Crossings in Topological Graphs

Jan Kynčl supervisor: Pavel Valtr

Graph: $G = (V, E), |V| < \infty, E \subseteq {V \choose 2}$

Graph:
$$G = (V, E), |V| < \infty, E \subseteq {V \choose 2}$$

Topological graph: drawing of an (abstract) graph in the plane

vertices = points

Graph:
$$G = (V, E), |V| < \infty, E \subseteq {V \choose 2}$$

Topological graph: drawing of an (abstract) graph in the plane

vertices = points

edges = simple curves

 edges do not pass through any vertices other than their end-points Graph: $G = (V, E), |V| < \infty, E \subseteq {V \choose 2}$

Topological graph: drawing of an (abstract) graph in the plane

vertices = points

- edges do not pass through any vertices other than their end-points
- any two edges have only finitely many common points

Graph:
$$G = (V, E), |V| < \infty, E \subseteq {V \choose 2}$$

Topological graph: drawing of an (abstract) graph in the plane

vertices = points

- edges do not pass through any vertices other than their end-points
- any two edges have only finitely many common points
- any intersection point of two edges is either a common end-point or a crossing (no touching allowed)

Graph:
$$G = (V, E), |V| < \infty, E \subseteq {V \choose 2}$$

Topological graph: drawing of an (abstract) graph in the plane

vertices = points

- edges do not pass through any vertices other than their end-points
- any two edges have only finitely many common points
- any intersection point of two edges is either a common end-point or a crossing (no touching allowed)
- at most two edges can intersect in one crossing

simple: any two edges have at most one common point complete: $E = {V \choose 2}$

simple: any two edges have at most one common point complete: $E = {V \choose 2}$



a simple complete topological graph

Problem: What is the smallest integer h(n), such that every simple complete topological graph on n vertices contains an edge crossing at most h(n) other edges?

Problem: What is the smallest integer h(n), such that every simple complete topological graph on n vertices contains an edge crossing at most h(n) other edges?

• An average number of crossings on a single edge is $\Theta(n^2)$ (since $\operatorname{cr}(K_n) = \Theta(n^4)$)

Problem: What is the smallest integer h(n), such that every simple complete topological graph on n vertices contains an edge crossing at most h(n) other edges?

- An average number of crossings on a single edge is $\Theta(n^2)$ (since $\operatorname{cr}(K_n) = \Theta(n^4)$)
- For geometric and extendable graphs h(n) = 0



Problem: What is the smallest integer h(n), such that every simple complete topological graph on n vertices contains an edge crossing at most h(n) other edges?

- An average number of crossings on a single edge is $\Theta(n^2)$ (since $\operatorname{cr}(K_n) = \Theta(n^4)$)
- For geometric and extendable graphs h(n) = 0



• h(n) = 0 for $n \le 7$, h(n) > 0 for $n \ge 8$ [H. Harborth, I. Mengersen, 1974] Previously known bounds for $\boldsymbol{h}(\boldsymbol{n})$:

Previously known bounds for h(n): Lower bound: $h(n) \ge \left(\frac{3}{4} + o(1)\right)n$ [H. Harborth, C.Thürmann, 1994] Previously known bounds for h(n): Lower bound: $h(n) \ge \left(\frac{3}{4} + o(1)\right)n$ [H. Harborth, C.Thürmann, 1994] $h(n) \ge \Omega(n^{3/2})$ [P. Valtr, 2005] Previously known bounds for h(n): Lower bound: $h(n) \ge \left(\frac{3}{4} + o(1)\right) n$ [H. Harborth, C.Thürmann, 1994] $h(n) \ge \Omega(n^{3/2})$ [P. Valtr, 2005] Upper bound: $h(n) = O(n^2)$ (trivial) Previously known bounds for h(n): Lower bound: $h(n) \ge \left(\frac{3}{4} + o(1)\right)n$ [H. Harborth, C.Thürmann, 1994] $h(n) \ge \Omega(n^{3/2})$ [P. Valtr, 2005] Upper bound: $h(n) = O(n^2)$ (trivial) Conjecture: $h(n) = o(n^2)$ [P. Brass, W. Moser, J. Pach, 2005]

Theorem:

$$h(n) \le O\left(\frac{n^2}{\log^{1/4} n}\right)$$

Theorem:

$$h(n) \le O\left(\frac{n^2}{\log^{1/4} n}\right)$$

Theorem: Every simple complete topological graph on n vertices with an induced extendable subgraph on cn vertices contains an edge with $O(n^{3/2} \log n)$ crossings.

