Spanning Trees Exercises 2.

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## Spanning tree

A spanning tree T of a connected, undirected graph G is a tree composed of all the vertices and some (or perhaps all) of the edges of G. In other words, a spanning tree of G is a selection of edges of G that form a tree spanning every vertex. That is, every vertex lies in the tree, but no cycles (or loops) are allowed. On the other hand, every bridge of G must belong to T. A spanning tree of a connected graph G can also be defined as a maximal set of edges of G that contains no cycle, or as a minimal set of edges that connect all vertices.

## Counting spanning trees

The number t(G) of spanning trees of a connected graph is a well-studied invariant. In some cases, it is easy to calculate t(G) directly. For example, if G is itself a tree, then t(G) = 1, while if G is the cycle graph  $C_n$  with n vertices, then t(G) = n. For any graph G, the number t(G) can be calculated using Kirchhoff's matrix-tree theorem.

Here are some known results concerning counting spanning trees of graphs.

- Complete graph  $K_n$ :  $t(K_n) = n^{n-2}$  (Cayley's formula),
- **2** Complete bipartite graph  $K_{n,m}$ :  $t(K_{n,m}) = m^{n-1}n^{m-1}$ ,

3 *n*-dimensional cube graph 
$$Q_n$$
:  $t(Q_n) = 2^{2^n - n - 1} \prod_{k=2}^n k^{\binom{n}{k}}$ .

# Spanning Trees

# Kirchhoff Matrix-Tree Theorem

The celebrated Kirchhoff Matrix-Tree Theorem is the following statement.

## Theorem (Kirchhoff (1847))

All cofactors of Laplacian matrix L(G) are equal to t(G).

More convenient form of this result were obtained by A. K. Kel'mans and V. M. Chelnokov.

## Theorem (Kel'mans, Chelnokov (1974))

Let  $0 = \lambda_1 \leq \lambda_2 \leq \cdots \leq \lambda_n$  denote the eigenvalues of the Laplace matrix L(X) of a n point graph X. Then

$$t(X) = \frac{1}{n} \prod_{k=2}^{n} \lambda_k.$$

# Spanning Trees

A generalisation of the Matrix-Tree-Theorem was obtained by Kelmans (1967) who gave a combinatorial interpretation to all the coefficients of  $\mu(X, x)$  in terms of the numbers of certain subforests of the graph. This result has been obtained even in greater generality (for weighted graphs) by Fiedler and Sedláček.

## Theorem (Kel'mans (1967))

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$$\mu(X,x) = x^{n} - c_{1}x^{n-1} + \ldots + (-1)^{i}c_{i}x^{n-i} + \ldots + (-1)^{n-1}c_{n-1}x^{n-1}$$

then

$$c_i = \sum_{S \subset V, |S|=n-i} t(X_S),$$

where t(H) is the number of spanning trees of H, and  $X_S$  is obtained from X by identifying all vertices of S to a single one.

From the last theorem we can derive useful corollary.

### Corollary

The degree of Laplacian polynomial  $\mu(X, x)$  is equal to n = |V(X)|. Its coefficients  $c_1$  and  $c_{n-1}$  are given by the formulas  $c_1 = 2|E(X)|$  and  $c_{n-1} = |V(X)| \cdot t(X)$ .

Hence, the number of vertices |V(X)|, number of edges |E(X)| and the number of spanning trees t(X) are uniquely defined by the Laplacian polynomial.

# Exercises

### Exercise 2.1.

Prove that the number of spanning trees for the path graph  $P_n$  is 1.

#### Exercise 2.2.

Prove that the number of spanning trees for the cyclic graph  $C_n$  is n.

#### Exercise 2.3.

Prove the Cayley formula for the number of spanning trees for the complete graph  $K_n$ :  $t(K_n) = n^{n-2}$ .

### Exercise 2.4.

Prove that the number of spanning trees for the complete bipartite graph  $K_{n,m}$  is given by the formula  $t(K_{n,m}) = m^{n-1}n^{m-1}$ .

