

Recent Progress on Hill's Conjecture

Martin Balko, Radoslav Fulek and Jan Kynčl

Charles University in Prague,
Czech Republic

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Preliminaries – Drawings

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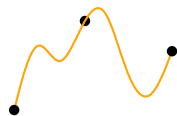
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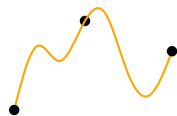
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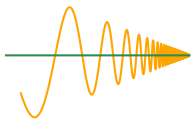
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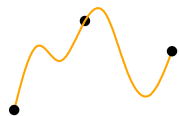
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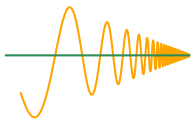
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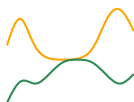
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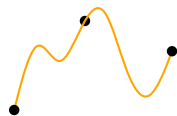
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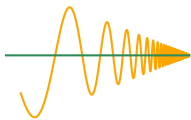
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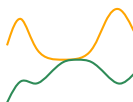
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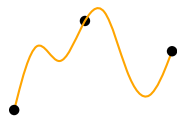
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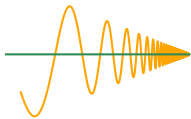
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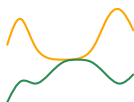
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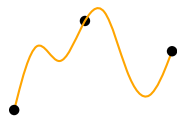


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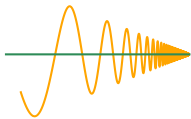
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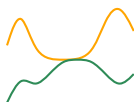
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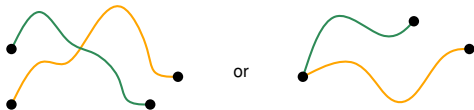


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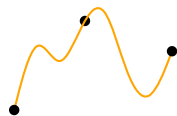
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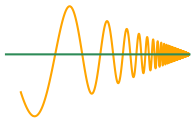
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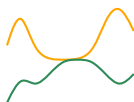
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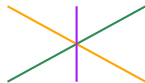
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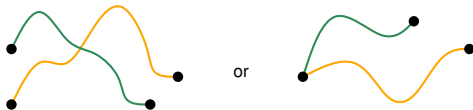


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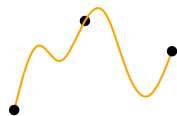


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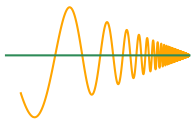
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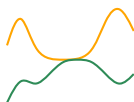
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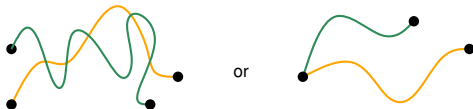


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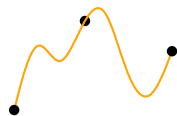
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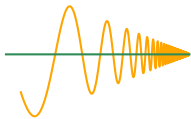
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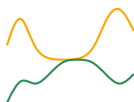
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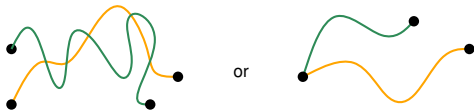


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- A drawing is called **x-monotone** if edges are x-monotone curves.

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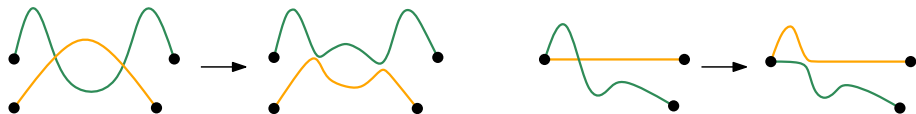


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Conjecture (Hill, 1958)

We have $\text{cr}(K_n) = Z(n) := \frac{1}{4} \lfloor \frac{n}{2} \rfloor \lfloor \frac{n-1}{2} \rfloor \lfloor \frac{n-2}{2} \rfloor \lfloor \frac{n-3}{2} \rfloor$ for every $n \in \mathbb{N}$.

