T)$  we have  $xTy$ , if and only if  $xRy$  for some  $R \in C$  with  $x, y \in D(R)$ . We check that  $T$  has the property defining  $C$ . Let  $A \subseteq D(T)$  be a proper downset in  $T$  and let  $b \in D(T) \setminus A$  be arbitrary. Thus  $b \in D(R)$  for some  $R \in C$ . We show that  $A \subseteq D(R)$ . If  $a \in A$  is arbitrary, then  $a \in D(S)$  for some  $S \in C$ . If  $D(S) \in D_R$ , then  $a \in D(R)$ . If  $D(R) \in D_S$  and  $aSb$ , then again  $a \in D(R)$ . The case  $bSa$  does not occur (for then one would have  $b \in A$ ). Hence

$A \subseteq D(R)$  and  $D(R) \setminus A \neq \emptyset$ . Therefore the element  $y = f(X \setminus A)$  is the minimum element in  $D(R) \setminus A$  and  $yRb$ . Since  $b$  was arbitrary,  $y$  is the minimum element in  $D(T) \setminus A$  and we see that  $T \in C$ .

In conclusion we show that  $D(T) = X$ , and  $T$  is therefore the sought-for well ordering of  $X$ . If  $D(T) \neq X$ , then we could extend  $T$  by the element  $x := f(X \setminus D(T))$  to  $R$ :

$$D(R) := D(T) \cup \{x\} \quad \text{and} \quad yRx \quad \text{for every } y \in D(R) \quad (1)$$

— we add to  $T$  a new maximum element. It is clear that  $R \in C$  (Exercise 16). Since  $R$  properly extends  $T$ , we have a contradiction with the maximality of  $T$ .  $\square$

The previous proof is taken from a manuscript of A. Pultr.

**Exercise 16** Show that the linear order  $R$  defined in equation (1) indeed belongs to  $C$ .

• *The prophet paradox.* Let  $(X, \leq_X)$  be a linear order. For any  $a \in X$  and any map  $f: X \rightarrow Y$  we denote by  $f|_a$  the restriction of  $f$  to the set

$$\{b \in X: b <_X a\} .$$

For a linear order  $(X, \leq_X)$  and a family  $\mathcal{F}$  of functions  $f: X \rightarrow Y$ , an  $(X, \mathcal{F})$ -prophet is a map

$$P: \{f|_a: f \in \mathcal{F}, a \in X\} \rightarrow Y .$$

The value  $P(f|_a) \in Y$  is the guess of  $P$  for the value  $f(a)$ . The prophet tries to guess from the values  $f(b)$  for all  $b <_X a$  the value of  $f$  at  $a$ . If  $P(f|_a) = f(a)$  then  $P$  succeeds for  $f$  at  $a$ , else  $P$  errs for  $f$  at  $a$ .

**Exercise 17** Let  $(X, \leq_X) = (\mathbb{R}, \leq)$  be the standard linear order of real numbers and let

$$\mathcal{F} = C(\mathbb{R}) = \{f: \mathbb{R} \rightarrow \mathbb{R}: f \text{ is continuous}\}$$

be the set of continuous real functions defined on  $\mathbb{R}$ . The exercise is to find an  $(\mathbb{R}, C(\mathbb{R}))$ -prophet that succeeds for  $f$  at  $a$  for every  $f \in C(\mathbb{R})$  and every  $a \in \mathbb{R}$ .

On the other hand we have the following equally simple result.

**Proposition 18 (all prophets err)** For  $n \in \mathbb{N}$ , let

$$(X, \leq_X) = ([n], \leq) = (\{1, 2, \dots, n\}, \leq)$$

be the usual linear order on the first  $n$  natural numbers and let

$$\mathcal{F} = Y^{[n]} = \{\text{all maps from } [n] \text{ to } Y\},$$

where  $Y$  is a set with at least two elements. Then it is true that for every  $([n], Y^{[n]})$ -prophet  $P$  there exists a function  $f \in Y^{[n]}$  such that

$$\forall a \in [n] : P(f|_a) \neq f(a).$$

Thus  $P$  errs for  $f$  at its every argument  $a \in [n]$ .

