

Optimal bounds for sign-representing the intersection of two halfspaces by polynomials

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The treshold degree of a function $f : \{0, 1\}^n \rightarrow \{-1, 1\}$ is the least degree of real polynomial p with $f(x) \equiv \text{sgn}p(x)$, denoted $\text{deg}_\pm(f)$.

Theorem 1 (Main theorem) *For $n = 1, 2, 3, \dots$, let $D(n)$ denote the maximum treshold degree of a function of the form $f(x) \wedge g(x)$, where $f, g : \{0, 1\}^n \rightarrow \{-1, 1\}$ are halfspaces. Then $D(n) = \Theta(n)$.*

We give a randomized algorithm which constructs two halfspaces on $\{0, 1\}^n$ whose intersection has treshold degree $\Theta(n)$.

Tools

The binary entropy function $H(p) = -p \log p - (1-p) \log(1-p)$ ($H : [0, 1] \rightarrow [0, 1]$) is strictly increasing on $[0, 1/2]$. Fact:

$$\sum_{i=0}^k \binom{n}{i} \leq 2^{H(k/n)n}, \quad k = 0, 1, 2, \dots, \lfloor n/2 \rfloor.$$

Fourier transform.

- inner product: $\langle f, g \rangle = 2^{-n} \sum_{x \in \{0,1\}^n} f(x)g(x)$
- $\chi_S : \{0, 1\}^n \rightarrow \{-1, 1\}$: $\chi_S(x) = (-1)^{\sum_{i \in S} x_i}$ yields an orthonormal basis
- a unique representation: $f = \sum_{S \subseteq \{1, \dots, n\}} \hat{f}(S) \chi_S$, where $\hat{f}(S) = \langle f, \chi_S \rangle$
- Parseval's identity: $\sum_{S \subseteq \{1, \dots, n\}} \hat{f}(S)^2 = \langle f, f \rangle$

Analysis of random halfspaces

Lemma 2 *Let $f, g : \{0, 1\}^n \rightarrow \{0, 1\}$ be given functions. Fix an integer k with $0 \leq k \leq n/2$. For a set $S \subseteq \{1, \dots, n\}$, define $F_S : \{0, 1\}^n \rightarrow \{0, 1\}$ by*

$$F_S(x) = f(x) \wedge \left(g(x) \oplus \bigoplus_{i \in S} x_i \right).$$

Fix a real $\zeta > 0$. Then with probability at least $1 - 2^{-n+H(k/n)n+2\zeta n}$ over a uniformly random choice of $S \in \mathcal{P}(\{1, 2, \dots, n\})$, one has

$$\left| \hat{F}_S(T) - \frac{1}{2} \hat{f}(T) \right| \leq 2^{-\zeta n-1}, \quad |T| \leq k.$$

Lemma 3 Fix an integer $k \geq 0$ and reals $\varepsilon, \zeta \in (0, 1/2)$. Choose sets $S_0, S_1, \dots, S_k \in \mathcal{P}(\{1, 2, \dots, n\})$ uniformly at random. Fix any integer s and define $f : \{0, 1\}^n \rightarrow \{0, 1\}$ by

$$f(x) = 1 \Leftrightarrow \sum_{i=0}^k 2^i \sum_{j \in S_i} x_j \equiv s \pmod{2^{k+1}}.$$

Then with probability at least $1 - (k+1)2^{-n+H(\varepsilon)n+2\zeta n}$ over the choice of S_0, S_1, \dots, S_k , one has

$$\left| \hat{f}(T) - \frac{\delta_{T, \emptyset}}{2^{k+1}} \right| \leq 2^{-\zeta n}, \quad |T| \leq \varepsilon n.$$