**Exercise 2.5.** Denote by G - e the graph obtained by removing edge e from the graph G. Let  $G \setminus e$  be the graph obtained from graph G by contracting edge e. In other words,  $G \setminus e$  is obtained by deleting edge e and identifying its ends. Prove the following formula

$$t(G) = t(G - e) + t(G \setminus e).$$

**Exercise 2.6.** Denote by  $G_{s,e}$  the graph resulting from subdivision of an edge *e* of a graph *G*. Then

$$t(G_{s,e}) = t(G \setminus e) + 2t(G - e) = t(G) + t(G - e).$$

### Exercise 2.7.

Find the number of spanning trees for the wheel graph  $W_n = K_1 * C_n$ . **Answer:** If *n* is odd then  $t(W_n) = \ell_k^2$ , if *n* is even then  $t(W_n) = 5f_n^2$ , where  $\ell_j$  is *j*-th Lukas number and  $f_k$  is *k*-th Fibonacci number. Note:

$$\begin{split} \ell_1 &= 1, \ \ell_2 = 3, \ \ell_{k+2} = \ell_{k+1} + \ell_k, \ k \geq 1. \\ f_1 &= 1, \ f_2 = 1, \ f_{k+2} = f_{k+1} + f_k, \ k \geq 1. \\ f_{2n} &= \ell_n \cdot f_n \text{ and } \ell_n = f_{n-1} + f_{n+1}. \end{split}$$

#### Exercise 2.8.

Find the number of spanning trees for the fan graph  $F_n = K_1 * P_n$ . **Answer:**  $t(F_n) = f_{2n}$ .

#### Exercise 2.9.

Find the number of spanning trees for the lattice graph  $L_{m,n} = K_m \times K_n$ . **Answer:**  $t(K_m \times K_n) = m^{m-2}n^{n-2}(m+n)^{(m-1)(n-1)}$ .

#### Exercise 2.10.

Prove that the following result by Boesch and Prodinger

$$t(K_m \times C_n) = \frac{n}{m} 2^{m-1} (T_n(1+\frac{m}{2})-1)^{m-1},$$

where  $T_n(x) = \cos(n \arccos x)$  is the Chebyshev polynomial of the first kind.

### Exercise 2.11.

Prove that the number of spanning trees for the prism  $P_2 \times C_n$  is given by the formula

$$t(P_2 \times C_n) = \frac{n}{2}((2+\sqrt{3})^n + (2-\sqrt{3})^n - 2).$$

#### Exercise 2.12.

Prove that the number of spanning trees for the Moebius ladder graph  $M_n$  is given by the formula

$$t(M_n) = \frac{n}{2}((2+\sqrt{3})^n + (2-\sqrt{3})^n + 2).$$

## Chebyshev polynomials

The Chebyshev polynomial of the first kind is defined by the formula

 $T_n(x) = \cos(n \arccos x).$ 

Equivalently,

$$T_n(x) = \frac{(x + \sqrt{x^2 - 1})^n + (x - \sqrt{x^2 - 1})^n}{2}.$$

Also,  $T_n(x)$  satisfies the recursive relation

 $T_0(x) = 1, \ T_1(x) = x, \ T_n(x) = 2x \cdot T_{n-1}(x) - T_{n-2}(x), \ n \ge 2.$ 

# Addendum

## Chebyshev polynomials

The Chebyshev polynomial of the second kind is defined by the formula

$$U_n(x) = \frac{\sin((n+1)\arccos x)}{\sin(\arccos x)}$$

Equivalently,

$$U_n(x) = \frac{(x + \sqrt{x^2 - 1})^{n+1} - (x - \sqrt{x^2 - 1})^{n+1}}{2\sqrt{x^2 - 1}}.$$

Also,  $U_n(x)$  satisfies the recursive relation

$$U_0(x) = 1, U_1(x) = 2x, U_n(x) = 2x \cdot U_{n-1}(x) - U_{n-2}(x), n \ge 2.$$

We have  $U_n(\cos \frac{k\pi}{n+1}) = 0, k = 1, 2, \dots, n$ . Hence

$$U_n(x) = 2^n \prod_{k=1}^n (x - \cos \frac{k\pi}{n+1}).$$

## Chebyshev polynomials Since

$$U_n(x) = (-1)^n U_n(-x) = 2^n \prod_{k=1}^n (x + \cos \frac{k\pi}{n+1})$$

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we obtain

$$U_n^2(x) = \prod_{k=1}^n (4x^2 - 4\cos^2\frac{k\pi}{n+1}).$$

Polynomials  $T_n(x)$  and  $U_{n-1}(x)$  are related by the following identity

$$T_n^2(x) + (x^2 - 1)U_{n-1}^2(x) = 1.$$

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# Chebyshev polynomials

Consider  $n \times n$  matrix

$$A_n(x) = \begin{pmatrix} 2x & -1 & 0 & 0 & \dots & 0 & 0 \\ -1 & 2x & -1 & 0 & \dots & 0 & 0 \\ 0 & -1 & 2x & -1 & \dots & 0 & 0 \\ 0 & 0 & -1 & 2x & \dots & 0 & 0 \\ \vdots & \vdots & \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & 0 & 0 & 0 & \dots & 2x & -1 \\ 0 & 0 & 0 & 0 & \dots & -1 & 2x \end{pmatrix}$$

Then det  $A_n(x) = U_n(x)$ .

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