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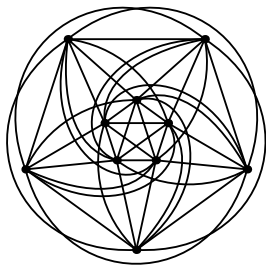
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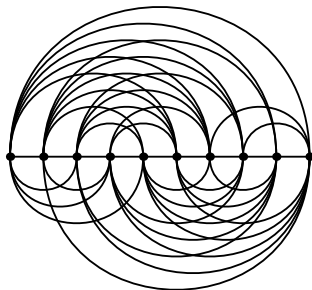
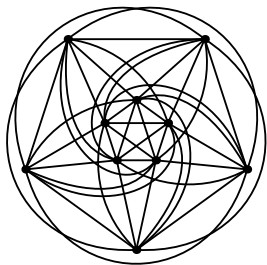


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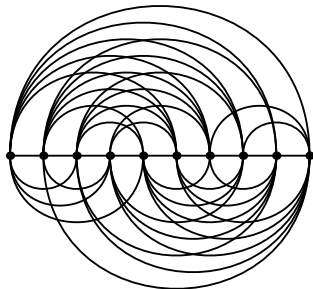
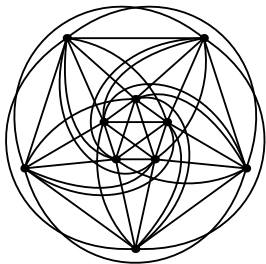


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- A drawing is **2-page** if the vertices are placed on a line ℓ and each edge is fully contained in a halfspace determined by ℓ .

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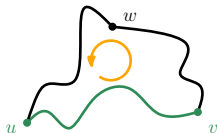
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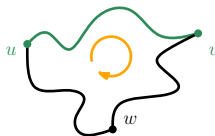
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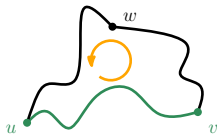
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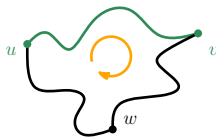
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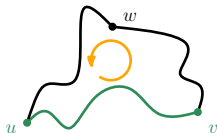


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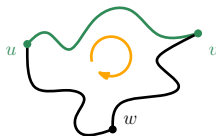
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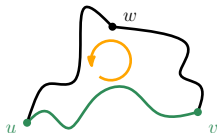


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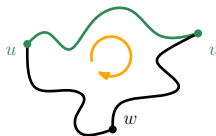
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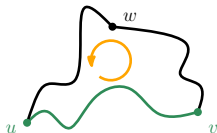


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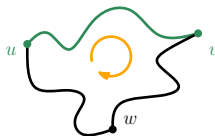
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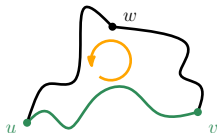
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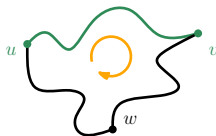


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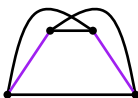


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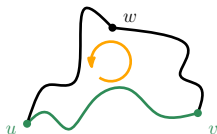
- A k -edge is an edge that has exactly k vertices on the same side.
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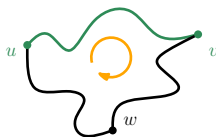
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Sketch of the Proof: Double Counting

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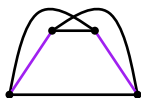


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Lemma

For a simple drawing D of K_n we get $\text{cr}(D) = 3\binom{n}{4} - \sum_{k=0}^{\lfloor n/2 \rfloor - 1} k(n-2-k)E_k(D)$.

