COUNTING PLANE GRAPHS: PERFECT MATCHINGS, SPANNING CYCLES, AND KASTELEYN'S TECHNIQUE

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Definition (sc(N)). For a set S of points in the plane, we denote by C(S) the set of all crossing-free straight-edge spanning cycles of S, and put sc(S) := |C(S)|. Moreover, we let sc(N) = $\max_{|S|=N} \operatorname{sc}(S)$.

Definition $(\operatorname{tr}(N))$. For a set S of points in the plane, we denote by $\mathcal{T}(S)$ the set of all triangulations of S, and put $\operatorname{tr}(S) := |\mathcal{T}(S)|$. Moreover, we let $\operatorname{tr}(N) = \max_{|S|=N} \operatorname{tr}(S)$.

Definition.

Hull edges/vertices: those which are part of the boundary of the convex hull of S Interior edges/vertices: those which are not part of the boundary of $\operatorname{conv}(S)$

h: number of hull vertices

n: number of interior vertices (n := N - h)

Theorem (Already known bounds).

$$tr(N) < 30^{N}$$

$$\operatorname{sc}(N) < 30^{N/4} \cdot \operatorname{tr}(S) \approx 2.3404^{N} \cdot \operatorname{tr}(S)$$

$$\Rightarrow$$
 sc(N) = $O(68.664^N)$

Definition $(sc_{\Delta}(N))$.

 $\mathcal{C}(T)$: set of spanning cycles contained in a triangulation T

 $\operatorname{sc}_{\Delta}(N) := \max_{|S|=N, T \in \mathcal{T}} |\mathcal{C}(T)|$ i. e. maximal number of plane spanning cycles that can be contained in any fixed triangulation of a set of N points in the plane.

Definition (Support). Given a plane edge graph G embedded over a set S of points in the plane, we say that G has a support of x if G is contained in exactly x triangulations of S; we write $\operatorname{supp}(G) = x$.

(1)
$$\operatorname{sc}(S) = \sum_{T \in \mathbb{T}(S)} \sum_{C \in \mathcal{C}(T)} \frac{1}{\operatorname{supp}(C)}$$

Definition (ps-flippable edge). Given a triangulation T, we say that a subset F of its edges is a set of ps-flippable edges, if F are diagonals of interior-disjoint convex polygons whose boundaries are also parts of T.

Lemma (L2.1). Every triangulation T over a set of N points in the plane contains a set F of N/2-2 ps-flippable edges. Also, there are triangulations with no larger sets of ps-flippable edges.

Lemma (L2.2). Consider a triangulation T, a set F of N/2-2 ps-flippable edges in T, and a graph $G \subseteq T$. If G does not contain j edges from F, then $\text{supp}(G) \ge 2^j$.

Definition (Convex decomposition). A conv. decomposition of a point set S is a crossing-free straight-edge graph D on S such that (i) D includes all the hull edges, (ii) each bounded face of D is a convex polygon, (iii) no point of S is isolated in D.

Lemma (L2.3). Let S be a set of points in the plane and let G be a crossing-free straight-edge graph over S that contains all the edges of the convex hull of S. Then G is a convex decomposition of S if and only if every interior vertex of S is valid with respect to G.

Date: 2014-01-02.

Lemma (L2.4). Let c > 1 be a constant such that every set S of an even number of points in the plane satisfies $sc(S) = O(c^{|S|})$. Then $sc(S) = O(c^{|S|})$ also holds for sets S of an odd number of points.

Definition (pm(N)). For a set S of points in the plane, we denote by $\mathcal{M}(S)$ the set of all perfect matchings of S, and put pm(S) := $|\mathcal{M}(S)|$. Moreover, we let pm(N) = $\max_{|S|=N} \text{pm}(S)$. $\mathcal{M}(G)$ is a set of perfect matchings contained in G.

Theorem (T3.1). For any set S of N points in the plane,

$$sc(S) = O(12^{N/4}) \cdot tr(S) = O(1.8613^{N}) \cdot tr(S)$$

Theorem (Kasteleyn's enhanced theorem). Every planar graph G can be oriented onto some digraph \vec{G} such that, for any real-valued weight function μ on its edges, we have

$$\left(\sum_{M \in \mathcal{M}(G)} \mu(M)\right)^2 = |\det(B_{\vec{G},\mu})|$$

Theorem (T4.1). For any set S of N points in the plane,

$$pm(S) \le 8 \cdot (3/2)^{N/4} \cdot tr(S) = O(1.1067^N) \cdot tr(S)$$

Lemma (L5.1). Let T be a triangulation over a set S of $N \ge 6$ points in the plane, such that N is even and S has a triangular convex hull; also, let $v_3(T) = tN$ be the number of interior vertices of degree 3 in T. Then

$$\sum_{C \in \mathcal{C}(T)} \frac{1}{\text{supp}(C)} < 8 \left(\frac{3}{2^t} \left(\frac{(2-t)(2-t/2)}{(1-t)^2} \right)^{1-t} \right)^{N/4}$$

Lemma (L5.2). Let c_{gon} be the maximum real number satisfying the following property: Every simple polygon P that has a triangulation T_P with k flippable and with $l \le k$ of these diagonals forming a ps-flippable set, has at least $2^l c_{gon}^{k-l}$ triangulations. Then $x \le c_{gon} \le 5/4$ with $x \approx 1.17965$ the unique real root of the polynomial $1 + 4x^2 - 4x^3$.

Lemma (L5.3). Consider a triangulation T with the number of flippable edges $\operatorname{flip}(T) = N/2 - 3 + \kappa N$, for some $\kappa \geq 0$, and let x be the unique real root of the polynomial $1 + 4x^2 - 4x^3$. Then

$$\sum_{C \in \mathcal{C}(T)} \frac{1}{\operatorname{supp}(C)} < 8 \left(\frac{(3 + (\gamma^2 - 1)(\kappa + 1/2))(4 + (x^2 - 1)\kappa)}{x^{4\kappa}} \right)^{N/4},$$

where $\gamma = x \cdot e^{-\frac{x^2-1}{4(4+(x^2-1)\kappa)}}$.

Lemma (L5.4). Let c > 1 be a constant such that every set S of an even number of points in the plane and a triangular convex hull satisfies $sc(S) = O(c^{|S|})$. Then $sc(S) = O(c^{|S|})$ also holds for every other finite point-set S in the plane.

Theorem (T5.5). For any set S of N points in the plane,

$$sc(S) = O(10.9247^{N/4}) \cdot tr(S) = O(1.8181^{N}) \cdot tr(S)$$

Corollary (C5.6). $sc(N) = O(54.5430^N)$.