

# 2-Cancellative Hypergraphs and Codes

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### 1 Definitions and notation

**Definition 1.**  $\mathcal{F}$  a family of sets is  $t$ -cancellative if for all  $t + 2$  sets  $A_1, \dots, A_t, B, C \in \mathcal{F}$  such that  $A_i \neq B, A_i \neq C$

$$A_1 \cup A_2 \cup \dots \cup A_t \cup B = A_1 \cup A_2 \cup \dots \cup A_t \cup C \rightarrow B = C.$$

Cancellative means 1-cancellative.

- $c(n)$  – the size of the largest cancellative family on  $n$  elements
- $c(n, r)$  – the size of the largest  $r$ -uniform cancellative family on  $n$  elements
- $f(n, P_1, P_2, \dots)$  – the maximum number of subsets of  $\{1, 2, \dots, n\}$  satisfying properties  $P_1, P_2, \dots$

**Observation 1.**  $c(n + m) \geq c(n)c(m)$

**Definition 2.** Hypergraph  $\mathbb{F} = (V, \mathcal{F})$  is  $(k)$  uniform if each edge has the same number of elements, and linear if  $|E \cap F| \leq 1$  for all edges  $E, F \in \mathcal{F}$ .

**Definition 3.** Associate each subset of  $\mathcal{F}$  to its characteristic vector. If it satisfies, that for each  $a$ -tuple of these vectors at least  $b$  different columns sum up to 1, it is locally  $(a, b)$ -thin.

**Definition 4.**  $\mathcal{F} \subseteq 2^n$  is  $g$ -cover-free if it is locally  $(g + 1, g + 1)$ -thin, i.e. it suffices  $A_0 \not\subseteq \bigcup_{i=1}^g A_i$  for all  $A_0, A_1, \dots, A_g \in \mathcal{F}$ .

- $C_g(n)$  – the size of the largest  $g$ -cover-free  $n$  vertex code
- $C_g(n, r)$  – the size of the largest  $g$ -cover-free  $r$ -uniform hypergraph on  $n$  vertices

**Observation 2.** •  $C_g(n) \leq C_{g-1}(n) \leq \dots \leq C_1(n)$

- $C_g(n, r) \leq C_{g-1}(n, r) \leq \dots \leq C_2(n, r)$
- $C_{t+1}(n, r) \leq c_t(n, r)$  because  $t + 1$ -cover-free  $\iff t$ -cancellative

### 2 Estimates

**Lemma 1.** (D'yachkov, Rykov )  $\exists \alpha_1, \alpha_2$  such that  $\alpha_1 \frac{1}{g^2} < \frac{\log(C_g(n))}{n} < \alpha_2 \frac{\log(g)}{g^2}$

**Theorem 1.**  $\exists \beta_1, \beta_2$  and  $n_0(t)$  such that  $\forall n > n_0(t), t \geq 2$  holds  $\beta_1 \frac{1}{t^2} < \frac{\log(c_t(n))}{n} < \beta_2 \frac{\log(t)}{t^2}$

**Theorem 2.**  $\forall n, k \in \mathbb{Z}: c_2(n, 2k) \leq \frac{\binom{n}{k}}{\frac{1}{2} \binom{2k}{k}}$

**Theorem 3.**  $c_t(n) < \alpha n^{\frac{t-1}{2}} \left( \frac{t+3}{t+2} \right)^n$

**Lemma 2.**  $\mathcal{F}$  an  $r$ -uniform hypergraph, then  $\exists \mathcal{F}^* \subseteq \mathcal{F}$  such that  $\mathcal{F}^*$  is  $r$ -partite and  $|\mathcal{F}^*| \geq \frac{r!}{r^r} |\mathcal{F}|$

**Theorem 4.**  $f_3(n, 7, 4) - \frac{2}{5}n \leq c_2(n, 3) \leq \frac{9}{2}f_3(n, 7, 4) + n$

**Lemma 3.** (F., Frankl)  $i(n, H) = \frac{1}{e(H)} \binom{n}{2} - o(n^2)$  where  $i(n, H)$  is the maximal number of almost disjoint induced copies of  $H$  that can be packed into any  $n$ -vertex graph.

