## NON-P-RECURSIVENESS OF NUMBERS OF MATCHINGS OR LINEAR CHORD DIAGRAMS WITH MANY CROSSINGS

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ABSTRACT. The number  $\operatorname{con}_n$  counts matchings X on  $\{1, 2, \ldots, 2n\}$ , which are partitions into n two-element blocks, such that the crossing graph of X is connected. Similarly,  $\operatorname{cro}_n$  counts matchings whose crossing graph has no isolated vertex. (If it has no edge, Catalan numbers arise.) We apply generating functions techniques and prove, using a more generally applicable criterion, that the sequences  $(\operatorname{con}_n)$  and  $(\operatorname{cro}_n)$  are not P-recursive. On the other hand, we show that the residues of  $\operatorname{con}_n$  and  $\operatorname{cro}_n$  modulo any fixed power of 2 can be determined P-recursively. We consider also the numbers  $\operatorname{sco}_n$  of symmetric connected matchings. Unfortunately, their generating function satisfies a complicated differential equation which we cannot handle.

Matchings on the vertex set  $[2n] = \{1, 2, ..., 2n\}$  consist of n mutually disjoint two-element edges. One finds easily that their number mat<sub>n</sub> equals  $(2n-1)!! = 1 \cdot 3 \cdot 5 \cdot ... \cdot (2n-1)$ . Another classical result tells us that the number ncr<sub>n</sub> of noncrossing matchings on [2n] (no two edges  $\{a, b\}$  and  $\{c, d\}$  satisfy a < c < b < d) is the *n*th Catalan number:  $\operatorname{ncr}_n = \frac{1}{n+1} \binom{2n}{n}$ . How many matchings are there if their crossings are restricted in a more complicated way? In the present article we investigate numbers of such matchings, namely the numbers  $\operatorname{connected} matchings$  in which each two edges can be connected by a chain of consecutively crossing edges, the numbers  $\operatorname{sco}_n$  of symmetric connected matchings which, in addition, are mirror symmetric, and the numbers  $\operatorname{con}_n$  of crossing matchings in which each edge crosses another edge. We concentrate only on P-recursiveness of these numbers. Also, we touch upon some modular properties. The sequences  $(\operatorname{mat}_n)$  and  $(\operatorname{ncr}_n)$  are trivially P-recursive but, as we prove, the sequences  $(\operatorname{con}_n)$  and  $(\operatorname{cro}_n)$  are not.

First we remind the definition of P-recursiveness and D-finiteness. Then we introduce  $D_A$ -finiteness and review some facts on power series. In Theorem 1 we prove that if a sequence of numbers has an OGF (ordinary generating function) that has zero convergence radius and satisfies a certain differential equation, then the sequence is far from being P-recursive. In Theorem 2 we apply this criterion to the sequences  $(con_n)$  and  $(cro_n)$ . In Theorem 3 a more complicated differential equation is derived for the OGF of the sequence  $(sco_n)$ . Finally, in Theorem 4 we show that modulo  $2^l$  the sequences  $(con_n)$  and  $(cro_n)$  coincide with certain P-recursive, in fact algebraic, sequences.

Symbols  $\mathbb{Z}$  and  $\mathbb{N}$  denote the sets of integers  $\{\ldots, -1, 0, 1, \ldots\}$  and  $\{1, 2, \ldots\}$ .  $\mathbb{C}$  denotes the set of complex numbers. A sequence of complex numbers  $(a_n)_{n\geq 0}$  is called *P*-recursive if there exist polynomials  $P_0, P_1, \ldots, P_j \in \mathbb{C}[x], P_0 \neq 0$ , such that

$$P_0(n)a_n + P_1(n)a_{n-1} + \dots + P_j(n)a_{n-j} = 0$$

holds for each integer  $n, n \ge j$ . Many combinatorial counting sequences are P-recursive, for instance  $(mat_n)$ ,  $(ncr_n)$ , Schröder numbers, and numbers of labelled k-regular graphs (Gessel [7]). But some are not, for instance Bell numbers and numbers of integer partitions.

