

(More about tree decomp.)

1) algo — given G, k — ask $tw(G) \leq k$? ←
 — NP-complete (Arnborg, Corneil, Proskurowski 1997)

$\forall k$ ex. a linear-time algo to det. tw
 (and find a tree dec!)
 $k \cdot n$ (better — approx.)

Open Problem: P for planar graphs

2) $\mathcal{G}_k := \{G : tw(G) \leq k\}$ — closed on minors ($H \leq_m G \in \mathcal{G}_k \Rightarrow H \in \mathcal{G}_k$)
 MCC — minor closed class
 → every MCC is $\text{Forb}_{\leq m}(F)$ for some F

Graph Minor Theorem ⇒ we can take F finite

GMT
 \leq_m is WQO on finite graphs

- \mathcal{G}_0 = edge-less graph = $\text{Forb}_{\leq m}(K_2)$
- \mathcal{G}_1 = forests = $\text{Forb}(K_3)$
- \mathcal{G}_2 = seq.-para. = $\text{Forb}(K_4)$
- \mathcal{G}_3 : $\text{Forb}(K_5, C_5 \square K_2, W_8, K_{2,2,2})$
 (octahedron)
- \mathcal{G}_4 — — — — — > 70 graphs

Def WQO — well-quasi-ordering: relation \leq — transitive & reflexive
 $\exists a_1, a_2, a_3 \dots \{ \exists i < j : a_i \leq a_j \}$ good param. (quasi-order)
 \Leftrightarrow no infinite antichain & no inf. dec. seq.
 (≡ ranked)

(3) Look for nice descr. of $\text{Forb}_{\leq n}(H)$ for arbitrary H

Thm (Robertson, Seymour, 2003)

$\forall n$... structural description of $\text{Forb}_{\leq n}(K_n)$ using tw

(4) $\text{tw}(\text{Grid}_n) = n \quad (\geq n-1, \leq n)$



$G \geq_n H \Rightarrow \text{tw}(G) \geq \text{tw}(H)$

Grid_n is planar

\nexists $K_5, K_{3,3}$ (Kuratowski)

$\text{Forb}_{\leq n}(H)$ can have unbdd tw

Thm (Grid theorem) $\forall r \exists k \forall G: \text{tw}(G) \geq k \Rightarrow G \geq_n \text{Grid}_r$

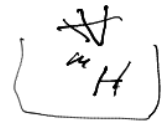
Coroll. tw is bounded on $\text{Forb}_{\leq n}(H)$ $\iff H$ is planar

Proof \Rightarrow , \Leftarrow : $\text{Grid}_n \in \text{Forb}_{\leq n}(H)$ (because minor of planar gr. is planar)
 $(H \text{ not planar}) \quad \text{tw}(\text{Grid}_n) \geq n$

\Leftarrow : $\forall k \exists G_k \in \text{Forb}_{\leq n}(H): \text{tw}(G_k) \geq k$

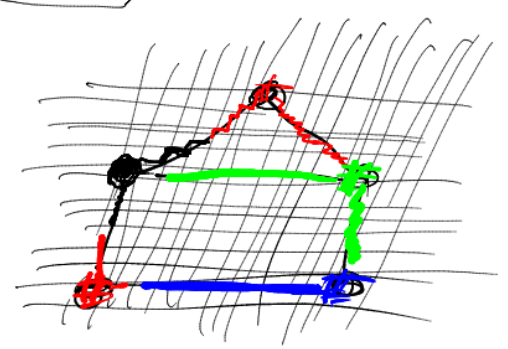
$\forall r \exists k \forall G: \text{tw}(G) \geq k \Rightarrow G \geq_n \text{Grid}_r$

$\forall r \exists k \quad G_k \geq_n \text{Grid}_r \geq_n H$



Assume H is planar

$\exists r \exists G \in \text{Forb}_{\leq n}(H) \not\geq_n \text{Grid}_r$



Szemerédi's regularity lemma

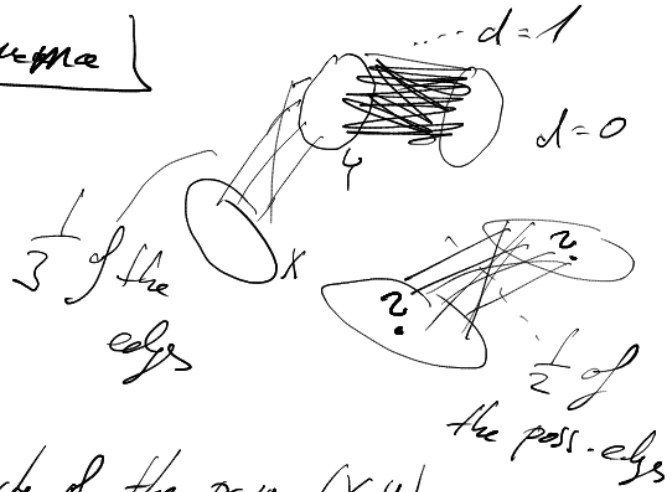
$G = (V, E)$... simple graph

$X, Y \subseteq V$ $X \cap Y = \emptyset$

$\|X, Y\| = \#$ of X - Y edges in G

$$d(X, Y) = \frac{\|X, Y\|}{|X| \cdot |Y|}$$

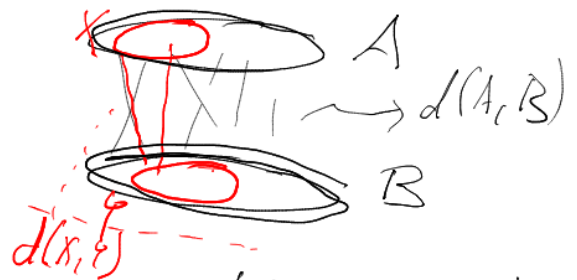
density of the pair (X, Y)
 $\in [0, 1]$



