

L 10

Thm. (P.L. Čebyshev) ①
 $\exists c_1 > c_2 > 0 \forall x \geq 2:$

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 $\doteq 185.0$

$$\frac{c_2 x}{\log x} < \pi(x) < \frac{c_1 x}{\log x}.$$

$\{p \in \mathbb{P} \mid p \leq x \}$.

Proof. $n \in \mathbb{N}$, then a) $\frac{4^n}{2n+1} \stackrel{\textcircled{1}}{\leq} \binom{2n}{n} \stackrel{\textcircled{2}}{\leq} 4^n$. This

follows from the binomial expansion $4^n = 2^{2n} =$

$$= (1+1)^{2n} = \sum_{i=0}^{2n} \binom{2n}{i} \text{ and the inequalities } \binom{2n}{i} \geq 0$$

and $\binom{2n}{i} \leq \binom{2n}{n}$. Also, b) $\pi(p) \stackrel{\textcircled{1}}{\leq} \binom{2n}{n} \stackrel{\textcircled{2}}{\leq} \pi(2n)$

$$n \leq p \leq 2n$$

The first \leq follows from the fact that $\binom{2n}{n} = \frac{(2n)!}{n! n!}$,

so in fact $(\pi(p)) \mid \binom{2n}{n}$. The second \leq was proven in

the previous lecture (the 2nd proof of P. Erdős): if

$\binom{2n}{n} = p_1^{a_1} p_2^{a_2} \cdots p_q^{a_q}$ then $p_i \leq 2n$ and even $p_i \leq 2n$. Combining $\stackrel{\textcircled{1}}{\leq}$ and $\stackrel{\textcircled{2}}{\leq}$ (which we did already in.) we get

$$\frac{4^n}{2n+1} \leq (2n)^{\pi(2n)} \text{ thus } \text{Thm: } \pi(2n) \geq \frac{2n \log 2 - \log(2n+1)}{\log(2n)} \frac{1}{\log(2n)}$$

$c_1 > (\log 2) \frac{2n}{\log(2n)} - 2$. Exercise for you: deduce from

this that indeed for some $c_2 > 0$ and every $x \geq 2$, $\pi(x) > \frac{c_2 x}{\log x}$.

Combining ① and ② we get that there:

(2)

$$\sum_{\substack{n < p \leq 2n \\ (\text{unique})}} \log p \leq 4^n \Rightarrow \sum_{n < p \leq 2n} \log p \leq n \log 4. \text{ For real } x > 2.$$

We take $2^j \in \mathbb{N}$ s.t. $2^j \leq x \leq 2^{j+1}$. Then $\sum_{p \leq x} \log p \leq \sum_{j=0}^{\infty} \sum_{\substack{p \\ 2^j < p \leq 2^{j+1}}} \log p \leq \sum_{j=0}^{\infty} 2^j \log 4 = (2^{j+1} - 1) \log 4 < 2 \cdot 2^j \log 4$

$\leq (2 \log 4) x$. So there: $\sum_{p \leq x} \log p \leq (2 \log 4) x$.

$x > 2$:

$$\Rightarrow (2 \log 4) x > \sum_{\sqrt{x} < p \leq x} \log p \geq (\pi(x) - \pi(\sqrt{x})) \log(\sqrt{x})$$

$0 \leq \sum_{p \leq x}$

$$\Rightarrow \pi(x) \leq \frac{(2 \log 4) x}{\log(\sqrt{x})} + \sqrt{x} = \frac{(4 \log 4) x}{\log x} + \sqrt{x}.$$

Exercise for you: deduce from this that indeed for some $C_1 > 0$ and every real $x > 2$:

$$\pi(x) < \frac{C_2 x}{\log x}.$$



I mention some further classical results from the theory of prime numbers, without proofs (almost).

PNT - The Prime Number Theorem (1896, J. Hadamard; de la Vallée-Poussin) For real $x \rightarrow +\infty$,

$$\pi(x) \sim \frac{x}{\log x}, \text{ i.e. } \pi(x) = (1 + o(1)) \frac{x}{\log x}.$$

Jacques Hadamard (1865-1963),

Charles J. de la Vallée Poussin (1866-1962),

$L_i(x) := \int_2^x \frac{dt}{\log t}$. Then, more precisely than

above, $\pi(x) = \underbrace{L_i(x)}_{\sim \frac{x}{\log x}} + O\left(x e^{-A(\log x)^{3/5}} / (\log \log x)^{1/5}\right)$, where $A = 0.2088$ (K. Ford, 2002)

• Kevin Ford (1967), [the Riemann Hypothesis]

(RH): $S(s) = \sum_{n=1}^{\infty} \frac{1}{n^s} : \mathbb{C} \setminus \{1\} \rightarrow \mathbb{C}, S(s) = 0$ with $\operatorname{Re}(s) > 0 \Rightarrow \operatorname{Re}(s) = \frac{1}{2}$. Explicit formula for $\pi(x)$.

H. von Koch (1901): RH $\Rightarrow \pi(x) = L_i(x) + O(\sqrt{x} \log x)$

three formulas of Franz Neumann (1840-
 $x \rightarrow \infty$ (≤ 1874) -1727)

① $\sum_{p \leq x} \frac{\log p}{p} = \log x + O(1)$. Proofs: See my LNs.

② $\sum_{p \leq x} \frac{1}{p} = \log(\log x) + c + O(1/\log x)$, where c is a constant.