We write  $\mathbb{C}[[x_1, \ldots, x_k]]$  for the ring of power series with complex coefficients and variables  $x_1, \ldots, x_k$ . A power series  $F \in \mathbb{C}[[x]]$  is *D*-finite if F solves the linear differential equation

$$R_0 F^{(m)} + R_1 F^{(m-1)} + \dots + R_{m-1} F' + R_m F + R_{m+1} = 0$$

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with polynomial coefficients  $R_i \in \mathbb{C}[x]$ ,  $R_0 \neq 0$ . A sequence is P-recursive if and only if its OGF is D-finite. For the (easy) proof of this equivalence and further information and references on P-recursiveness and D-finiteness we refer the reader to Stanley [22, chapter 6]. At first these concepts were systematicly investigated in Stanley [21].

The ring  $\mathbb{C}[[x_1, \ldots, x_k]]$  is contained in the ring of Laurent series  $\mathbb{C}((x_1, \ldots, x_k))$  whose elements are formal sums

$$F = \sum_{i_1,\dots,i_k \in \mathbb{Z}} a_{i_1,\dots,i_k} x_1^{i_1} \dots x_k^{i_k}$$

with  $a_{i_1,\ldots,i_k} \in \mathbb{C}$  and such that, for some  $p \in \mathbb{N}$  (depending on F),  $(x_1 \ldots x_k)^p F \in \mathbb{C}[[x_1, \ldots, x_k]]$ . For k = 1 the ring  $\mathbb{C}((x_1, \ldots, x_k))$  is a field (but not for k > 1). We use notation

$$[x_1^{i_1} \dots x_k^{i_k}]F = a_{i_1,\dots,i_k} \text{ and} \text{ ord}_{x_j}(F) = \min\{i_j \in \mathbb{Z} : a_{i_1,\dots,i_k} \neq 0 \text{ for some } i_1,\dots,i_{j-1},i_{j+1},\dots,i_k \in \mathbb{Z}\}.$$

By the definition, for every j and nonzero  $F \in \mathbb{C}((x_1, \ldots, x_k))$  we have  $\operatorname{ord}_{x_j}(F) > -\infty$ . We set  $\operatorname{ord}_{x_j}(0) = \infty$ .

A power series  $F \in \mathbb{C}[[x_1, \ldots, x_k]]$  is analytic if there is a real constant  $\delta > 0$  such that the series  $F(\alpha_1, \ldots, \alpha_k)$  absolutely converges for all  $\alpha_i \in \mathbb{C}$ ,  $|\alpha_i| < \delta$ . Analytic power series form a ring denoted  $\mathbb{C}\{x_1, \ldots, x_k\}$ . This ring is closed also on division (if defined) and derivatives. For more information see Fischer [5] or Ruiz [19]. A Laurent series  $F \in \mathbb{C}((x_1, \ldots, x_k))$  is analytic if, for some  $p \in \mathbb{N}$ ,  $(x_1 \ldots x_k)^p F \in \mathbb{C}\{x_1, \ldots, x_k\}$ . Analytic Laurent series form again a ring closed on division and derivatives. We say that  $F \in \mathbb{C}[[x]]$  is  $D_A$ -finite if F solves the linear differential equation

$$R_0 F^{(m)} + R_1 F^{(m-1)} + \dots + R_{m-1} F' + R_m F + R_{m+1} = 0$$

with analytic coefficients  $R_i \in \mathbb{C}\{x\}, R_0 \neq 0$ .

**Fact.** If  $G \in \mathbb{C}[[x, y]]$ , G(0, 0) = 0, is nonzero and analytic, then every solution  $F \in \mathbb{C}[[x]]$ , F(0) = 0, of the equation

$$G(x,F) = 0$$

is analytic.