Theorem:

$$h(n) \le O\left(\frac{n^2}{\log^{1/4} n}\right)$$

Theorem: Every simple complete topological graph on n vertices with an induced extendable subgraph on cn vertices contains an edge with $O(n^{3/2} \log n)$ crossings.

Theorem: There exists a complete topological graph on n vertices, such that

- edges with common end-point do not cross
- every two edges have at most two common points
- every edge crosses $\Omega(n^2)$ other edges.

Theorem: On every connected compact surface except for the sphere there exists a simple complete topological graph with each edge having at least $\Omega(n^2)$ crossings.

Theorem: On every connected compact surface except for the sphere there exists a simple complete topological graph with each edge having at least $\Omega(n^2)$ crossings.

torus:



projective plane: [A. Pór]



Recognition of simply realizable complete abstract topological graphs

Abstract topological graph (AT-graph):

A = (G, R); G = (V, E) is a graph, $R \subseteq {E \choose 2}$

Recognition of simply realizable complete abstract topological graphs

Abstract topological graph (AT-graph):

A = (G, R); G = (V, E) is a graph, $R \subseteq {E \choose 2}$

A is simply realizable if there exists a simple topological graph T, which is a drawing of G, and exactly the pairs of edges from R cross in T

Recognition of simply realizable complete abstract topological graphs

Abstract topological graph (AT-graph):

A = (G, R); G = (V, E) is a graph, $R \subseteq {E \choose 2}$

A is simply realizable if there exists a simple topological graph T, which is a drawing of G, and exactly the pairs of edges from R cross in T

Theorem: The problem, whether a given complete AT-graph is simply realizable, is solvable in polynomial time.

Recognition of simply realizable complete abstract topological graphs

Abstract topological graph (AT-graph):

A = (G, R); G = (V, E) is a graph, $R \subseteq {E \choose 2}$

A is simply realizable if there exists a simple topological graph T, which is a drawing of G, and exactly the pairs of edges from R cross in T

Theorem: The problem, whether a given complete AT-graph is simply realizable, is solvable in polynomial time.

Remark: Other similar problems are NP-hard

[J. Kratochvíl, 1991; J.K., 2006]

Topological graphs G, H are

isomorphic if there exists a homeomorphism (of the sphere) which maps G onto H

Topological graphs G, H are

isomorphic if there exists a homeomorphism (of the sphere) which maps G onto H

Topological graphs G, H are

isomorphic if there exists a homeomorphism (of the sphere) which maps G onto H



Topological graphs G, H are

isomorphic if there exists a homeomorphism (of the sphere) which maps G onto H



Topological graphs G, H are

isomorphic if there exists a homeomorphism (of the sphere) which maps G onto H



T(n) = number of non-isomorphic $T_w(n) =$ number of weakly non-isomorphic simple complete topological graphs on n vertices T(n) = number of non-isomorphic $T_w(n) =$ number of weakly non-isomorphic simple complete topological graphs on n vertices **Theorem:**

$$2^{\Omega(n^2)} \le T_w(n) \le 2^{O(n^2 \log n)}$$

[J. Pach, G. Tóth, 2004]

$$T(n) = 2^{\Theta(n^4)}$$

Lower bounds are attained even for extendable graphs

T(n) = number of non-isomorphic $T_w(n) =$ number of weakly non-isomorphic simple complete topological graphs on n vertices **Theorem:**

$$2^{\Omega(n^2)} \le T_w(n) \le 2^{O(n^2 \log n)}$$

[J. Pach, G. Tóth, 2004]

$$T(n) = 2^{\Theta(n^4)}$$

Lower bounds are attained even for extendable graphs

Remark: The number of weakly non-isomorphic complete geometric graphs on *n* vertices is $2^{O(n \log n)}$