# Arrangements of pseudocircles and circles 

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A line arrangement


## A pseudoline arrangement



## Pseudoline arrangements

A pseudoline is the image of a line under a homeomorphism of $\mathbb{R}^{2}$. A pseudoline arrangement satisfies the following.

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The study of such objects apparently goes back at least to 1826 (J. Steiner, vol. 1 of Crelle's).

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F.W. Levi (1926). Hint: consult Pappus of Alexandria (c. 340 AD).

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Ringel (1956).

## Small pseudoline arrangements are stretchable

Goodman and Pollack (1980), after a conjecture of Grünbaum, showed every arrangement of eight pseudolines is equivalent to a line arrangement. In other words, every arrangement of eight pseudolines is stretchable.

Richter-Gebert (1989) showed Ringel's example is the unique non-stretchable simple arrangement of nine pseudolines.

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Theorem (Mnëv, 1988, cf. Shor, 1991)
SIMPLE STRETCHABILITY is NP-hard.
Note 1: The theorem may be viewed as a corollary to a deep topological theorem of Mnëv, but Shor gave a simpler and direct proof.

Note 2: In fact, SIMPLE STRETCHABILITY is complete for the computational class "existential theory of the reals" $\dagger$.

[^0]
## A circle arrangement



A pseudocircle arrangement


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3. No point lies on three pseudocircles.

Note for context: one may interpret simple line arrangements as arrangements of great circles on a sphere.

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Linhart and Ortner (2005) conjectured this to be the smallest non-circleable arrangement.

## Small pseudocircle arrangements are circleable

## Theorem (K and Müller)

Every pseudocircle arrangement on four pseudocircles is circleable.
The (lengthy) case analysis to prove this is much more involved than Goodman and Pollack's (compact) argument for pseudoline arrangements, but fortunately it is shortened by the use of circle inversions.

## Circle inversions



## Circle inversions



Fun facts about an inversion in a circle $C$ with centre $p$ : It

- exchanges the interior and exterior of $C$;
- maps circles through $p$ to lines, circles not through $p$ to circles, lines through $p$ to lines, lines not through $p$ to circles;
- is conformal.


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CONVEX CIRCLEABILITY is NP-hard.

## CIRCLEABILITY reductions

Starting with a simple pseudoline arrangment, perform an inversion on a circle with centre not on any pseudoline.

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## A non-convexible arrangement

Linhart and Ortner (2008) asked for a non-convexible arrangement:
Take Ringel's construction and perform the last reductions to obtain a non-convexible arrangement of eighteen pseudocircles.

What is the smallest one?

## One more pseudocircle result

We learned of the following conjecture, of Russian folklore, from Artem Pyatkin.

Given an arrangement of convex pseudocircles, the corresponding intersection graph can be realised as the intersection graph of a collection of circles.
(The intersection graph has vertex set the collection of pseudocircles, two vertices adjacent if the corresponding pseudocircles intersect.)

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(The intersection graph has vertex set the collection of pseudocircles, two vertices adjacent if the corresponding pseudocircles intersect.)
Theorem (K and Müller)
The following depicts a counter-example to the above conjecture.

## One more pseudocircle result



## Open problems

1. What are all the smallest non-circleable (convex) pseudocircle arrangements?
2. What is the size of a smallest non-convexible (simple) arrangement of pseudocircles?
3. What is the size of a smallest convex pseudocircle arrangement, the corresponding intersection graph of which cannot be realised as the intersection graph of a collection of circles?
4. What is the computational complexity of recognising if a convex pseudocircle arrangement is realisable as the intersection graph of a collection of circles?

## Thank you!




[^0]:    ${ }^{\dagger}$ Given an existential first-order sentence over the real numbers, is it true?