It is easy to see that the condition G(0,0) = 0, F(0) = 0 can be replaced by "the substitution of F(x) for y in G(x, y) makes formal sense". Also,  $\mathbb{C}[[x, y]]$  can be replaced by  $\mathbb{C}((x, y))$ . With the condition  $\frac{\partial G}{\partial y}(0,0) \neq 0$  this is the implicit function theorem for analytic functions (power series) and F is unique; see, for example, Hille [8] or Markushevich [14]. Without it one proceeds as follows. By the analytic version of the Weierstrass preparation theorem, see [5, p. 107] or [19, p. 14], G(x, F) = 0 is equivalent to a polynomial equation P(x, F) = 0 where  $P \in \mathbb{C}\{x\}[y]$  is monic. By the standard results of algebraic geometry on local parametrizations (Puiseux series), all solutions of P(x, F) = 0 are again analytic. A very readable account on these matters is Fischer [5, chapters 6 and 7].

**Theorem 1.** Let  $F \in \mathbb{C}[[x]]$  satisfy the differential equation

(1) 
$$F' = G(x, F)$$

where  $G \in \mathbb{C}((x, y))$ . Suppose that (i) F is not analytic, (ii) G is analytic, and (iii)  $\operatorname{ord}_y(G) < 0$ . Then F is not  $D_A$ -finite, the less D-finite.

*Proof.* Differentiating repeatedly (1) and substituting G(x, F) for F', we express the derivatives of F as

$$F^{(k)} = G_k(x, F)$$
 where  $G_{k+1}(x, y) = \frac{\partial G_k}{\partial x} + G \cdot \frac{\partial G_k}{\partial y}$  (and  $G_1 = G$ ).

By condition (iii),  $p = \operatorname{ord}_y(G) < 0$ . Using condition (ii) and induction on k, we see that each  $G_k \in \mathbb{C}((x, y))$  is analytic and

$$\operatorname{ord}_y(G_k) = k(p-1) + 1 < 0.$$

We suppose for the contradiction that F is  $D_A$ -finite and satisfies an equation

$$R_0 F^{(m)} + R_1 F^{(m-1)} + \dots + R_m F + R_{m+1} = 0$$

with  $R_i \in \mathbb{C}\{x\}$ ,  $R_0 \neq 0$ . We may assume that m > 0 because m = 0 implies  $F = -R_1/R_0 \in \mathbb{C}\{x\}$ and this contradicts condition (i). Replacing each  $F^{(k)}$  by  $G_k(x, F)$ , we obtain equation

$$H(x,F) = 0$$
 where  $H(x,y) = R_0(x)G_m(x,y) + \dots + R_m(x)y + R_{m+1}(x)$ .

The analyticity of the  $G_k$ s and  $R_i$ s implies the analyticity of H. It is clear that  $H \neq 0$  because  $\operatorname{ord}_y(H) = \operatorname{ord}_y(R_0G_m) = m(p-1)+1 < 0$ . Thus, by the above Fact,  $F \in \mathbb{C}\{x\}$ , in the contradiction with condition (i).

We have introduced the class of  $D_A$ -finite power series primarily to put our arguments in a natural setting — why to restrict to polynomials when the key lies in the analytic×nonanalytic dichotomy — but one could investigate the closure properties of the class as well. Clearly,  $D_A$ -finite power series are closed to addition and multiplication, the proof being the same as for D-finite power series, see [22, chapter 6]. D-finite power series are closed also to the Hadamard product that is defined as the coefficientwise multiplication. Are  $D_A$ -finite power series closed to the Hadamard product?

To the "non-Catalan" aspect of matching enumeration (many crossings, nonanalytic OGFs) are relevant articles (listed in chronological order) by Touchard [28, 29, 30], Kleitman [12], Hsieh [9], Riordan [18], Stein [23], Stein and Everett [24], Nijenhuis and Wilf [15], Read [17], Ismail, Stanton and Viennot [10], Penaud [16], Li and Sun [13], Cori and Marcus [4], Flajolet and Noy [6], Stoimenow [27], Broadhurst and Kreimer [3], Sawada [20], and Klazar [11]. We attempted to make this list of references complete or nearly complete.