**Theorem 5.**  $c_2(n, 4) = \frac{1}{6}n^2 - o(n^2)$

**Definition 5.**  $\mathcal{P} = \{G_1 = (V_1, E_1), G_2 = (V_2, E_2), \dots\}$  is a packing if all the graphs are edge-disjoint subgraphs of some  $G = (V, \mathcal{E})$ .

An induced packing is called almost disjoint induced packing into  $G$ , if  $|V_i \cap V_j| \leq 2$  for all  $i \neq j$ , i.e. any induced  $G[V_i], G[V_j]$  are vertex disjoint, share 1 vertex, or the set of their intersection is not a subset of any edge of  $G$ .

**Theorem 6.**  $c_2(n, 2k) \geq \frac{n^k}{(2k)^k} - o(n^k)$

**Definition 6.** Symmetric polynomial is defined as  $\sigma_i(X) = \sum_{I \subseteq [s], |I|=i} \prod_{\alpha \in I} x_\alpha$ , where  $X = \{x_1, x_2, \dots, x_s\}$ ,  $0 \leq i \leq s$  and  $\sigma_0(X) = 1$ .

5  $X_1, \dots, X_l$  disjoint,  $|X_j| = t_j, 0 < t_j < k, \sum_j (k - t_j) = k$ , matrix  $M(X_1, \dots, X_l) \dots$

**Observation 3.** •  $\det(M(X_1, \dots, X_l))$  is non-vanishing

- $\mathbf{F}_q^s$  a field,  $q$  power of a prime,  $Z(p) = \{(x_1, \dots, x_s) \in \mathbf{F}_q^s : p(x_1, \dots, x_s) = 0\}$
- $\mathcal{P}_{<k} = \{a_0 + a_1x + \dots + a_{k-1}x^{k-1} : a_i \in F\}$
- $p_Z(x) = \prod_{z \in Z} (x - z)$

**Definition 7.**  $p_1(x), \dots, p_l(x)$  are  $(k_1, \dots, k_l)$ -independent if  $f_1(x)p_1(x) + \dots + f_l(x)p_l(x) = 0, \deg(f_i) < k_i \rightarrow f_i(x) = 0$  for all  $i$ .

- let  $l \geq 2, k_i \in \mathbb{Z}, \sum k_i = k, x_i \in \mathbf{F}_q, 1 \leq i \leq (l-1)k$ , then multiset

$$X_1 = \{x_s : 1 \leq s \leq k - k_1\}, \quad X_j = \{x_s : \sum_{i < j} (k - k_i) < s \leq \sum_{i \leq j} (k - k_i)\}$$

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**Lemma 4.** Polynomials  $p_{X_1}(x), \dots, p_{X_l}(x)$  are  $(k_1, \dots, k_l)$ -independent for all but at most  $\binom{l}{2} q^{(l-1)k-1}$  sequences.

**Lemma 5.**  $\forall k \exists q_0(k) : \text{if } q > q_0(k) \text{ then } \exists S \subseteq \mathbf{F}_q, |S| = 2k \text{ such that polynomials}$

$$p_X(x), p_Y(x), p_W(x)$$

are  $(k - |X|, k - |Y|, k - |W|)$ -independent for each partition

$$S = X \cup Y \cup W,$$

$$1 \leq |X|, |Y|, |W| \leq k, |X| + |Y| + |W| = 2k.$$

**Theorem 7.**  $\forall n \geq r \geq 2$  holds:  $c(n, r) > \frac{\gamma_0}{2^r} \binom{n}{r}$ , where  $\gamma_0 = \prod_{k \geq 1} \frac{2^k - 1}{2^k} = 0,2887 \dots$