**Proof.** Let

$$P : \{g : g : [m] \rightarrow Y, m \in \{0, 1, \dots, n-1\}\} \rightarrow Y$$

be an  $([n], Y^{[n]})$ -prophet. We set  $[0] = \emptyset$ . We define the values  $f(m)$  of the required function  $f : [n] \rightarrow Y$  by induction on  $m = 1, 2, \dots, n$ . At the start we take  $f(1) \in Y$  so that  $f(1) \neq P(\emptyset)$ , which is possible as  $|Y| \geq 2$ . If  $m \in [n]$ ,  $m > 1$  and  $f(1), f(2), \dots, f(m-1)$  are already defined, we take

$$f(m) \in Y \setminus \{P(f|_m)\}.$$

Again, this is possible as  $|Y| \geq 2$ . It is clear that  $P$  errs for the function  $f$  at its every argument.  $\square$

**Exercise 19** What happens when  $|Y| \leq 1$ ?

One might think that when in Exercise 17 the family of continuous functions is extended to the family  $\mathcal{F} = \mathbb{R}^{\mathbb{R}}$  of all real function, one obtains a result similar to the previous proposition, namely that every prophet has to err for a troublesome function very often. Surprisingly, quite the opposite is the case under the assumption of AC. There exists a prophet that for every real function almost never errs.

**Theorem 20 (the prophet paradox)** Let  $(X, \leq_X) = (\mathbb{R}, \leq)$  be the usual linear order of real numbers and let

$$\mathcal{F} = \mathbb{R}^{\mathbb{R}} = \{\text{all functions from } \mathbb{R} \text{ to } \mathbb{R}\}.$$

Then there exists an  $(\mathbb{R}, \mathbb{R}^{\mathbb{R}})$ -prophet  $P$  such that

$\forall f \in \mathbb{R}^{\mathbb{R}} : \text{the set } \{a \in \mathbb{R} : P \text{ errs for } f \text{ at } a\} \text{ is at most countable.}$

**Proof.** We define  $P$  by means of the well ordering

$$(\mathbb{R}^{\mathbb{R}}, \preceq)$$

that exists by Theorem 15 under the assumption of AC. For  $g \in \mathbb{R}^{\mathbb{R}}$  and  $a \in \mathbb{R}$  we set

$$P(g|_a) = g_0(a) \text{ where } g_0 = \min_{\preceq}(\{h \in \mathbb{R}^{\mathbb{R}} : h|_a = g|_a\}).$$

Now let an  $f \in \mathbb{R}^{\mathbb{R}}$  be given. We take the set

$$X = \{a \in \mathbb{R} : P(f|_a) \neq f(a)\}$$

of errors of  $P$  for  $f$ . Let  $a < b$  with  $a \in X$  be two real numbers,

$$g_a = \min_{\preceq}(\underbrace{\{g \in \mathbb{R}^{\mathbb{R}} : g|_a = f|_a\}}_{M_a}) \text{ and } g_b := \min_{\preceq}(\underbrace{\{g \in \mathbb{R}^{\mathbb{R}} : g|_b = f|_b\}}_{M_b}).$$

From  $a < b$  we get that  $M_b \subseteq M_a$  and  $g_a \preceq g_b$ . From

$$g_a(a) = P(f|_a) \neq f(a) = g_b(a)$$

we see that  $g_a \neq g_b$ . Thus  $g_a \prec g_b$ . We see that the linear order  $(X, \leq)$  (with the usual order  $\leq$  of real numbers) is a well ordering. Else, by Exercise 13, we would have in  $(X, \leq)$  an infinite strictly descending chain  $a_1 > a_2 > \dots$ , which would yield by the last argument an infinite strictly descending chain  $g_{a_1} \succ g_{a_2} \succ \dots$  in  $(\mathbb{R}^{\mathbb{R}}, \preceq)$ . But the last chain does not exist because  $(\mathbb{R}^{\mathbb{R}}, \preceq)$  is a well ordering. Since  $(X, \leq)$  is a well ordering, by the next Exercise 21 the set  $X$  is at most countable.  $\square$

**Exercise 21** Let  $(\mathbb{R}, \leq)$  be the usual linear order of real numbers and let  $X \subseteq \mathbb{R}$  be such that the linear suborder  $(X, \leq)$  is a well ordering. Show that then  $X$  is at most countable.

The last theorem is taken from the book

Ch. S. Hardin and A. D. Taylor, *The Mathematics of Coordinated Inference*, Springer, 2013.

THANK YOU FOR YOUR ATTENTION!

HOMEWORK: Exercises 6, 7, 13 and 21. Deadline is the end of the coming Monday. Please, send me your solutions by e-mail to [klazar@kam.mff.cuni.cz](mailto:klazar@kam.mff.cuni.cz). To get credits for the tutorial, you should solve (or at least send in attempted solutions of) at least half of the homework exercises.