Matchings are also called *complete pairings* or *linear(ized) chord diagrams*. In about 1993, their close relatives (circular) chord diagrams started to play, together with chord diagram algebras, an important role in studying knot invariants, especially Vassiliev knot invariants. By now more than 30 articles applying chord diagrams in knot theory have been published, of which we mention only Bar-Natan [1] and Stoimenow [25, 26]. This lower bound could be probably improved because Bar-Natan's bibliography on Vassiliev invariants, see [2], contains more than 400 items. Chord diagrams are orbits of the 2*n*-element cyclic group acting on matchings by cyclically reordering vertices in [2*n*]. They appear in knot theory as follows. Suppose K is a singular knot, which is a smooth embedding of the circle  $S^1$  in  $\mathbb{R}^3$  with n transversal self-intersections (double points). Fix an orientation of  $S^1$  and associate with K a chord diagram X: mark on  $S^1$  the 2*n* preimages of the self-intersections of K and connect by chord each two of them that correspond to one self-intersection. In fact, in these applications only those chord diagrams are important in which each chord crosses another chord (cf. cro<sub>n</sub>). We refer to [1] for further information.

Recall that  $\operatorname{con}_n$  and  $\operatorname{cro}_n$  are numbers of connected and crossing matchings with n edges. We can rephrase their definition in terms of *crossing graphs*. For a matching X, this is a graph G(X) = (X, E) with the vertex set X and the edge set  $E = \{\{A, B\} : A, B \in X \text{ cross }\}$ . Connected (crossing) matchings are matchings X for which G(X) is connected (has no isolated vertex).

Theorem 2. The power series

$$E = \sum_{n \ge 1} \operatorname{con}_n x^n = x + x^2 + 4x^3 + 27x^4 + \cdots$$

and

$$F = \sum_{n \ge 0} \operatorname{cro}_n x^n = 1 + x^2 + 4x^3 + 31x^4 + \cdots$$

satisfy the differential equations

(2) 
$$E' = \frac{E^2 + E - x}{2xE}$$

and

(3) 
$$F' = \frac{-x^2 F^3 + F - 1}{2x^3 F^2 + 2x^2 F}.$$

By Theorem 1, E and F are not  $D_A$ -finite. In particular, the sequences  $(con_n)$  and  $(cro_n)$  are not *P*-recursive.

*Proof.* By the spaces in a matching X on [2n] we mean the 2n - 1 inner spaces between the consecutive elements of [2n] and the two spaces before and after [2n]. Alltogether X has 2n + 1 spaces. Let X be a connected matching on [2n] with the first edge  $\{1, a\} \in X$ . Deleting it from X, the graph G(X) falls apart into k components to which correspond nonempty connected matchings  $X_1, \ldots, X_k$ .  $X_i$  has  $n_i > 0$  edges and  $X_1, \ldots, X_k$  are nested, which means that  $X_{i+1}$  lies in one of the  $2n_i - 1$  inner spaces of  $X_i$ . Also, a lies in one of the  $2n_k - 1$  inner spaces of  $X_k$ . There are no other restrictions on the  $X_i$ s and a. In terms of E,

$$E = x \sum_{k>0} (2xE' - E)^k.$$

Summing the geometric series and solving the result for E', we obtain equation (2).

Let X be a generic crossing matching on [2n]. We shall construct X by considering the edges which cross only the first edge  $\{1, a\} \in X$ . Suppose there are k such edges:  $\{b_i, c_i\} \in X$ ,  $1 < b_1 < b_2 < \cdots < b_k < a < c_k < \cdots < c_2 < c_1 \leq 2n$ . The remaining edges of X lie either inside or outside each  $\{b_i, c_i\}$  and the  $n_i \geq 0$  edges sandwiched between  $\{b_{i-1}, c_{i-1}\}$  and  $\{b_i, c_i\}$  form, for  $i = 1, 2, \ldots, k + 1$ , a (possibly empty) crossing matching  $X_i$ . We distinguish the cases k > 0 and k = 0.

