On ordered Ramsey numbers of bounded-degree graphs

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Ramsey's theorem for graphs

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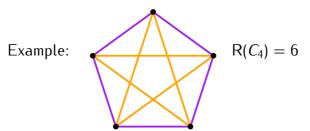
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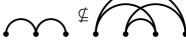
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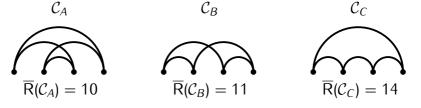


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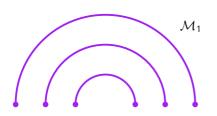
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Conlon et al. showed that this holds for almost every ordered matching.

• There are *n*-vertex ordered matchings \mathcal{M} with $\overline{\mathbb{R}}(\mathcal{M})$ linear in n.

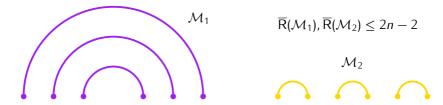
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$$\overline{R}(\mathcal{M}_1)$$
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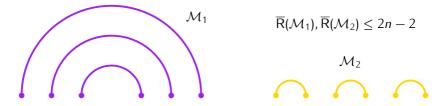


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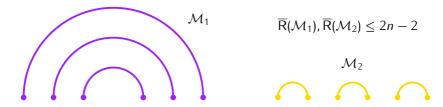
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Problem (Conlon, Fox, Lee, Sudakov, 2014)

Do random 3-regular graphs have superlinear ordered Ramsey numbers for all orderings?

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• For 3-regular graphs, we obtain min- $\overline{R}(G) \ge \frac{n^{7/6}}{4 \log n \log \log n}$.

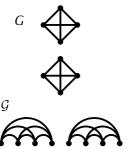
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- Key lemma: Almost every such G satisfies the following: for every partition of V(G) into few sets X_1, \ldots, X_t , each of size at most s, there are many pairs (X_i, X_j) with an edge between them. Here, $t = n/(2 \log n \log \log n)$ and $s = n^{1/2-1/d}/2$.

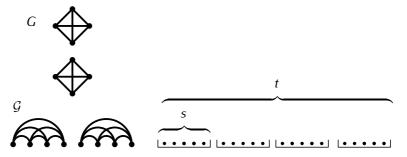
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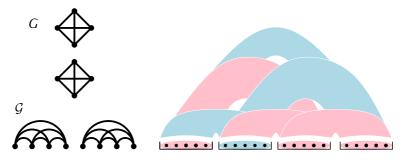
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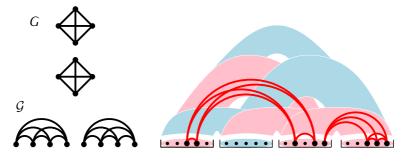
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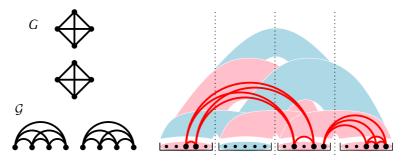
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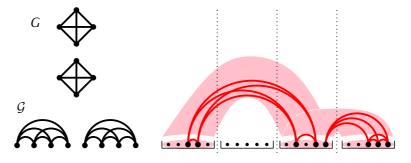
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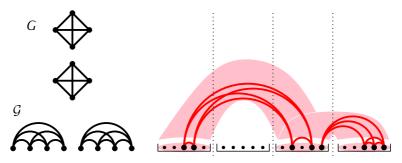
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 We use an estimate by Bender and Canfield and by Wormald for the number of d-regular graphs on n vertices.

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- How about 2-regular graphs?

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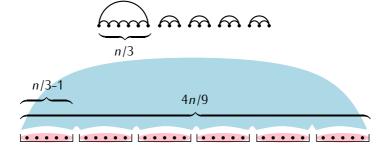
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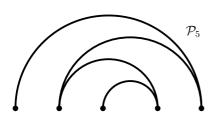
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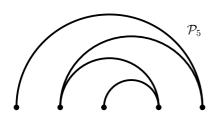


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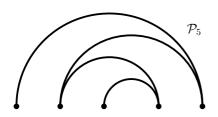


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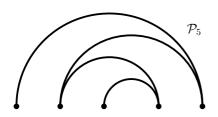


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For every $\varepsilon > 0$ and every $n \in \mathbb{N}$, every ordered graph on $N \geq n/\varepsilon$ vertices with at least εN^2 edges contains \mathcal{P}_n as an ordered subgraph.

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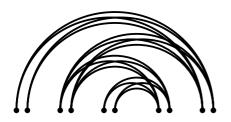


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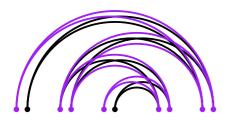


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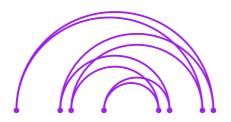


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