③ $\prod_{p \leq x} \left(1 - \frac{1}{p}\right) = \frac{d}{\log x} \left(1 + O(1/\log x)\right)$, where $d > 0$ is a constant,

For $\prod p_i^{a_i} = n$ we define $w(n) := \sum a_i$, $\Omega(n) := \sum_{d|n} d$, $\tau(n) = (1+a_1)(1+a_2)\dots(1+a_k)$ (number of divisors of n). For real $x \geq 2$,

- $\sum_{u \leq x} w(u) = x \cdot \log(\log x) + c_1 x + O(x/\log x)$
- $\sum_{u \leq x} \Omega(u) = x \cdot \log(\log x) + c_2 x + O(x/\log x)$ where $c_i, i=1,2$, are constants.

Almost every n where
 $\uparrow (\log n)^{\log 2 - \varepsilon} < \tau(n) < \downarrow x^{\log f + \varepsilon}$

Theorem (Hardy-Ramanujan, 1917)

$\forall \varepsilon > 0 \exists x_0 = x_0(\varepsilon) > 0$ s.t.

$x > x_0 \Rightarrow \#\{u \leq x \mid |w(u) - \log(\log u)| < \varepsilon \log(\log u)\} > \frac{x}{(1-\varepsilon)x}$.

the multiplicative table result (or $\log(\log u)$) does not matter.
 Also holds for $\Omega(u)$ too.

Godfrey H. Hardy (1877-1947) • Thangal Ramanujan (1887-1920)

Proof. (see my LNs) by the 2nd moment method
 by Pan (Tuvor) (1910-1976), Drin (On-Space)
 J(Pál) • The Probabilistic Method

Chapter 5 - Congruences

(5)

$a \not\equiv 0 \pmod{p}$

Theory of quadratic residues

$a \in \mathbb{Z}, p \in \mathbb{P} : a \text{ is a (quadratic) residue mod } p \Leftrightarrow$
 $x^2 \equiv a \pmod{p}$ has a solution $x \in \mathbb{Z}$. Else a is a (quadratic) non-residue.

For example, for $p=11$
 The residues are $1, 4, 9, 5, 3$. The rest $-2, 6, 7, 8, 10$ are q. non-residues. For $p=2$, 1 is a q.r. and there is no q. non-residue.

Proposition For any prime $p > 2$, the # of q.r. mod $p =$ the # of q. non-r. mod $p = \frac{p-1}{2}$.

Proof. $\mathbb{Z}_p^\times := \mathbb{Z}_p \setminus \{0\}$, the map $\mathbb{Z}_p^\times \ni x \mapsto x^2 \in \mathbb{Z}_p^\times$ is two-to-one: $x^2 \equiv y^2 \pmod{p} \Leftrightarrow (x-y)(x+y) \equiv 0$ (and $x \not\equiv -y \pmod{p}$ as $p > 2$). \square $\blacksquare \pmod{p}$

Standard notation: $a \in \mathbb{Z}, p > 2$ a prime

- Legendre's symbol: $\left(\frac{a}{p}\right) := \begin{cases} 1 & \text{if } a \text{ is q.r.} \\ -1 & \text{if } a \text{ is q. non-r.} \\ 0 & \text{if } a \equiv 0 \pmod{p} \end{cases}$

Trivially: $a \equiv b \pmod{p} \Rightarrow \left(\frac{a}{p}\right) = \left(\frac{b}{p}\right)$,
 $\left(\frac{ab}{p}\right) = \left(\frac{a}{p}\right), \left(\frac{b}{p}\right)$, ($b \not\equiv 0 \pmod{p}$).

Proposition (Euler's criterion) $\forall a \in \mathbb{Z} \setminus \{0\}, \forall p > 2$: (6)

$$\left(\frac{a}{p}\right) \equiv a^{\frac{p-1}{2}} \pmod{p}.$$

Proof. Recall the Little theorem of Fermat: $a \neq 0 \Rightarrow a^{p-1} \equiv 1$. So

$(a^{\frac{p-1}{2}} - 1)(a^{\frac{p-1}{2}} + 1) \equiv 0$ and $a^{\frac{p-1}{2}} \equiv \pm 1$. If $a \equiv 0$ then $0 \equiv 0 \pmod{p}$. \checkmark . If $a \neq 0$ and is a q.r., then

$b^2 \equiv a \pmod{p}$, so $a^{\frac{p-1}{2}} = (b^2)^{\frac{p-1}{2}} = b^{p-1} \equiv 1$ (by FET).

It remains to prove that if a is a q.-non-r., $\left(\frac{a}{p}\right) \equiv -1$.

then $a^{\frac{p-1}{2}} \equiv -1 \pmod{p}$. This follows from the

overn in algebra that $\# \{x \in F \mid f(x) = 0\}_F \leq \deg(f)$. Eu

our case $f(x) = x^{\frac{p-1}{2}} - 1$ and $F = \mathbb{Z}_p = \mathbb{Z}/p\mathbb{Z}$.

$x^{\frac{p-1}{2}} - 1 = 0$ has $\frac{p-1}{2}$ solutions in \mathbb{Z}_p (namely the q.

r. (see 9 and the previous prop.) $\Rightarrow x^{\frac{p-1}{2}} = 1$

~~if x is a q.-non-r.~~ \square

Proposition $\forall a, b \in \mathbb{Z} \setminus \{0\}, \forall p > 2$: $\left(\frac{ab}{p}\right) = \left(\frac{a}{p}\right) \left(\frac{b}{p}\right)$.

~~P.~~ $\left(\frac{ab}{p}\right) = (ab) \stackrel{\frac{p-1}{2}}{=} a^{\frac{p-1}{2}} b^{\frac{p-1}{2}} \equiv \left(\frac{a}{p}\right) \left(\frac{b}{p}\right)$, hence \square . \square