Let k > 0. X arises by choosing a crossing matching  $X_1$  with  $n_1 \ge 0$  edges, inserting  $\{b_1, c_1\}$  in one of the  $2n_1 + 1$  spaces of  $X_1$ , inserting a crossing matching  $X_2$  inside  $\{b_1, c_1\}$ , inserting  $\{b_2, c_2\}$ in one of the  $2n_2 + 1$  spaces of  $X_2$  (inside  $\{b_1, c_1\}$ ), inserting a crossing matching  $X_3$  inside  $\{b_2, c_2\}$ , and so on. In the end  $X_{k+1}$  is inserted inside  $\{b_k, c_k\}$ , 1 is put in the beginning, and a is inserted in one of the  $2n_{k+1} + 1$  spaces of  $X_{k+1}$  (and inside  $\{b_k, c_k\}$ , of course).

If k = 0, just  $\{1, a\}$  is added to  $X_1$ . However, now the difference is that a cannot be inserted in each of the  $2n_1 + 1$  spaces of  $X_1$ . Since  $\{1, a\}$  must cross an edge, allowed are only spaces not separating  $X_1$  into two crossing matchings.

In terms of F,

$$F = 1 + \sum_{k \ge 0} (x(2xF' + F))^{k+1} - xF^2.$$

Term  $xF^2$  subtracts the number of spaces forbidden for a if k = 0. Summing the geometric series and solving the result for F', we obtain equation (3).

The conditions of Theorem 1 are easy to check. The right hand sides G of (2) and (3) are  $(2xy)^{-1}(-x+y+y^2)$  and  $(2x^2y)^{-1}(-1+y-x^2y^3)(1+xy)^{-1}$ , respectively. Thus both G belong to  $\mathbb{C}(x,y) \cap \mathbb{C}((x,y))$  (are rational Laurent series) and for both  $\operatorname{ord}_y(G) = -1$ . Conditions (ii) and (iii) are satisfied. Clearly,  $\operatorname{con}_n, \operatorname{cro}_n \geq (n-1)!$  for every  $n \in \mathbb{N}$ ; just intersect one edge by the others. Condition (i) is satisfied too and Theorem 1 applies.

Writing (2) as  $E = E \cdot (2xE' - E) + x$ , we obtain the recurrence

$$con_n = \sum_{k=1}^{n-1} (2k-1) \cdot con_k \cdot con_{n-k} 
= (n-1) \sum_{k=1}^{n-1} con_k \cdot con_{n-k} \quad (n \ge 2 \text{ and } con_1 = 1)$$

found in [23]. It was derived first by Stein and Riordan and then a simple bijective proof was provided in [15].

We give an alternative derivation for (2). Every matching X with n > 0 edges decomposes uniquely into a connected matching Y (the component of G(X) containing the vertex corresponding to the first edge of X) with m > 0 edges and 2m matchings  $Z_i$ ,  $i = 1 \dots 2m$ , which have  $l_i$  edges each and are inserted in the 2m + 1 spaces of Y except the first one. Each  $Z_i$  may be empty, no  $Z_i$  lies before Y, and  $m + l_1 + \dots + l_{2m} = n$ . Thus, if E is the above OGF of  $(\operatorname{con}_n)$  and  $M = \sum_{n \ge 0} \operatorname{mat}_n x^n = 1 + x + 3x^2 + 15x^3 + \dots$ ,

$$(4) E(xM^2) = M - 1.$$

Formula  $mat_n = (2n - 1)!!$  implies

(5) 
$$M' = \frac{(1-x)M - 1}{2x^2}$$

Here  $\operatorname{ord}_y(G)$  is nonnegative, of course.

We write  $E_0$  for  $E(xM^2)$  and  $E_1$  for  $E'(xM^2)$ . Thus (4) reads as  $M = E_0 + 1$ . Differentiating (4), solving the result for M', and replacing M by  $E_0 + 1$ , we obtain the equation  $M' = E_1(E_0 + 1)^2/(1 - 2xE_1(E_0 + 1))$ . Substituting in (5) for M' the latter expression and for M the expression  $E_0 + 1$ , we obtain a polynomial equation  $P(x, E_0, E_1) = 0$ . Rewriting it as  $Q(x(E_0 + 1)^2, E_0, E_1) = 0$  and substituting for x the power series inverse of  $xM^2 = x(E_0 + 1)^2$ , we get a polynomial equation Q(x, E(x), E(x)') = 0. It turns out to be (2).