Def $\epsilon > 0$, $A, B \subseteq V$, A, B disjoint

We say (A, B) is an ϵ -regular pair

$\forall X \subseteq A \forall Y \subseteq B$: $|X| \geq \epsilon |A|, |Y| \geq \epsilon |B| \Rightarrow |d(X, Y) - d(A, B)| \leq \epsilon$



$V = V_0 \cup V_1 \cup \dots \cup V_k$ (disjoint sets) are ϵ -regular partitions

1) $|V_0| \leq \epsilon |V|$

2) $|V_1| = \dots = |V_k|$

3) all but $\leq \epsilon k^2$ of the pairs (V_i, V_j) (with $1 \leq i < j \leq k$)

are ϵ -regular



Theorem (Regularity Lemma)

$\forall \epsilon > 0 \forall m \in \mathbb{N} \exists M \in \mathbb{N} \forall G : |G| \geq m \Rightarrow \exists \epsilon$ -regular part. of G into k parts with $m \leq k \leq M$.

\rightarrow part. into single vertices ... obv. ϵ -reg. ... need for M

\rightarrow part. with $k=1$ $V_1 = V$... need for m

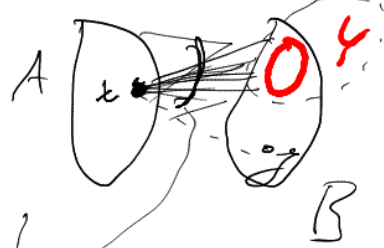
\rightarrow triv. true for sparse graphs $\frac{|E|}{|V|}$ small



Lemma (A, B) is an ϵ -regular pair of density d .

(in some G)

$Y \subseteq B, |Y| \geq \epsilon |B|$



Then all but $< \epsilon |A|$ vertices in A have each $\geq (d - \epsilon) |Y|$ neighbours in Y .

Take random bip. graph A, B & prob. of each edge = p
 $d = p$ w.h.p.

Proof $X = \{x \in A : |\{x, y\}, Y| < (d - \epsilon) \cdot |Y|\}$

If $|X| < \epsilon |A|$, we are done.

Otherwise: $|X| \geq \epsilon |A|, |Y| \geq \epsilon |B| \implies |d(x, Y) - d(A, B)| \leq \epsilon$

$$\frac{|X, Y|}{|X| \cdot |Y|} = \frac{\sum_{x \in X} |\{x, y\}, Y|}{|X| \cdot |Y|} < \frac{(d - \epsilon) \cdot |Y| \cdot |X|}{|X| \cdot |Y|} = \underline{d - \epsilon}$$

Given G & its ϵ -regular part. V_1, V_2, \dots, V_k ; $d \in [0, 1]$

we define R ... regularity graphs of G with ϵ, k, d

$|V_1| = \dots = |V_k|$

vertices $\{V_1, \dots, V_k\}$

edges $V_i, V_j \dots (V_i, V_j)$ is ϵ -regular

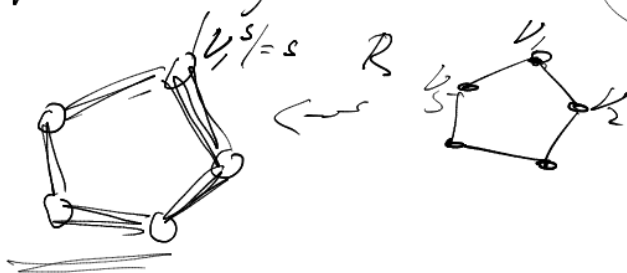
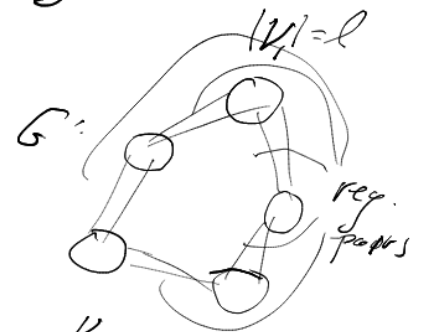
& $d(V_i, V_j) \geq d$.

replace

$V_i \rightsquigarrow \underline{V_i^s}$... set of s vertices

$V_i, V_j \rightsquigarrow$ complete bip. graph V_i^s to V_j^s

R_s ... resulting graph



if $R = K_k$

then R_s is complete

k -part. graph.

(Blow-up Lemma; Kozlov-Sarkozy, Szemerédi, 1997)

Lemma $\forall d \in (0, 1] \forall \Delta \geq 1 \exists \varepsilon_0 > 0 :$

$\forall G$ graph $\forall H$ graph $\Delta(H) \leq \Delta \quad \forall s \in \mathbb{N}$

if R ... regul. graph of G with $\varepsilon \leq \varepsilon_0, d \geq \frac{2s}{\Delta}, d$
then

$$\underline{H \subseteq R_s} \Rightarrow \underline{H \subseteq G}$$

Turán thm: many edges \Rightarrow compl. subgr. of G

Kodó's-stone $\underline{\text{---}} \Rightarrow \underline{H \subseteq G}$

H bipartite ... separate problem