Sketch of the Proof: Main Trick

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Lemma

For every simple drawing D of K_n we have

$$\text{cr}(D) = 2 \sum_{k=0}^{\lfloor n/2 \rfloor - 2} E_{\leq k}(D) - \frac{1}{2} \binom{n}{2} \left\lfloor \frac{n-2}{2} \right\rfloor - \frac{1}{2} (1 + (-1)^n) E_{\leq \lfloor n/2 \rfloor - 2}(D).$$

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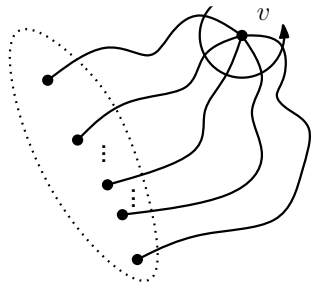
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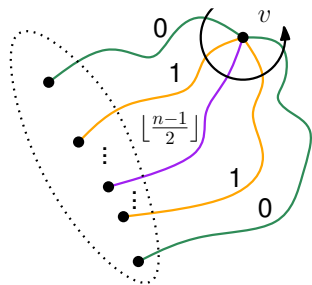
Sketch of the Proof: Structure of k -edges

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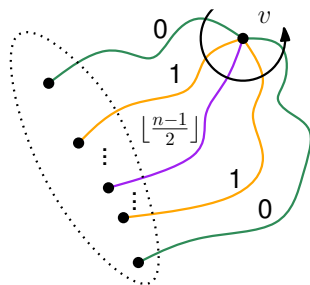
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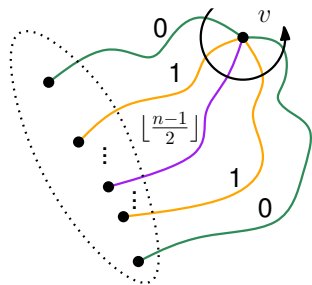
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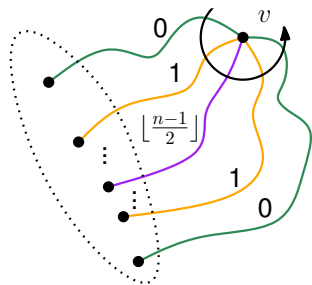
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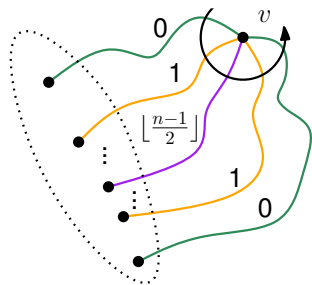
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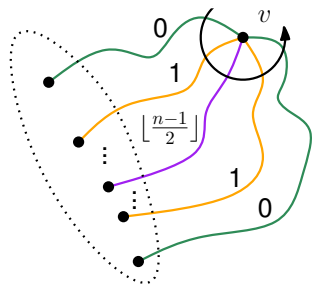
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For a simple x -monotone D we have $E_{\leq k}(D, D') \geq \sum_{i=1}^{k+1} (k+2-i) = \binom{k+2}{2}$.

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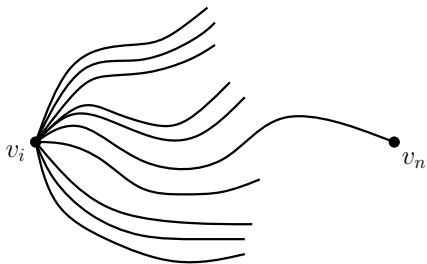
- For $0 \leq k \leq (n-3)/2$ and every $i \in [k+1]$, the $k+2-i$ bottommost and $k+2-i$ topmost right edges at v_i are j -edges for some $j \leq k$.

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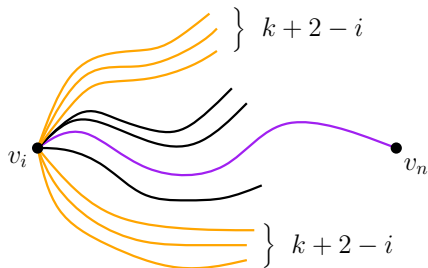
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Let $n \geq 3$ and let D be a simple x -monotone drawing of K_n . Then for every k , $0 \leq k < n/2 - 1$, we have $E_{\leq k}(D) \geq 3 \binom{k+3}{3}$.