This method may look complicated but its advantage is that it applies without much change to the natural generalization of the problem to r-matchings. These are partitions, for a fixed  $r \ge 2$ , of  $\{1, 2, \ldots, rn\}$  into r-element sets. Generalization of (non)crossing for r > 2 is easy (viz. noncrossing partitions). The formula for the number  $mat_n^r$  of all r-matchings and the analogue of (5) are straightforward. So is the analogue of (4). A differential equation for the OGF  $E_r$  of numbers  $con_n^r$ of connected r-matchings then can be derived by the method we have just indicated for r = 2. The equation involves derivatives of  $E_r$  to the order r - 1. The more combinatorial argument from the first part of the proof of Theorem 2 is still possible but for r > 2 becomes quite cumbersome.

From (4) also a relation between the numbers  $\operatorname{mat}_n$  and  $\operatorname{con}_n$  can be derived that, though standard, seems not to appear in the literature. Namely, writing  $M_0$  for  $xM^2$ , (4) can be restated as  $x(1+E(M_0))^2 = M_0$  and the Lagrange inversion formula gives  $[x^n]M_0(x) = n^{-1}[x^{n-1}](1+E(x))^{2n}$ . In terms of the coefficients,

$$(n+1)\sum_{i=0}^{n} (2i-1)!! \cdot (2n-2i-1)!! = \sum_{i=0}^{n} \operatorname{con}_{a_1} \cdot \operatorname{con}_{a_2} \cdot \ldots \cdot \operatorname{con}_{a_{2n+2}}.$$

Here  $(-1)!! = \operatorname{con}_0 = 1$  and the latter summation goes over all 2n + 2-tuples such that  $a_i \ge 0$  and  $a_1 + a_2 + \cdots + a_{2n+2} = n$ .

Restating equation (3) as  $F \cdot (1 + xF) \cdot (1 - xF - 2x^2F') = 1$  leads, after rearrangements and shifts of the indices, to the recurrence

$$\operatorname{cro}_n = \sum (2m+1) \cdot \operatorname{cro}_k \cdot \operatorname{cro}_l \cdot \operatorname{cro}_m \quad (n \ge 1 \text{ and } \operatorname{cro}_{-1} = \operatorname{cro}_0 = 1)$$

where we sum over  $-1 \le k, l, m \le n-1$  &  $k \ne -1$  & k+l+m = n-2. This is a more explicit relation than the formulas in Stoimenow [27, p. 217] where  $\operatorname{cro}_n$  is denoted  $\overline{\psi}_n$ . The first few values of these numbers are:

$$(\operatorname{con}_n)_{n\geq 2} = (1, 4, 27, 248, 2830, 38232, 593859, 10401712, \ldots)$$
  
 $(\operatorname{cro}_n)_{n\geq 2} = (1, 4, 31, 288, 3272, 43580, 666143, 11491696, \ldots).$ 

The matchings X on [2n] such that the graph G(X) is connected and X does not change when the linear order of [2n] is reversed were investigated first in [23]. Stein calls them *irreducible symmetric diagrams* and denotes their number, sco<sub>n</sub> in our notation, by  $\sigma_{2n}$ .