Theorem 4 (Key property of random halfspaces) Fix an integer $k \geq 0$ and reals $\varepsilon, \zeta \in (0, 1/2)$. Choose integers w_1, w_2, \dots, w_n uniformly at random from $\{0, 1, \dots, 2^{k+1} - 1\}$. For $s \in \mathbb{Z}$, define $f_s : \{0, 1\}^n \rightarrow \{0, 1\}$ by

$$f_s(x) = 1 \Leftrightarrow \sum_{i=0}^n w_i x_i \equiv s \pmod{2^{k+1}}.$$

Then with probability at least $1 - (k+1)2^{-n+H(\varepsilon)n+2\zeta n+k+1}$ over the choice of w_1, w_2, \dots, w_n , one has

$$\left| \hat{f}_s(T) - \frac{\delta_{T, \emptyset}}{2^{k+1}} \right| \leq 2^{-\zeta n}, \quad |T| \leq \varepsilon n, \quad s \in \mathbb{Z}.$$

Zeroing out correlations by a change of distribution

$f, g : X \rightarrow (R)$, X finite, then $\langle f, g \rangle = \frac{1}{|X|} \sum_{x \in X} f(x)g(x)$.

Theorem 5 Let $f, \chi_1, \dots, \chi_k : X \rightarrow \{-1, +1\}$ be given functions on a finite set X . Suppose that

$$\sum_{i=1}^k |\langle f, \chi_i \rangle| < \frac{1}{2},$$

$$\sum_{j=1, j \neq i}^k |\langle \chi_i, \chi_j \rangle| \leq \frac{1}{2}, \quad i = 1, 2, \dots, k.$$

Then there exists a probability distribution μ on X such that

$$\mathbb{E}_\mu[f(x)\chi_i(x)] = 0, \quad i = 1, 2, \dots, k.$$

Theorem 6 Let $\alpha > 0$ be a sufficiently small absolute constant. Choose integers w_1, w_2, \dots, w_n uniformly at random from $\{0, 1, \dots, 2^{\lfloor \alpha n \rfloor + 1} - 1\}$. For $s \in \mathbb{Z}$, define

$$X_s = \left\{ x \in \{0, 1\}^n : \sum_{i=1}^n w_i x_i \equiv s \pmod{2^{\lfloor \alpha n \rfloor + 1}} \right\}.$$

Then with probability at least $1 - e^{-n/3}$ over the choice of w_1, w_2, \dots, w_n , there is a distribution μ_s on X_s (for each s) such that

$$\mathbb{E}_{\mu_s}[p(x)] = \mathbb{E}_{\mu_t}[p(x)]$$

for any $s, t \in \mathbb{Z}$ and any polynomial p of degree at most $\lfloor \alpha n \rfloor$.

Reduction to a univariate problem

Another tools: rational approximation.

- $\deg p(x)/q(x) := \max\{\deg(p), \deg(q)\}$, where p, q are polynomials on \mathbb{R}^n
- $f : X \rightarrow \{-1, +1\}$, where $X \subseteq \mathbb{R}^n$. For $d \geq 0$ define

$$R(f, d) = \inf_{p, q} \sup_{x \in X} \left| f(x) - \frac{p(x)}{q(x)} \right|,$$

where the infimum is over all polynomials p, q of degree up to d such that $q|_X \neq 0$.

- $R^+(f, d) = \inf_{p, q} \sup_{x \in X} \left| f(x) - \frac{p(x)}{q(x)} \right|$ where q is positive on X
- $R^+(f, 2d) \leq R(f, d) \leq R^+(f, d)$
- $S \subseteq \mathbb{R}$: $R^+(S, d) = \inf_{p, q} \sup_{x \in S} \left| \operatorname{sgn} x - \frac{p(x)}{q(x)} \right|$

Theorem 7 (Sherstov) Let n, d be positive integers, $R = R^+(\{\pm 1, \pm 2, \dots, \pm n\}, d)$. For $1 \leq d \leq \log n$,

$$\exp \left\{ -\Theta \left(\frac{1}{n^{1/(2d)}} \right) \right\} \leq R < \exp \left\{ -\frac{1}{n^{1/d}} \right\}.$$

For $\log n < d < n$,

$$R = \exp \left\{ -\Theta \left(\frac{1}{\log(2n/d)} \right) \right\}.$$

For $d \geq n$, $R = 0$.

