## Tight lower bounds for the size of epsilon-nets

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## Introduction

Couple  $(X, \mathcal{R})$  is called range space in universe  $\mathcal{U}$  if  $X \subset \mathcal{U}$  is a finite set and  $\mathcal{R} \subset 2^{\mathcal{U}}$  is a system of sets. The set  $A \subset X$  is called shattered if for every  $B \subset A$  there is a set  $R_B \in \mathcal{R}$  such that  $R_B \cap A = B$ . The size of the largest shattered subset of X is called the dimension of the range space  $(X, \mathcal{R})$ .

For every  $\varepsilon > 0$ , the set  $S \subset X$  is called the  $\varepsilon$ -net for the range space  $(X, \mathcal{R})$  if every range  $R \in \mathcal{R}$  with  $|R \cap X| \ge \varepsilon |X|$  contains at least one element of S.

**Theorem** (Matoušek, Seidl, Welzl, 1990-2) All range spaces  $(X, \mathcal{R})$ , where X is a finite set of points in  $\mathbb{R}^3$  and  $\mathcal{R}$  consists of half-spaces, admit  $\varepsilon$ -nets of size  $O(\frac{1}{\varepsilon})$ .

**Theorem** (Aronov, Ezra, Sharir, 2010) All range spaces  $(X, \mathcal{R})$ , where X is a finite set of points in  $\mathbb{R}^2$  (or  $\mathbb{R}^3$ ) and  $\mathcal{R}$  consists of axis-parallel rectangles (boxes), admit  $\varepsilon$ -nets of size  $O(\frac{1}{\varepsilon} \log \log \frac{1}{\varepsilon})$ .

Let  $(X, \mathcal{R})$  be a range space with ranges from  $\mathbb{R}^m$ . The dual range space  $(\mathbb{R}, \bigcup_{x \in \mathbb{R}^m} \mathcal{R}_x)$  is defined as a system on the underlying set  $\mathcal{R}$  consisting of the sets  $\mathcal{R}_x = \{R \mid x \in R \in \mathcal{R}\}$ , for all  $x \in \mathbb{R}^m$ .

## Main results

**Theorem 1** For any  $\varepsilon > 0$  and for any sufficiently large integer  $n \ge n_0(\varepsilon)$ , there exists a dual range space  $\sum^*$  of VC-dimension 2, induced by n axis-parallel rectangles in  $\mathbb{R}^2$ , in which the minimum size of an  $\varepsilon$ -net is at least  $C^{\frac{1}{\varepsilon}} \log \frac{1}{\varepsilon}$ . Here C > 0 is an absolute constant.

**Theorem 2** For any  $\varepsilon > 0$  and for any sufficiently large integer  $n \geq n_0(\varepsilon)$ , there exists a (primal) range space  $\sum = (X, \mathcal{R})$  of VC-dimension 2, where X is a set of n points of  $\mathbb{R}^4$ ,  $\mathcal{R}$  consists of axis-parallel boxes with one of their vertices at the origin, and in which the size of the smallest  $\varepsilon$ -net is at least  $C^{\frac{1}{\varepsilon}} \log \frac{1}{\varepsilon}$ . Here C > 0 is an absolute constant.

**Theorem 3** For any  $\varepsilon > 0$  and for any sufficiently large integer  $n \geq n_0(\varepsilon)$ , there exists a (primal) range space  $\sum = (X, \mathcal{R})$  of VC-dimension 2, where X is a set of n points of  $\mathbb{R}^4$ ,  $\mathcal{R}$  consists of half-spaces, and in which the size of the smallest  $\varepsilon$ -net is at least  $C^{\frac{1}{\varepsilon}} \log \frac{1}{\varepsilon}$ . Here C > 0 is an absolute constant.

**Theorem 4** For any  $\varepsilon > 0$  and for any sufficiently large integer  $n \ge n_0(\varepsilon)$ , there exists a (primal) range space  $\sum = (X, \mathcal{R})$ , where X is a set of n points in the plane,  $\mathcal{R}$  consists of axis-parallel rectangles, and in which the size of the smallest  $\varepsilon$ -net is at least  $C^{\frac{1}{\varepsilon}} \log \log \frac{1}{\varepsilon}$ . Here C > 0 is an absolute constant.

The structure of the proofs in the paper is the following:

$$\begin{array}{c} \text{Lemma 1} \\ \text{Lemma 2} \end{array} \} \Longrightarrow \text{Theorem 1} \Longrightarrow \text{ Theorem 2} \\ \text{Lemma 3} \end{array} \} \Longrightarrow \text{Theorem 3}$$
 
$$\text{Lemma 4} \Longrightarrow \text{Theorem 4}$$

## Useful tools

Let  $c \geq 2$  and  $d \geq 1$  be integers. Let  $x \in [c]^k = \{0, 1, \ldots, c-1\}^k$ , that is  $x = x_1 x_2 \ldots x_k$ ,  $k \in [d]$ . Expanding x as a c-ary fraction we define  $\bar{x} = \sum_{j=1}^k x_j/c^k$ . For any  $0 \leq k \leq d$ ,  $u \in [c]^k$  and  $v \in [c]^{d-k}$  we define an open axis-parallel rectangle  $R_{u,v}^k$  in the unit square as

$$R_{u,v}^k = (\bar{u}, \bar{u} + c^{-k}) \times (\bar{v}, \bar{v} + c^{k-d})$$

and consider the family

$$\mathcal{R} = \mathcal{R}(c, d) = \{ R_{u,v}^k \mid 0 \le k \le d, u \in [c]^k, v \in [c]^{d-k}, u_k = v_{d-k} \}$$

Clearly  $|\mathcal{R}| = (d+1)c^{d-1}$ . Finally  $\sum = \sum (c,d)$  be the infinite range space  $(\mathbb{R}^2, \mathcal{R})$  and let  $\sum^* = \sum^* (c,d)$  denote its dual range space.

**Lemma 1** Let  $d \ge 1$ ,  $r \ge 2$ ,  $c \ge 3$  and let  $\sum^* = \sum^* (c, d)$ . If a subset  $I \subset \mathcal{R}(c, d)$  contains no r-element range of  $\sum^*$  then

$$|I| \le (r-1)\frac{c-1}{c-2}c^{d-1}.$$

**Lemma 2** Both  $\sum$  and  $\sum^*$  have VC-dimension 2.

**Lemma 3** Let P be a finite set of points in the positive orthant of  $\mathbb{R}^d$ . To each  $p \in P$  we can assign a point p' in the positive orthant of  $\mathbb{R}^d$  so that the set  $P' = \{p' \mid p \in P\}$  satisfies the following condition. For any axis-parallel box  $B \subset \mathbb{R}^d$  that contains the origin, there is a half-space  $H_B \subset \mathbb{R}^d$  which contains the origin and for which  $\{p' \mid p \in B \cap P\} = P' \cap H_B$ .

**Lemma 4** Let  $n \geq 2$ ,  $r = \lceil \log \log n/5 \rceil$  be integers and  $\varepsilon = r/n$ . Let X be a set of n randomly uniformly selected points of unit square, and  $\mathcal{R}$  denote the family of all axis-parallel rectngles of the form  $[j/2^t, (j+1)/2^t) \times [a,b]$ , where  $j, t \in \mathbb{N}_0$  and a < b are reals. Then, with probability tending to 1, the range space  $(X, \mathcal{R})$  does not admit an  $\varepsilon$ -net of size at most n/2.