**Theorem 3.** The power series

$$F = \sum_{n \ge 1} \operatorname{sco}_n x^n = x + x^2 + 2x^3 + 7x^4 + \cdots$$

is related to the power series E of Theorem 2 by

(6) 
$$F^{2} - (1+x)F + E(x^{2}) \cdot (2xF' - F + 1) + x = 0$$

and satisfies the second order differential equation

(7) 
$$F''A_0 + (F')^3A_1 + (F')^2A_2 + F'A_3 + A_4 = 0$$

where  $A_i \in \mathbb{C}[x, F]$ , namely  $A_0 = -2x^2F^4 + (4x^2 + 4x^3)F^3 - (2x^2 + 8x^3 + 2x^4)F^2 + (4x^3 + 4x^4)F - 2x^4$ ,  $A_1 = 8x^5$ ,  $A_2 = 4x^2F^3 - (2x^2 + 6x^3)F^2 - (2x^2 - 4x^3 + 10x^4)F + 2x^3 + 10x^4$ ,  $A_3 = -5xF^4 + (7x + 7x^2)F^3 + (x - 11x^2 + 4x^3)F^2 + (-3x + x^2 - 8x^3)F - x + 3x^2 + 4x^3$ , and  $A_4 = F^5 - (2 + x)F^4 + (2x - x^2)F^3 + (2 + 3x^2)F^2 - (1 + 2x + 3x^2)F + x + x^2$ .

*Proof.* We say that a matching X is a *chain* if the connected matchings corresponding to the components of G(X) are nested as in the proof of Theorem 2. If X is in addition symmetric, we say that X is a *symmetric chain*. Now let X be a generic symmetric connected matching on [2n] with the first edge  $\alpha = \{1, a\}$  and the last edge  $\beta = \{2n - a + 1, 2n\}$ . Note that for n > 1 always  $\alpha \neq \beta$ . We delete  $\alpha$  and  $\beta$  from X and see what happens. After a while we see (since  $X \setminus \{\alpha, \beta\}$  must be symmetric) that the connected matchings corresponding to the components of  $G(X) \setminus \{\alpha, \beta\}$  have the following structure.

- On the top there is a (possibly empty) symmetric chain consisting of  $k \ge 0$  nonempty symmetric connected matchings  $X_1, X_2, \ldots, X_k$  where  $X_k$  is the innermost one.
- Then there are two (possibly empty) chains Y and Z such that Y precedes Z, Y and Z lie in a centrally symmetric pair of (not necessarily distinct) inner spaces of  $X_k$ , and Y and Z are reflections of one another.
- The endvertex of  $\alpha$  lies in an inner space of the innermost component of Y and the first vertex of  $\beta$  lies in an inner space of the innermost component of Z, or vice versa, and these spaces move one on the other upon reflection.
- If  $\alpha$  and  $\beta$  do not cross then  $k \ge 1$  (for X to be connected) else  $k \ge 0$ .

This is a complete description of the structure of X, that is to say, arranging  $\alpha$ ,  $\beta$ ,  $X_i$ , Y, and Z in any way meeting the above conditions, we get a symmetric connected matching. After another while we see that the conditions translate into the OGFs F and E as

$$F = 2E(x^{2}) \cdot xF' \cdot \sum_{k \ge 1} F^{k-1} + E(x^{2}) + x.$$

Factor 2 accounts for the crossing and noncrossing of  $\alpha$  and  $\beta$ , the first  $E(x^2)$  counts  $Y \cup Z \cup \{\alpha, \beta\}$ , the term xF' counts the positions of Y and Z in  $X_k$ ,  $F^{k-1}$  and xF' count the nonempty top symmetric chains, the second  $E(x^2)$  accounts for the case k = 0, and the last x accounts for n = 1. Summing the geometric series and rearranging the result, we obtain equation (6).

Equation (7) arises by solving (6) for  $E(x^2)$ , differentiating the result, and substituting in (2) for  $E(x^2)$  and  $E'(x^2)$  the expressions obtained.

Equating the coefficient at  $x^n$  in (6) to zero leads to the formula

$$\operatorname{sco}_{n} = \sum_{i=1}^{n-2} \operatorname{sco}_{i} \cdot \operatorname{sco}_{n-i} + \sum_{i=1}^{\lfloor n/2 \rfloor} (2n - 4i - 1) \cdot \operatorname{con}_{i} \cdot \operatorname{sco}_{n-2i}$$
  
where  $n \ge 2$ ,  $\operatorname{sco}_{0} = -1$ , and  $\operatorname{sco}_{1} = 1$ .