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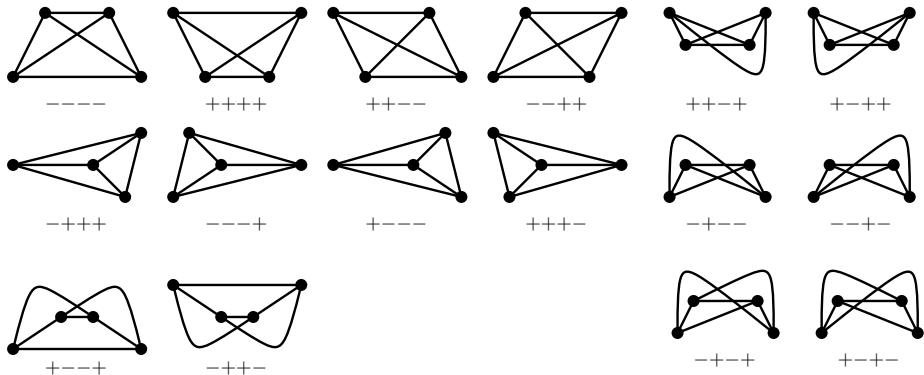
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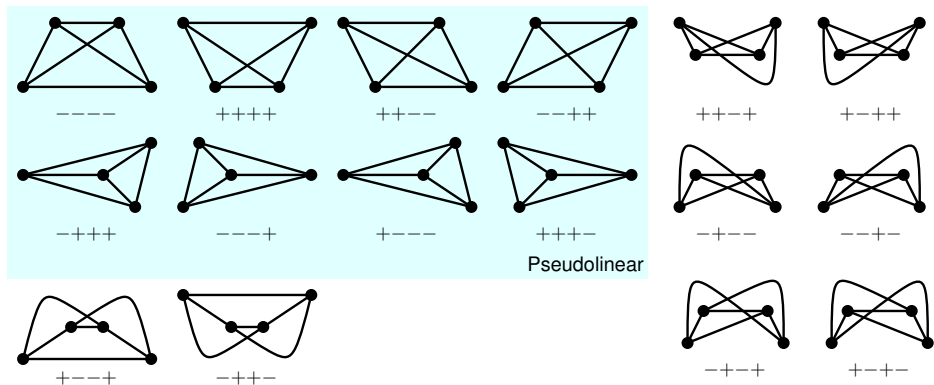
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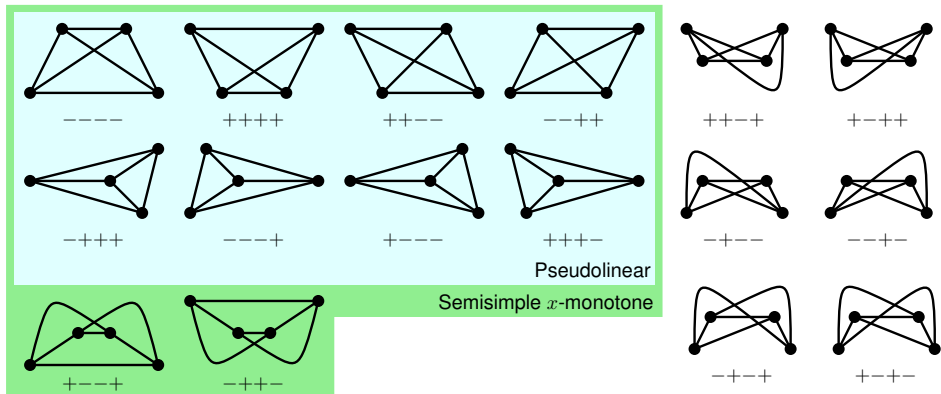
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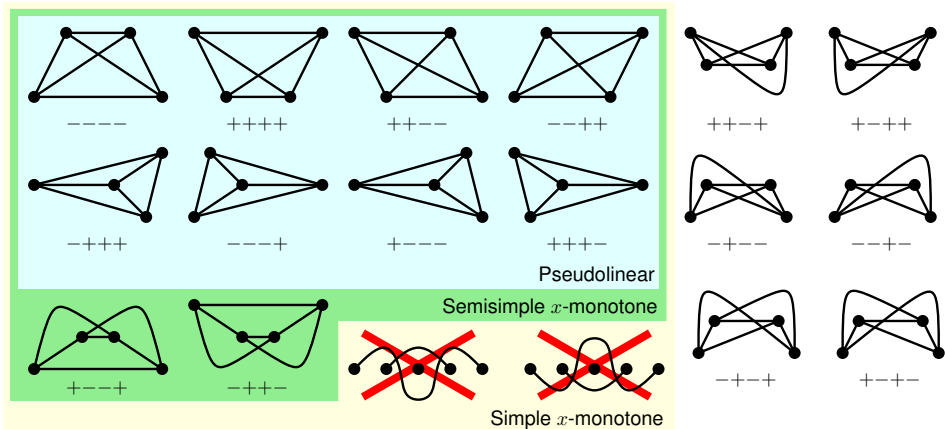
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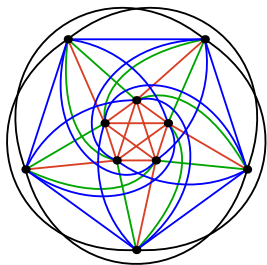
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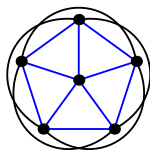
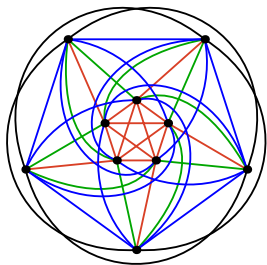