Theorem 8 (Sherstov) Let $f : X \rightarrow \{-1, +1\}$ and $g : Y \rightarrow \{-1, +1\}$ be given functions, where $X, Y \subseteq \mathbb{R}^n$ are arbitrary finite sets. Assume that f and g are not identically false. Let $d = \deg_{\pm}(f \wedge g)$. Then

$$R^+(f, 4d) + R^+(g, 2d) < 1.$$

Proposition 9 Let n_1, \dots, n_k be positive integers, $|x| := x_1 + x_2 + \dots + x_n$. Consider a function $F : \{0, 1\}^{n_1} \times \dots \times \{0, 1\}^{n_k} \rightarrow \{-1, +1\}$ such that $F(x_1, \dots, x_k) \equiv f(|x_1|, \dots, |x_k|)$ for some $f : \{0, 1, \dots, n_1\} \times \dots \times \{0, 1, \dots, n_k\} \rightarrow \{-1, +1\}$. Then for all d , $R^+(F, d) = R^+(f, d)$.

Theorem 10 (Reduction to a univariate problem) Put $k = \lfloor \alpha n \rfloor$, where $\alpha > 0$ is the absolute constant from Theorem 6. Choose $w_1, w_2, \dots, w_n \in \{0, 1, \dots, 2^{\lfloor \alpha n \rfloor + 1} - 1\}$ uniformly at random. Define $f : \{0, 1\}^n \times \{0, 1, 2, \dots, n\} \rightarrow \{-1, +1\}$ by

$$f(x, t) = \operatorname{sgn} \left(\frac{1}{2} + \sum_{i=1}^n w_i x_i - 2^{k+1} t \right).$$

Then with probability at least $1 - e^{-n/3}$ over the choice of w_1, w_2, \dots, w_n , one has

$$R^+(f, d) \geq R^+(\{\pm 1, \pm 2, \dots, \pm 2^k\}, d), \quad d = 0, 1, \dots, k.$$

Theorem 11 Put $k = \lfloor \alpha n \rfloor$, where $\alpha > 0$ is the absolute constant from Theorem 6. Choose w_1, w_2, \dots, w_n uniformly at random from $\{0, 1, \dots, 2^{\lfloor \alpha n \rfloor + 1} - 1\}$. Define $f : \{0, 1\}^{2n} \rightarrow \{-1, +1\}$ by

$$f(x) = \operatorname{sgn} \left(\frac{1}{2} + \sum_{i=1}^n w_i x_i - 2^{k+1} \sum_{i=n+1}^{2n} x_i \right).$$

Then with probability at least $1 - e^{-n/3}$ over the choice of w_1, w_2, \dots, w_n , one has

$$R^+(f, d) \geq R^+(\{\pm 1, \pm 2, \dots, \pm 2^k\}, d), \quad d = 0, 1, \dots, k.$$

Now it is easy to prove the main result.

Theorem 12 (Main result) Fix sufficiently small absolute constants $\alpha > 0$ and $\beta = \beta(\alpha) > 0$. Choose integers $w_1, w_2, \dots, w_n \in \{0, 1, \dots, 2^{\lfloor \alpha n \rfloor + 1} - 1\}$ uniformly at random. Then with probability at least $1 - e^{-n/3}$, the function $f : \{0, 1\}^{2n} \rightarrow \{-1, +1\}$ given by

$$f(x) = \operatorname{sgn} \left(\frac{1}{2} + \sum_{i=1}^n w_i x_i - 2^{\lfloor \alpha n \rfloor + 1} \sum_{i=n+1}^{2n} x_i \right)$$

obeys $\deg_{\pm}(f \wedge f) \geq \lfloor \beta n \rfloor$.