This recurrence is simpler and more transparent than the one in the end of [23]. We have

 $(sco_n)_{n\geq 2} = (1, 2, 7, 22, 96, 380, 1853, 8510, 44940, \ldots),$ 

in agreement with the table in [23]. A much more complicated recurrence but only in terms of the numbers  $sco_n$  themselves can be obtained from (7).

We say that a sequence  $(a_n)$  of integers is *P*-recursive modulo  $m \in \mathbb{N}$  if there is a P-recursive sequence of integers  $(b_n)$  such that, for each  $n, a_n \equiv b_n \mod m$ .

**Theorem 4.** The sequences  $(con_n)$  and  $(cro_n)$  are *P*-recursive modulo  $2^k$  for each  $k \in \mathbb{N}$ .

Proof. We rewrite equation (2) for  $E = \sum_{n \ge 1} \operatorname{con}_n x^n$  as  $E^2 + E - x = 2xEE'$ . Differentiating it, replacing E' on the left by  $(E^2 + E - x)/(2xE)$ , and multiplying all by 2xE, we obtain equation  $2E^3 + 3E^2 + (1 - 4x)E - x = P_2$  where  $P_2 \in 4\mathbb{Z}[x, E, E', E'']$ . Continuing this way, we derive for each  $k \in \mathbb{N}$  an equation

$$Q_k(x, E) = P_k(x, E, E', \dots, E^{(k)})$$

where  $Q_k$  and  $P_k$  are integral polynomials and  $P_k$  vanishes identically modulo  $2^k$ . It is easy to check that the polynomials  $[y^1]Q_k \in \mathbb{Z}[x]$  and  $[y^0]Q_k \in \mathbb{Z}[x]$  have constant terms 1 and 0, respectively. Thus the equation

$$Q_k(x, A_k) = 0$$

has a unique power series solution  $A_k \in \mathbb{Z}[[x]]$  with  $A_k(0) = 0$ . It follows by the induction on n that  $\operatorname{con}_n = [x^n]E \equiv [x^n]A_k \mod 2^k$  for each n.  $A_k$  is algebraic over  $\mathbb{Z}(x)$  and thus D-finite ([22, Theorem 6.4.6]). Hence the coefficients of  $A_k$  form a P-recursive sequence and  $(\operatorname{con}_n)$  is P-recursive modulo  $2^k$ .

The proof for the sequence  $(cro_n)$  is similar and we omit it.

The Lagrange inversion formula yields some explicit congruences. We give one example. Reducing (3) modulo 2 we obtain equation  $x^2F^3 - F + 1 = 0$ . With G = xF it becomes  $G = x(1 - G^2)^{-1}$  and the Lagrange inversion shows that  $\operatorname{cro}_{2n+1}$  is always even and

$$\operatorname{cro}_{2n} \equiv \frac{1}{2n+1} \binom{3n}{n} \equiv \binom{3n}{n} \mod 2.$$

The well known result saying that the highest power of a prime p dividing n! has the exponent  $\lfloor n/p \rfloor + \lfloor n/p^2 \rfloor + \lfloor n/p^3 \rfloor + \cdots$  gives a more explicit criterion:  $\operatorname{cro}_{2n}$  is even if and only if there is an  $r \in \mathbb{N}$  such that the residue of n on division by  $2^r$  is greater than  $2^{r+1}/3$ . For example,  $\operatorname{cro}_{304}$  is even (set r = 5) but  $\operatorname{cro}_{296}$  is odd.

**Conclusion.** As we have indicated, the problem of enumeration of matchings with restricted crossings can be generalized to set partitions with *r*-element blocks (or even to set partitions without any restriction on block sizes). However, then differential equations of more complicated types than (1) arise, such as (7). Probably, these generalizations require less amateur approach with some techniques from the differential algebra. Or maybe from the geometry, as suggested by an anonymous referee. We conjecture that neither the sequence  $(sco_n)$  is P-recursive. Another research direction is to investigate the behaviour of  $con_n$  and  $cro_n$  with respect to other moduli. We conjecture that these numbers are P-recursive modulo m only if  $m = 2^k$ .

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