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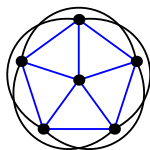
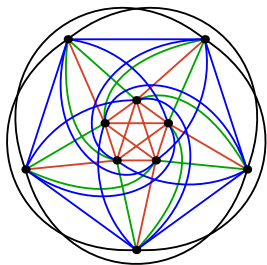
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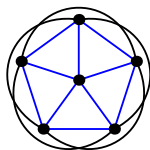
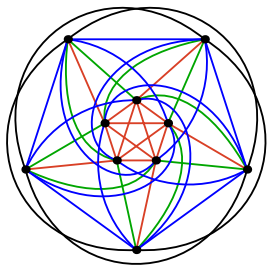
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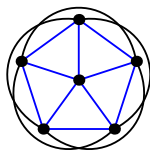
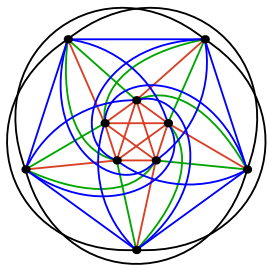
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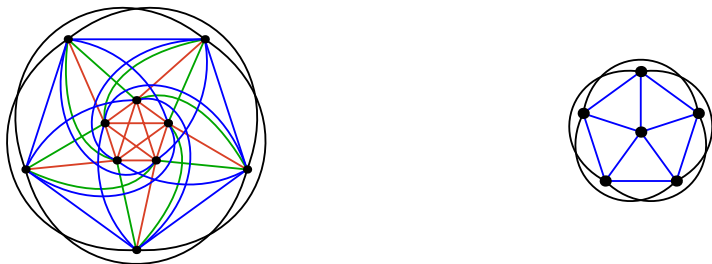
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- Implies Hill's conjecture. All drawings we have found satisfy this conjecture.